

Huffman Coding :

- It is an entropy based coding technique.
- An assumption is made and only one code is encoded at a time.
- It is an optimal coding method.
- It has a variable length output.

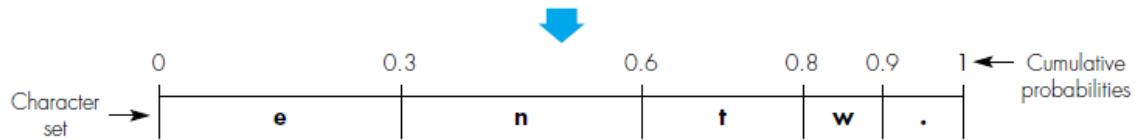
Arithmetic Coding :

- It is also an entropy based coding.
- No assumptions are made and based on probability of chances, the codes are encoded.
- All the codes are encoded with arithmetic values from 0 to 1.
- It is a much more slower technique than Huffman coding.
- The efficiency of arithmetic coding is higher than Huffman coding.

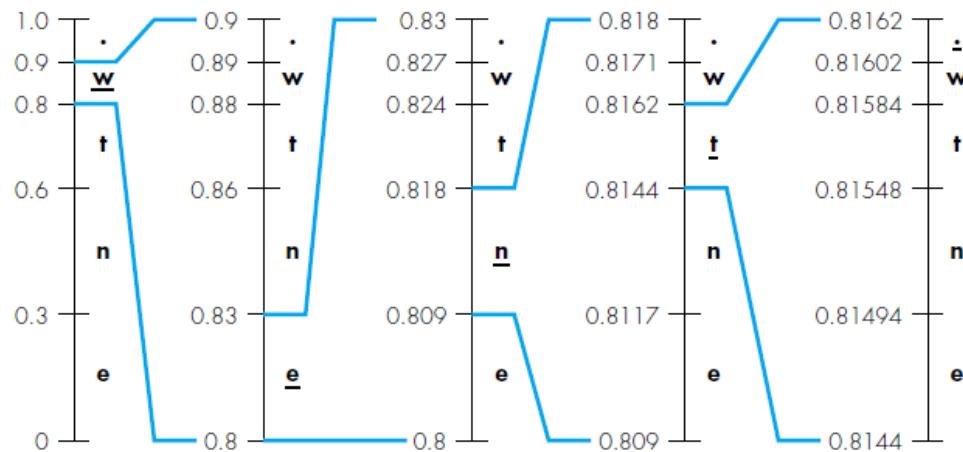
(a)

Example character set and their probabilities:

$$e = 0.3, n = 0.3, t = 0.2, w = 0.1, . = 0.1$$



(b)



Encoded version of the character string **went.** is a single codeword in the range $0.816\ 02 \leqslant \text{codeword} < 0.8162$

2.

So the Probability of code word in terms of accuracy is ,

$$0.8160 < \text{codeword} < 0.8162$$

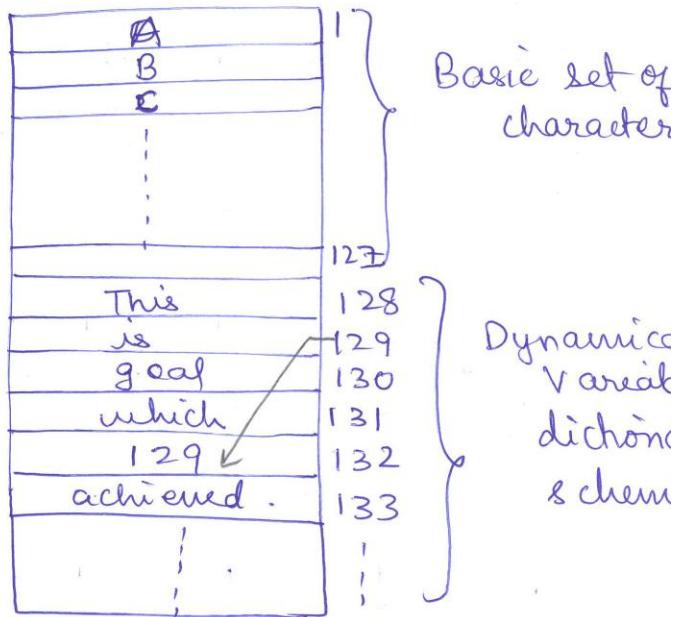
Arithmetic code word .

Q-6

LZW Algorithm

Lempel Ziv Welsch algorithm is designed with a dynamically changing dictionary and is used for compression of images.

