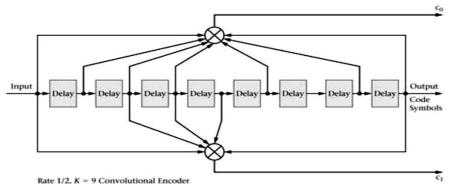
CMR INSTITUTE OF USN						AND TEARS
TECHNOLOGY						* CMR INSTITUTE OF TECHNOLOGY, BENGALURU. ACCREDITED WITH A+ GRADE BY NAAC

Date:	21/05/20	18	Duratio		ins Ma	ax Mark			1		1			
							ss:	50	Sem:	: 8th	Branch:	I	ECE (A	,B)
				Ans	wer Any	FIVE I	FULL	Ques	tions			·		
												Marks	OE	
												IVIAIKS	CO	RB'
	Draw and d n 802.11. A				-		mana	ageme	nt MA	C fram	e format	[10]	CO5	L
	The IEEE 802 this standard i MAC frame of a frame check tion, address, frame type. T structure of a address 3, see an individual 2 (Bytes) Frame Control	is able to pro- consists of the k sequence (and sequence) The FCS con a MAC fram quence cont frame field 2 Duration/ ID	operly con ne followin (FCS). The ce control tains an IE ne format a rol, address , there typi	Istruct fram Ing basic con- e MAC hea- information EEE 32-bit co- and a mana ss 4, and fra- ically exist s	mes for fran mponents: ader consiss n. The fran cyclic redu agement M ame body a	a MAC asts of sev ne body o indancy c IAC fram are only p hat are us	heade weral fi contain code (C ne exa presen sed to p	r, a vari fields in ns infor CRC). F imple. T at in cer	iable len cluding mation t Figure 9– The field tain type	agth frame frame contracts frame contracts frame contracts frame -8 shows Is labeled es of fram hal inform	e body, and ntrol, dura- ecific to the the general address 2, nes. Within nation.	2m		
	7			Manageme	ent MAC Fra	ıme Forma	ıt							
	2 (Bytes) Frame	2 Duration	6	6	6	2 Sequenc	e	0-2312	1	4				
	Control		DA	SA	BSSID	Control		Frame	Body	FC	5	3m		
Red	•	/	MAC 1				→							
F	igure 9–8 Exa	amples of IEE	E 802.11 M	AC frame fo	ormats (Cou	rtesy of IE	EEE).							
	2 (Bits)	2	(95)) (95))	4	1	1	1	1	1	1 1	1			
	Protocol Version	Туре		Subtype	То	From	More	Retry	Pwr Mo	ore WFP				
			19. 20 10 10 10	Markel Bar	DS	DS	Frag		Mgt Da	ata	1.00	2m		
	-		1	Frame Contro	ol Field (2 D.	tes _ 16 D	te) -			/	1	2111		

Figure 9-9 shows the structure of the frame control field (i.e., the first 2 bytes of the MAC frame). As one can see, further information can be encoded into the control frame subfields that can even consist of 1bit fields. For further details of the meanings and possible encodings for these fields one should look at the most recent version of the IEEE 802.11 standard. This work will not go into that fine amount of detail. Returning to the general MAC frame format shown by Figure 9-8, a few comments about the address, sequence, and frame body fields are appropriate here. The four address fields in the MAC frame format are used to indicate the basic service set identifier (BSSID), destination address (DA), source address (SA), receiver address (RA), and transmitter address (TA) (although not all at the same time). Furthermore, some types of MAC frames may not contain some of the address fields just mentioned. Each address field is 48 bits in length and can therefore use 48-bit IEEE 802 MAC addresses to indicate an individual station on the network or a group address. The group address can be one of two types, either a multicast group or a broadcast group (i.e., all of the stations presently active in the wireless LAN). The BSSID field is used to uniquely identify each BSS. For a typical wireless LAN, the value of this field is the MAC address currently in use by the 3m station portion of the AP or APs of the WLAN. The sequence field consists of 16 bits that are composed of two subfields of 4 bits and 12 bits. The 12 bit field provides a sequence number for each MSDU and the 4-bit field provides a MSDU fragment number, if needed. The frame body field has a minimum length of 0 bytes and as shown in the figure can be as long as 2312 bytes. [10] **CO**4 L2 2. Discuss various coding techniques used in wireless communication. Soln: Wireless radio channel is most <u>unreliable and random characteristics channels</u>. Hence it is necessary to make the signal more robust before it is transmitting through wireless channels. At transmitter increase the transmitted signal's immunity to radio channel noise and other channel impairments like fading and multipath spread. In digitally based systems, techniques correspond to an attempt to realize a reduction in bit errors and frame errors. The best strategy is to employ some form of error detection and correction codes to reduce the required number of requests for retransmission by the system when errors cannot be corrected. Error correction may generally be realized in two different ways: ■ Automatic repeat request (ARQ) (sometimes also referred to as backward error correction): This is an error control technique whereby an error detection scheme is combined with requests for retransmission of erroneous data. Every block of data received is checked using the error detection code used, and if the check fails, retransmission of the data is requested - this may be done repeatedly, until the data can be verified. ■ Forward error correction (FEC): The sender <u>encodes the data</u> using an error-correcting code (ECC) prior to transmission. The additional information (redundancy) added by the code is used by the receiver to 1mrecover the original data. Different codes are used to enhance the transmission of packet data over wireless systems. **Block codes** • **Convolutional codes Turbo codes Speech coders** Block codes: System takes a block of data bits and encodes them into another block of bits with some additional bits that are used to detect or combat errors. Block codes are processed on a block-by-block basis.