LZW scheme





$$\begin{aligned}x &= 0 \\y &= 1 \\z &= 2\end{aligned}$$

- In this scheme, initially the code words are assigned with values such as $x=0, y=1, z=2$.
- $s \rightarrow s$ is the first character in the string
 $c \rightarrow$ next character in the string.
- $s = s + c$ If code word already exists
 $s = c$ If code word does not exist.

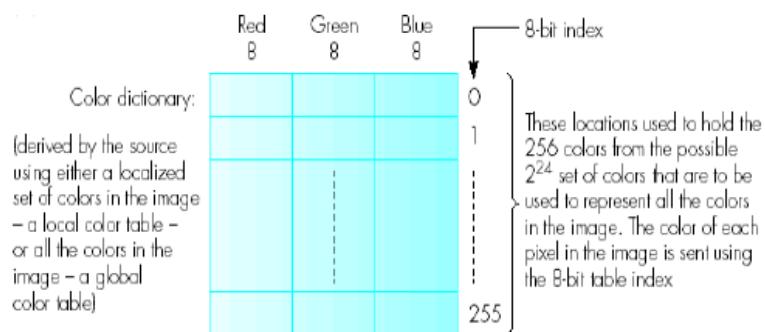
String: XY XY XY X YZ XY XY XX.

s	c	old	new	string table
X	Y	0	4	XY
Y	X	1	5	YX
X	Y			
XY	Y	4	6	XYY
Y	X			
YX	Y	5	7	YXY
Y	Z	1	8	YZ
Z	X	2	9	ZX
X	Y			
XY	X	4	10	XYX
X	Y			
XY	Y			
XYY	X	6	11	XYYX
X	EOF			

The output code for LZW algorithm
is 01451246 //

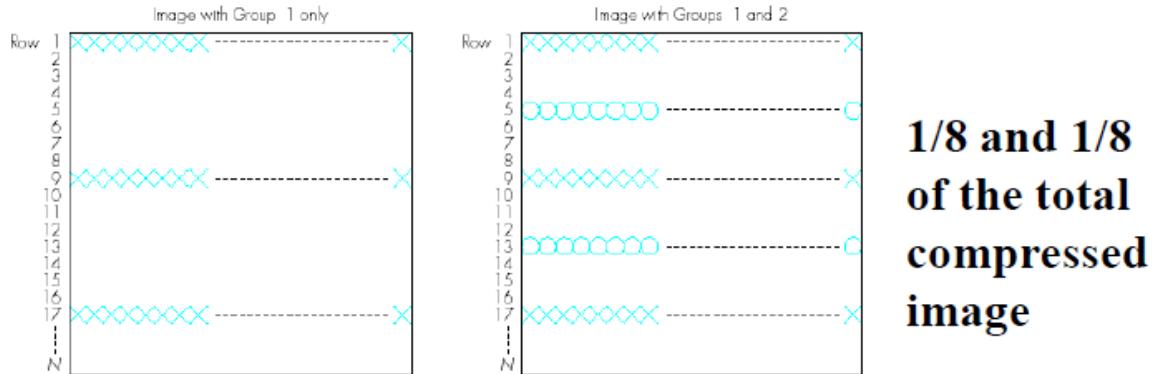
2b.

Image Compression – GIF compression Principles



- The graphics interchange format is used extensively with the Internet for the representation and compression of graphical images

Image Compression – GIF interlaced mode



- GIF also allows an image to be stored and subsequently transferred over the network in an ***interlaced mode***; useful over either low bit rate channels or the Internet which provides a *variable transmission rate*

Image Compression – GIF

- Although colour images comprising 24-bit pixels are supported GIF reduces the number of possible colours that are present by choosing 256 entries from the original set of 2^{24} colours that match closely to the original image
- Hence instead of sending as 24-bit colour values only 8-bit index to the table entry that contains the closest match to the original is sent. This results in a 3:1 compression ratio
- The contents of the table are sent in addition to the screen size and aspect ratio information
- The image can also be transferred over the network using the interlaced mode