- Using more sophisticated techniques, additional bits may be generated through a <u>matrix or polynomial generator</u> and added to the original block of bits to form a <u>codeword</u> that will be eventually transmitted by the system.
- A codeword generated by a polynomial is a form of cyclic code codes of this type are known as <u>Cyclic Redundancy Check (CRC) codes.</u>
- Consequently, the block coder is a *memoryless* device
- In block coding, divide message into <u>blocks</u>, each of k bits, called <u>datawords</u> and add 'r' <u>redundant bits</u> to each block to make the length n = k + r. The resulting n-bit blocks are called <u>codeword</u>
- <u>Additional bits 'r'</u> may be generated through <u>a matrix or Polynomial</u> <u>generator (eg. CRC code)</u> and added to the original block of bits to form a <u>codeword</u> that will be eventually transmitted by a system.
- Depending upon the type of coding level employed these schemes can <u>both</u> <u>detect and correct</u> limited numbers of errors.
 Convolutional codes
- Convolutional codes are applied in applications that require good performance with low implementation complexity.
- They operate on code streams (not in blocks)
- It maps information to code bits sequentially by convolving a sequence of information bits with "generator".
- k & n are very small (usually k=1-3, n=2-6)
- Input depends not only on current set of k input bits, but also on past input.
- In cdma2000 system a convolutional encoder with R=1/3 and K=9 is used.
- In practice, the use of convolutional encoders provides better FEC capabilities than available from block codes.
- Figure shows in block diagram form an implementation of a convolutional encoder (with K=9 and R=1/2) specified for use in cdma2000



Turbo encoder

- Turbo encoders are a modified form of combined convolutional encoders that can be used to create a new class of <u>enhanced error correction codes</u>.
- It is constructed from <u>two systematic</u>, <u>recursive convolutional encoders</u> connected in parallel with an <u>interleaver</u> preceding the input to the second convolutional encoder.
- The output bit steams of the two convolutional encoders are <u>multiplexed</u> <u>together</u> and repeated to form the final code symbols.
- For cdma2000, Rate 1/2, 1/3, 1/4 and 1/5 turbo encoders are employed instead of convolutional encoders for various higher-bit transfer rates and radio configurations

Speech coders:

The speech coders used for both GSM and CDMA wireless system. Speech coder take 20-msec segments and process it into lower-bit-rate digitally encoded speech in preparation for its transmission over the air interface