Q-1

JPEG Encoder

Entropy Encoding

Quantization

Block Diagram

Image/Blok preparation

Tables

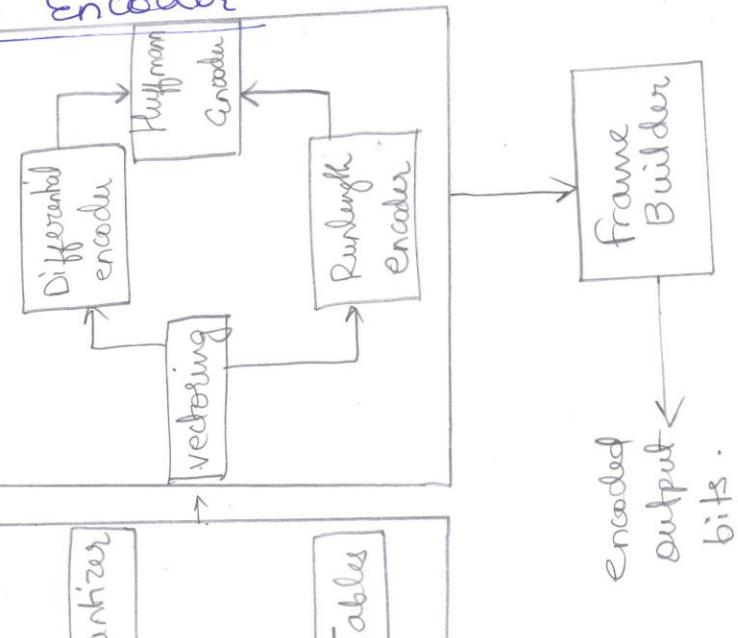
Quantizer

Block Preparation

Forward DCT

Source Image

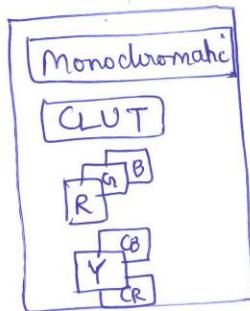
Image Preparation



→ The main components of a JPEG are :

- ① Image / Block Preparation
- ② Forward DCT
- ③ Quantization
- ④ Entropy Encoding
- ⑤ Frame Building.

① Image / Block Preparation :



In this block, the input about the image is converted as CLUT, RGB and YCrCb for the preparation.

② Forward DCT

→ The image data is transformed using the Forward DCT function.

$$F[i, j] = \frac{1}{4} c(i)c(j) \sum_{x=0}^7 \sum_{y=0}^7 P[x, y] \cdot \cos \left(\frac{2\pi}{16} xy \right)$$

where $P[x, y]$ is the input for 8×8 and $F[i, j]$ is the transformed output.

The DCT coefficients are then processed before quantization.

DCT coefficient	Difference	SSS value
13	13	4
12	-1	1
14	2	2
13	-1	1

The SSS value is the no. of bits used to represent the difference in the DCT Coefficients.

③ Quantisation :

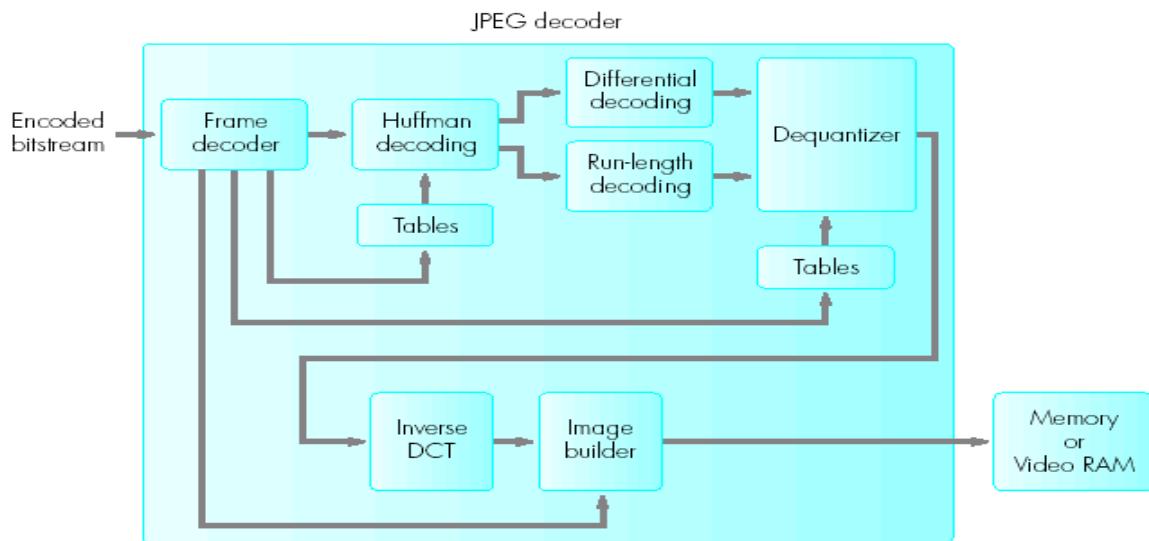
- The 64 DCT coefficient must be quantised in order to get a compressed image.
- Thus discrete levels are used for quantization specified in the quantization table.
- After all the coefficients are quantised, it goes into the entropy encoding unit.