3m

2m

1. Wavefor		ch coders:			
		hat is called "waveform" coding			
		form of the original voice			
	veform.	h a compression ratio as the			
	er category of speech c	h a compression ratio as the			
	ample PCM at the 64kb				
2. Vocoder	1	ps data fate.			
		ec that analyzes and synthesizes			
	n voice signal for audi		2m		
	0	voice transformation etc			
_		eech by measuring how its spectral			
character	istics change over time				
		citation-Long Term Prediction (RPE-LTP)			
		reduce the amount of data sent between			
the	mobile station (MS) ar	nd base transceiver station (BTS).			
1 • 1 • 0		r.	[6]		T
xplain briefly ser	vice provided by GSM	l.	[6]	CO3	L
States Internet	CSM Talana	to esta contrate facility resident			
	GSM Telese	tvices			
	GSM Bearer S				
	GSWI bearer S	bervices			
· · · · · · · · · · ·	GSM Bearer S	Services			
A CONTRACTOR OF THE OWNER					
End User	GSM	Transmission Fnd User			
End User					
End User	GSM	Transmission Fnd User	2m		
End User	GSM	Transmission Fnd User	2m		
Relationship of teleservices	GSM	Transmission Fnd User	2m		
Relationship of teleservices and bearer	GSM	Transmission Fnd User	2m		
Relationship of teleservices and bearer services to the GSM system	GSM	Transmission Fnd User	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of	GSM	Transmission Fnd User	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI).	GSM Network	Transmission Networks End User	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of	GSM Network	Transmission Fnd User	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI).	GSM Network	Additional Details Full rate at 13 kbps voice	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI).	GSM Network Service Telephony Emergency calls Short Message Service	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number	2m		
Relationship of teleservices and bearer services to the GSM system Courtesy of ETSI). Service Category	GSM Network Service Telephony Emergency calls Short Message Service Videotext access	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number Point-to-point (between two users) and cell	2m		
Relationship of teleservices and bearer services to the GSM system Courtesy of TSI). Service Category GSM Teleservices	GSM Network Service Telephony Emergency calls Short Message Service	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI). Service Category GSM Teleservices	GSM Network Service Telephony Emergency calls Short Message Service Videotext access Teletex, FAX, etc. Asynchronous data	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number Point-to-point (between two users) and cell broadcast types 300–9600 bps (transparent/nontransparent)	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI). Service Category GSM Teleservices	GSM Network Service Telephony Emergency calls Short Message Service Videotext access Teletex, FAX, etc. Asynchronous data Synchronous data	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number Point-to-point (between two users) and cell	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI). Service Category GSM Teleservices	GSM Network Service Telephony Emergency calls Short Message Service Videotext access Teletex, FAX, etc. Asynchronous data	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number Point-to-point (between two users) and cell broadcast types 300–9600 bps (transparent/nontransparent)	2m		
Relationship of teleservices and bearer services to the GSM system (Courtesy of ETSI). Service Category GSM Teleservices	GSM Network Service Telephony Emergency calls Short Message Service Videotext access Teletex, FAX, etc. Asynchronous data Synchronous data Synchronous data Synchronous data	Additional Details Full rate at 13 kbps voice "112" is GSM-wide emergency number Point-to-point (between two users) and cell broadcast types 300–9600 bps (transparent/nontransparent) 2400–9600 bps transparent	2m		

TABLE 5–1 Phase 1 GSM services (Courtesy of ETSI).

	Service protection of the	Additional	Details	A De La Martin				
GSM Teleservices	Half-rate speech coder Enhanced full rate	Optional in	mplementation					
Supplementary Services	Calling line identification Connected line identifica Call waiting Call hold Multiparty communicatio Closed user group Advice of charge Operator determined cal barring	tion the caller's Presentatio the called I Incoming c Put current Up to five c in one com Restriction	on or restriction o ID call during currer t call on hold to a ongoing calls car	of displaying at conversation answer another a be included es from	2	2m		
TABLE 5–2 Phase	2 GSM services (Courtesy c	of ETSI).		5 (-5-19-3484 1 8009 (16872)				
Draw and expla	in GSM TDMA frame	e with logical	channel.		[[4]	CO3	
	Tani (con - TI	OMA Frame		1.32521 - 1724				
TSO	TS1 TS2 TS	3 TS4	TS5	TS6 TS	7 2	2m		
sona está i	head and the strength of the		ellist. (c					
e da atenana da e	Physical Chanr (One Timeslot		nada na	a ar suite an				
	carrier frequency, a c	hannel consis			occurs			
			1	0 3 7 17 1				
•	frame of eight times meslots represents a p	· •	•	efer Mullet)	0	2m		
Each of these ti	0	hysical chann	nel			2m [10]	CO3	
Each of these ti Draw and desc	meslots represents a p	hysical chann TDMA fram	nel				CO3	
Each of these ti Draw and desc	meslots represents a pribe the structure of	ohysical chann TDMA fram s. perframes = 271	nel ne, multifran 15648 TDMA F	ne, superframe			CO3	
Each of these ti Draw and desc	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su	ohysical chann TDMA fram s. perframes = 271 3 seconds and 70	ne, multifran 15648 TDMA F 60 microsecor	ne, superframe	e and [CO3	
Each of these ti Draw and desc hyperframe. Spe	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 132 Either Fifty-One 26-I	TDMA fram TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 5-TQMA Frames	ne, multifran 15648 TDMA F 60 microsecon 2042 2043 2 (≈6.12 second es or Twenty-S	ne, superframe rames ids) 044 2045 2046	e and [[10]	CO3	
Each of these ti Draw and desc hyperframe. Spe	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 132 Either Fifty-One 26-I	TDMA fram TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 5 TDMA Frames frame Multifram	ne, multifran 15648 TDMA F 60 microsecon 2042 2043 2 (≈6.12 second es or Twenty-S	ne, superframe rames ids) 044 2045 2046	e and [[10]	CO3	
Each of these ti Draw and desc hyperframe. Spe	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 132 Either Fifty-One 26-I 51-Fr	TDMA fram TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 5 TDMA Frames frame Multifram	tiel 15648 TDMA F 15648 TDMA F 15648 TDMA F 12042 2043 2 $(\approx 6.12 \text{ second}$ es or Twenty-S 146 47 48	ne, superframe rames ids) 044 2045 2046 s) six 49 50 Multifr	e and [[10]	CO3	
Each of these ti Draw and desc hyperframe. Spo 0 1 2 3	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 132 Either Fifty-One 26-I 51-Fr	TDMA fram TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 5 TDMA Frames frame Multifram	tiel tie, multifran 15648 TDMA F 60 microsecon 2042 2043 2 (≈ 6.12 second es or Twenty-S 46 47 48 24 1 (51 fram	ne, superframe rames ids) 044 2045 2046 s) six 49 50 Multifr	e and [[10]	CO3	
Each of these ti Draw and desc hyperframe. Spo 0 1 2 3	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 132 Either Fifty-One 26-I 51-Fr 0 1 2 3 4 0 1 2	hysical chann TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 Control 2040 2041 Control 2040 2041	tiel tie, multifran 15648 TDMA F 60 microsecon 2042 2043 2 (≈ 6.12 second es or Twenty-S 46 47 48 24 1 (51 fram	ne, superframe rames ids) 044 2045 2046 is) ix 49 50 Multifr he) Multiframe =	ames	[10]	CO3	
Each of these ti Draw and desc hyperframe. Specific 0 1 2 3	meslots represents a pribe the structure of eacify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 1324 Either Fifty-One 26-F 51-Fr 0 1 2 3 4 0 1 2 ne) Multiframe = Frames (120 ms) 22 23 24 25	hysical chann TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 5 TDMA Frames ame Multiframe ame Multiframe	hel he, multifran 15648 TDMA F 60 microsecon 2042 2043 2 (≈ 6.12 second es or Twenty-S 46 47 48 24 1 (51 fram 51 TDMA 0 1 2 3	ne, superframe rames ids) 044 2045 2046 is) ix 49 50 25 Multifr he) Multiframe = Frames (235 ms) 47 48 4	ames	[10]	CO3	
Each of these ti Draw and desc hyperframe. Specific 0 1 2 3	meslots represents a p ribe the structure of ecify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 132 Either Fifty-One 26-1 51-Fr 0 1 2 3 4 0 1 2 he) Multiframe = Frames (120 ms) 22 23 24 25	hysical chann TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 STIDMA Frames rame Multiframe ame Multiframe Timeslots (120/	the ine, multifram 15648 TDMA F 60 microsecon 2042 2043 2 (≈ 6.12 second es or Twenty-S 46 47 48 24 1 (51 fram 51 TDMA 0 1 2 3 26 ≈ 4.615 ms	ne, superframe rames ids) 044 2045 2046 (s) ix 49 50 Multifr e) Multiframe = Frames (235 ms) 47 48 4	ames	[10]	CO3	
Each of these ti Draw and desc hyperframe. Spectrum 0 1 2 3 1 (26 fram) 0 1 2 3	meslots represents a pribe the structure of eacify their time length 1 Hyperframe = 2048 Su (3 hours 28 minutes 5 4 5 6 7 1 Superframe = 1324 Either Fifty-One 26-F 51-Fr 0 1 2 3 4 0 1 2 ne) Multiframe = Frames (120 ms) 22 23 24 25	hysical chann TDMA fram s. perframes = 271 3 seconds and 7 2040 2041 5 TDMA Frames rame Multiframe ame Multiframe Timeslots (120/7 TS3 TS4 TS	11 2042 2043 2 12 2042 2043 2 12 2042 2043 2 12 2042 2043 2 14 47 48 48 1 (51 fram 51 TDMA 0 1 2 3 26<=	ne, superframe rames ids) 044 2045 2046 is) ix 49 50 25 Multifr he) Multiframe = Frames (235 ms) 47 48 4	e and [2047] ames	[10]	CO3	