④ Entropy Encoding :

Before three levels of

0	(0,0)	(q1)	(0,7)
1	(1,0)					
1						
1						
1						
7	(7,1)					(7,7)

encoding, vectoring
vector table with 64
with index running
{ 63, 62, 61, ..., }



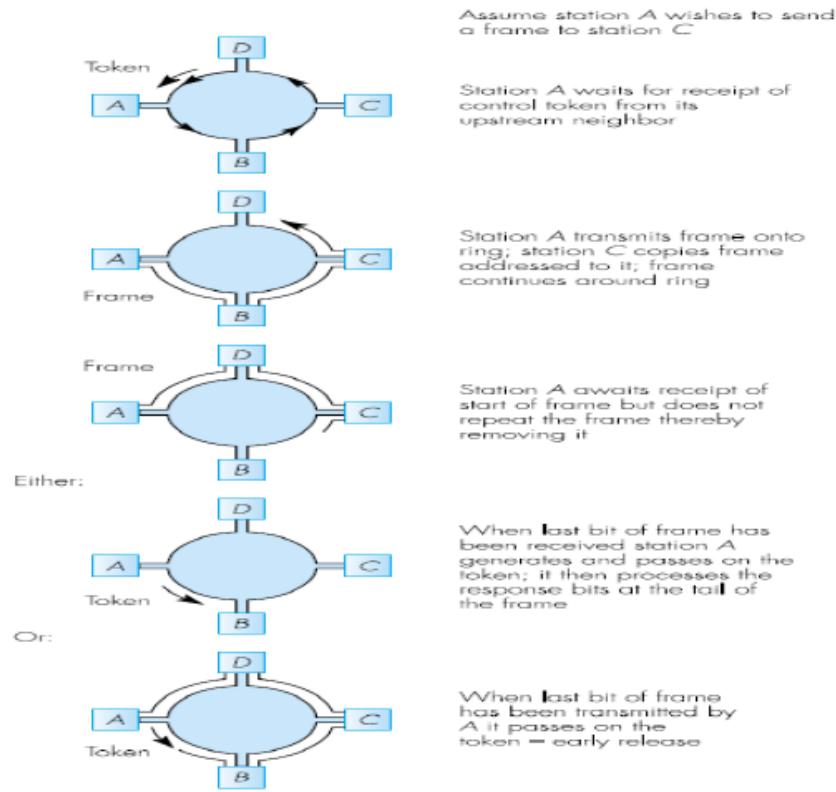
JPEG decoding

- The two decompressed streams containing the DC and AC coefficients of each block are then passed to the differential and run-length decoders
- The resulting matrix of values is then dequantized using either the default or the preloaded values in the quantization table
- Each resulting block of 8X8 spatial frequency coefficient is passed in turn to the **inverse DCT** which in turn transforms it back to their spatial form
- The image builder then reconstructs the image from these blocks using the control information passed to it by the frame decoder

Token ring

- All the stations are connected together by a set of unidirectional links in the form of a ring and all frame transmissions between any of the stations take place over it by circulating the frame around the ring
- Only one frame transfer can be in progress over the ring at a time

Token ring network operation



Token ring wiring configuration

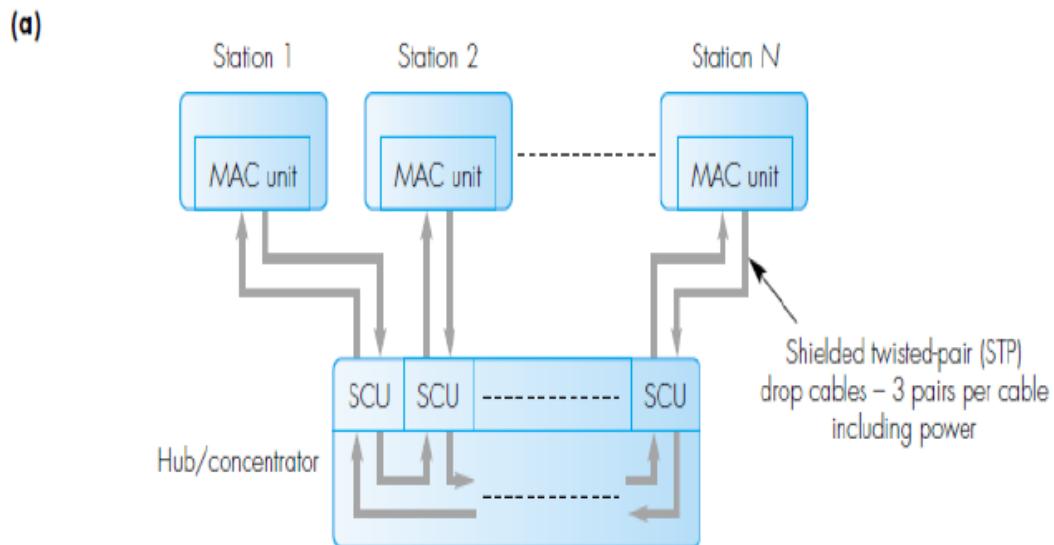
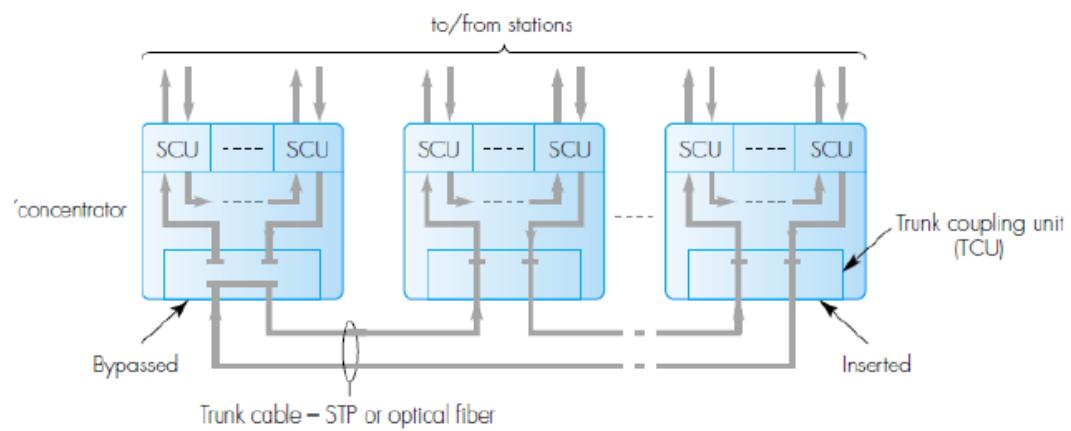
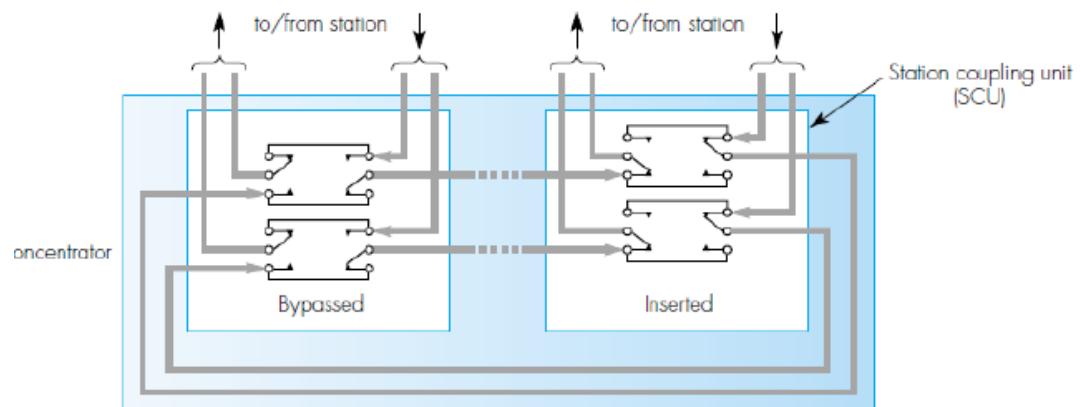


Figure 8.6 Token ring wiring configurations: (a) single hub; (b) station coupling unit; (c) multiple hubs/concentrators.



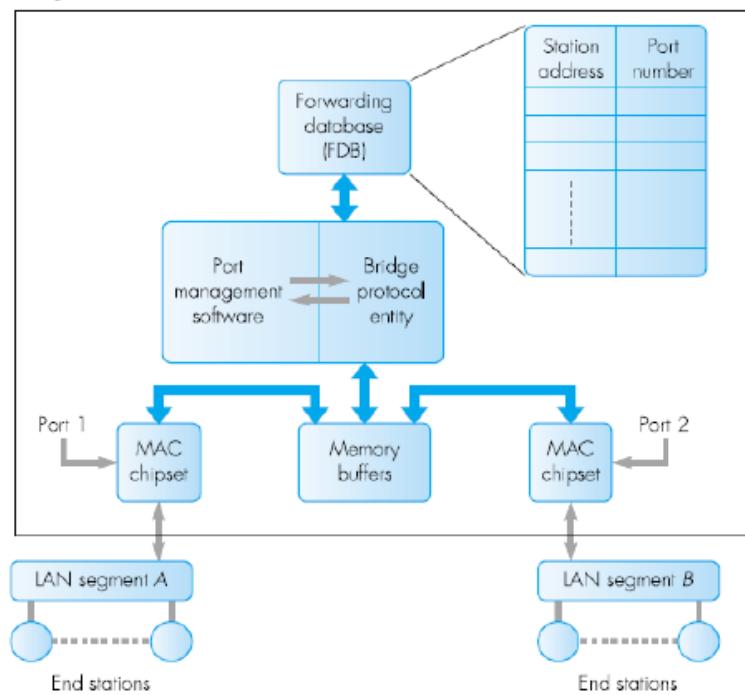
5.