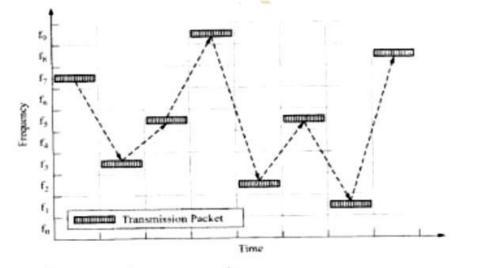
5. Briefly explain the spread spectrum modulation techniques: FHSS and DSSS.

There are two basic ways of implementing spread spectrum transmission:

- 1. Frequency Hopping Spread Spectrum (FHSS)
- 2. Direct Sequence Spread Spectrum (DSSS).

Frequency-hopping spread spectrum (FHSS) transmission is the <u>repeated</u> <u>switching of frequencies</u> during radio transmission to reduce interference and avoid interception. It is useful to counter eavesdropping, or to obstruct jamming of telecommunications. It can minimize the effects of unintentional interference. It consists of a system that changes the center frequency of transmission on a periodic basis with a pseudorandom sequence.

Here data are transmitted through number of different carrier frequencies hops. All the carrier frequencies hop independent from one another. For the system to work both the transmitter and receiver must have prior knowledge of the hopping sequence. In a frequency-hopping spread spectrum (FHSS) system, the transmitted signal is spread across multiple channels. Ex: The full bandwidth is divided into 8 channels, centered at f_0 through f_9 . The signal "hops" between them in the following sequence: f_7 , f_3 , f_5 , f_9 , f_2 , f_5 , f_1 , f_8 .



Frequency hopping spread spectrum example.

- Direct Sequence Spread Spectrum (DSSS) is a spread spectrum technique whereby the original data signal is <u>multiplied with a pseudo random noise</u> <u>spreading code</u>.
- This spreading code has a <u>higher chip rate</u> (this the bit rate of the code), which results in a <u>wideband time continuous scrambled signal</u>.
- DSSS significantly improves <u>protection against interfering</u> (or jamming) signals, especially narrowband and <u>makes the signal less noticeable</u>.
- It also <u>provides security of transmission</u> if the code is not known to the public.
- These reasons make DSSS very popular by the military. In fact, DSSS was first used in the 1940s by the military

4m

3m

3m

[10]

6. Explain working of BSS, DS and ESS network with neat diagrams.

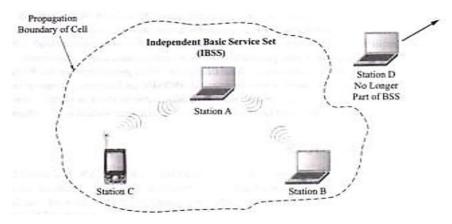


Figure 9-2 A typical independent basic service set.

The **Basic Service Set** (BSS) is the simplest and most fundamental structure of an IEEE 802.11x WLAN.

Architecture description:

There is no backbone infrastructure and the network consists of at <u>least two</u> <u>wireless stations</u>. BSS structure is referred to as a <u>peer-to-peer or ad hoc</u> <u>wireless network</u>. The propagation boundary will exist but its exact extent and shape are <u>subject to many variables</u>. Simulation software exists that can provide some <u>reasonable estimates of RSS</u> for typical multi-floor architectural layouts and various building materials. It is also possible to have <u>two or more</u> of these IBSSs in existence and operational <u>within the same general area</u> but <u>not in communication</u> with one another. An STA may be turned on or off or come into or go out of range of the BSS an <u>unlimited number of times</u>. The STA becomes a member of the BSS structure when it becomes associated with the BSS.