8.5.1:Transparent bridges

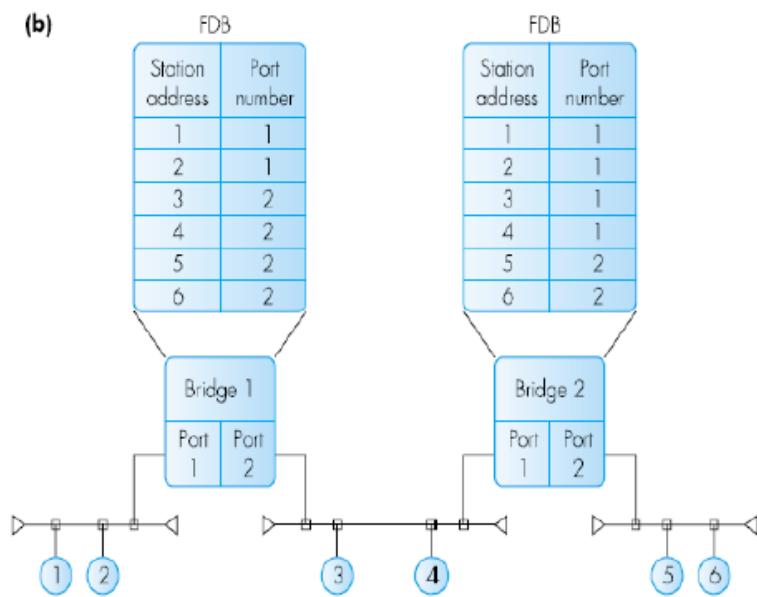
- With a transparent bridge, as with a repeater, the presence of one (or more) bridges in a route between two communicating stations is transparent to the two stations . All routing decisions are made exclusively by the bridge(s)
- Fig 8.12
- Two steps- frame forwarding(filtering), learning
- A bridge maintains a forwarding database(routing directory)
- Bridge learning
 - Forwarding database to be created in advance

Figure 8.12 Transparent bridge schematic: (a) architecture; (b) application example.

(a) Bridge



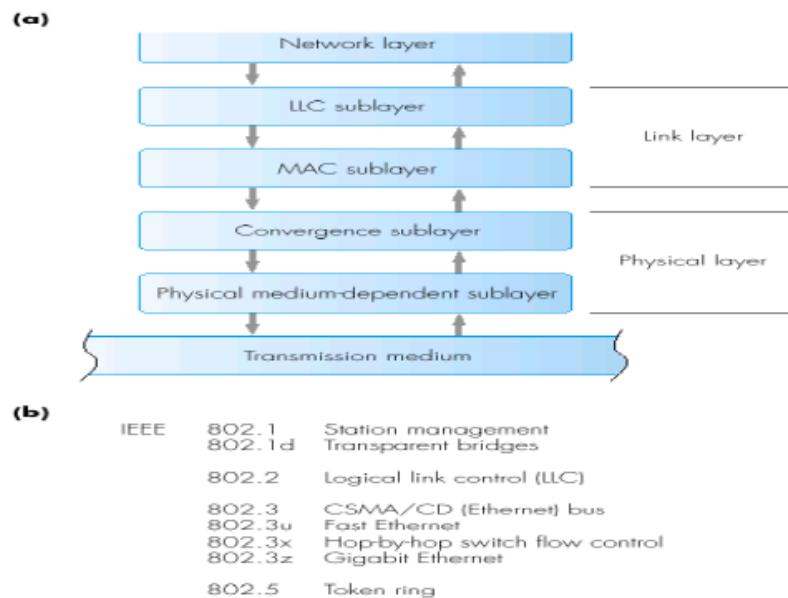
(b)



6.

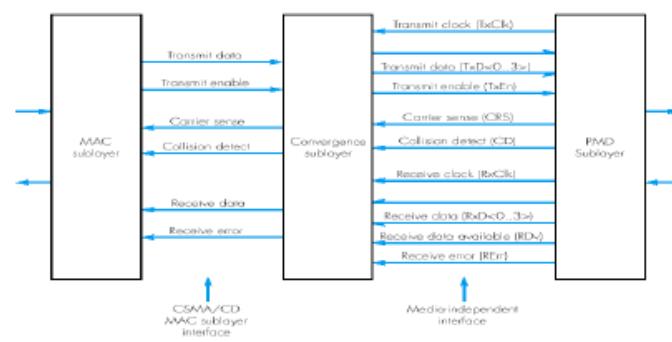
8.8:LAN protocol

Figure 8.31 LAN protocols: (a) protocol framework; (b) examples.



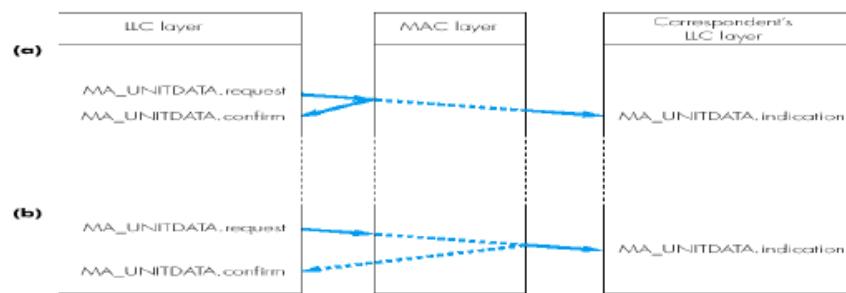
8.8.1:Physical layer

Figure 8.32 Fast Ethernet media-independent interface.



8.8.2:MAC sublayer

Figure 8.33 MAC user service primitives: (a) CSMA/CD; (b) token ring.

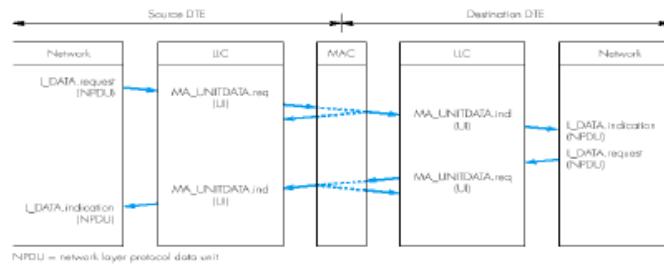


MAC Sub Layer

- Standard service primitives are
- MA_UNITDATA.request - includes required destination address, service data unit and the required class of service
- MA_UNITDATA.indication
- MA_UNITDATA.confirm - includes a parameter that specifies a success or failure of data primitive

8.8.3:LLC sublayer

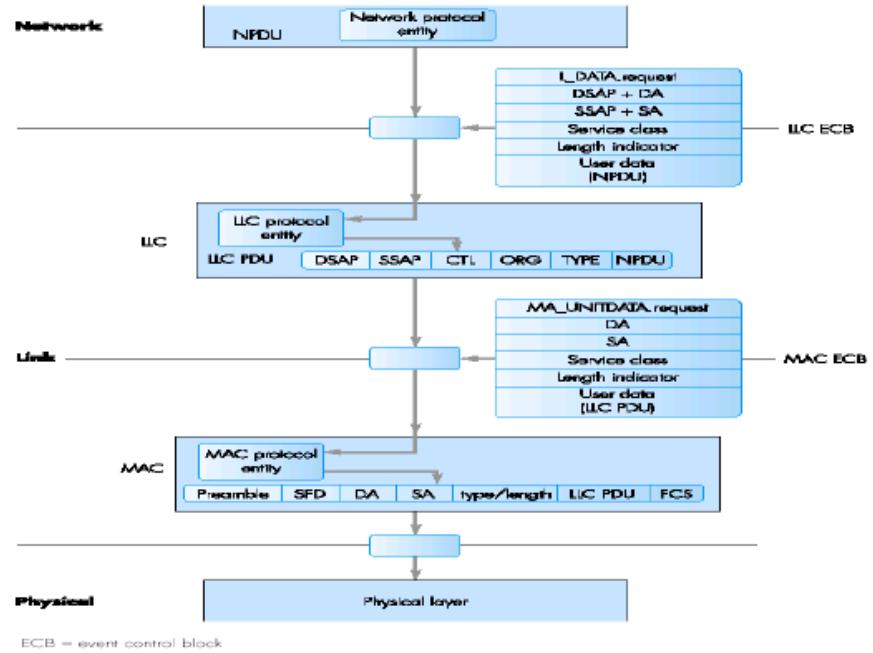
Figure 8.34 LLC/MAC sublayer interactions.



LLC Sub Layer

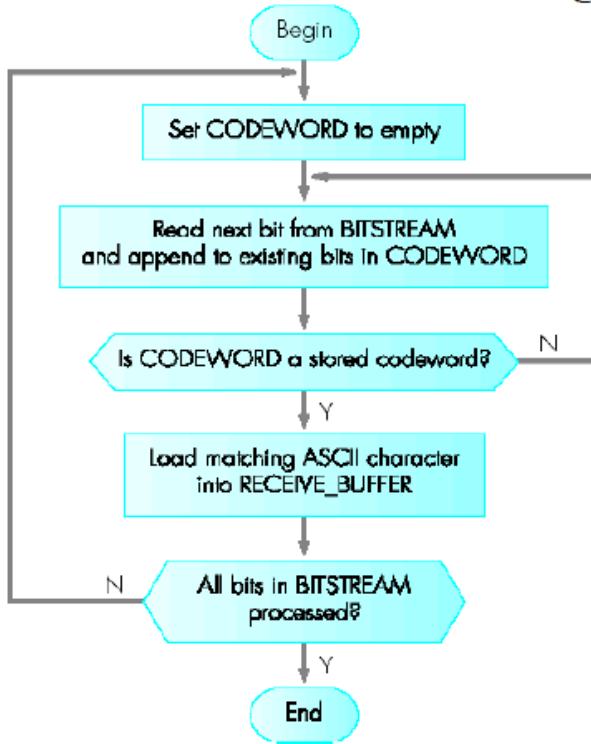
- LLC protocol is based on high level data link control (HDLC) protocol and supports both connectionless and connection oriented mode
- In the connectionless mode, the service primitive is L_DATA.request and data is transferred in an unnumbered information frame
- The parameters are source and destination address and the user data
- For internet applications, two field form is called subnet access protocol - SNAP header

Figure 8.35 Interlayer primitives and parameters.

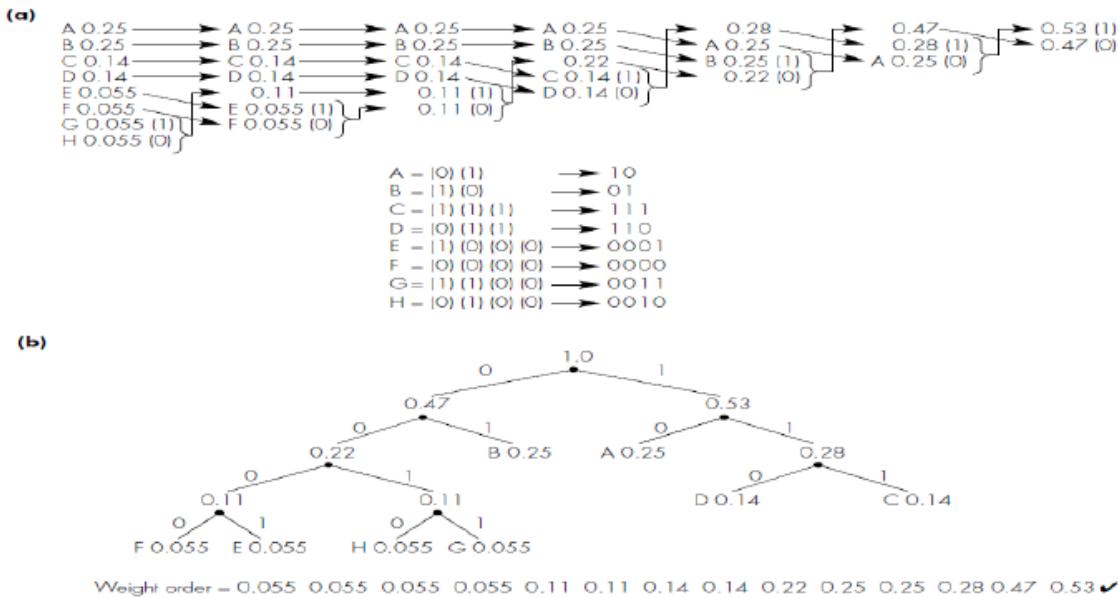


- The interaction between LLC and MAC is shown
- LLC sublayer reads the destination on source LLC service access point address, DSAP and SSAP from event control block ECB and adds at the head of LLC PDU with 8 bit control field, type field + LPDU and passed to MAC sublayer
- The service class used by the MAC sublayer protocol entity defines the priority for the frame
- On receipt of the request, the MAC protocol entity creates a frame ready for transmission on the link

Text Compression – Flow chart of a suitable decoding algorithm



Decoding of received bitstream assuming codewords derived: decoding algorithm



Ex 3.3

AABBCAAD

A=2bits

B=2bits

C=3bits

D=3bits

$$\begin{aligned} \text{No .of bits} &= 2 \times 2 + 2 \times 2 + 3 \times 1 + 2 \times 1 + 3 \times 1 \\ &= 16 \text{bits} \end{aligned}$$