Distribution System (DS)

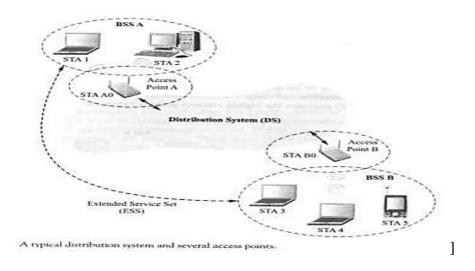
It provides an <u>extended wireless network</u> consisting of multiple BSSs.To provide flexibility to the WLAN architecture, IEEE 802.11 <u>logically separates</u> the wireless medium (WM) from the Distribution System Medium (DSM). The function of the DS is <u>enable mobile device support</u>. It provides <u>seamless</u> integration of multiple BSSs. This function is physically performed by a device known as an <u>access point (AP)</u>. The AP provides access in the DS by providing DS services and at the same time performing the STA function within the BSS.

Data transfers occur between stations within a BSS and the DS via an AP. All the APs are also stations and as such have addresses. However, the address used by an AP for data communications on the WM side and the one used on the DSM side are not necessarily one and the same. This DS structure gives rise to the use of <u>APs as bridges to extend the reach of a network</u>.

1m

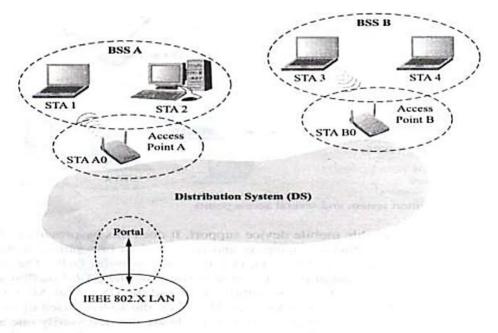
[10]





Extended Service Set (ESS)

The IEEE 802.11 standard provides for the use of <u>multiple BSSs and a DS to</u> <u>create a wireless network of arbitrary size</u> and complexity networks are known as <u>extended service set (ESS) networks</u>. ESS networks provide advantages, so that stations within an ESS network may communicate with one another and <u>mobile stations may move transparently</u> from one BSS to another as long as they are all part of the same ESS network. The most basic BSS consists of one AP and one STA.



A wireless LAN with a connection to an IEEE 802.x wired LAN.

Due to use of an ESS network the following situations may occur:

4m

- □ BSSs may overlap to provide continuous coverage areas or BSSs can be physically separate entities
- □ BSSs may be physically collocated for redundancy reasons, and one or more ESS networks may be physically located in the same area.
- □ The above situation can commonly occur when separate organizations set up their own WLANs in close proximity to one another.

7. Describe Bluetooth protocol stack with relevant figures.

Applications/Profiles

IEEE 802.15.1 Bluetooth/WPAN

Figure 10-7 The Bluetooth protocol stack (Courtesy of IEEE).

The Bluetooth standard call for a <u>set of communication protocols and a set of interoperable application</u> that are used to support the usages address in the specifications. The <u>link manager protocol (LMP)</u> and the <u>Logical Link Control And Adaptation(L2CAP) layer protocol</u> are Bluetooth specific whereas the protocols within the "Other" box are not. Some of these other protocols are the point-to-point protocol (PPP) and wireless application protocol (WAP).

Physical radio layer: It is for Tx and Rx data and voice.

Baseband layer: It enables RF link between Bluetooth devices.

Link manager: It is the protocol that handles link establishment b/w Bluetooth devices which include authentication and encryptions.

LLC and L2CAP: It is connection based communication protocol that implements multiplexing. No flow control. But provide reliable base band link. **Audio profile:** It responsible for managing connection for Tx /Rx data from audio devices.

Control: For control signal generations for various activities.

Other LLC: Link controller for optional device, fax, headsets like cordless phone etc.

8. Explain RAKE receiver with a neat block diagram.

A rake receiver is a radio receiver designed to <u>counter the effects of multipath</u> <u>fading</u>. It does this by using <u>several "sub-receivers" called fingers</u>, that is, <u>several correlators</u> each assigned to a different multipath component. <u>Each finger independently decodes a single multipath component</u>; at a later stage the contribution of all fingers are **combined** in order to make the most use of the different transmission characteristics of each transmission path. This could very well result in <u>higher signal-to-noise ratio</u> (or Eb/N0) in a multipath environment than in a "clean" environment.

Since each component contains the original information, if the magnitude and time-of-arrival (phase) of each component is computed at the receiver (through a process called channel estimation), then all the components can be added. The outputs of each correlator are weighted to provide better estimate of the transmitted signal than is provided by a single component. Demodulation and bit decisions are then based on the weighted outputs of the *M* correlators.

4m

[10]

5m

CO4

L4

CO5 | L3

[10]

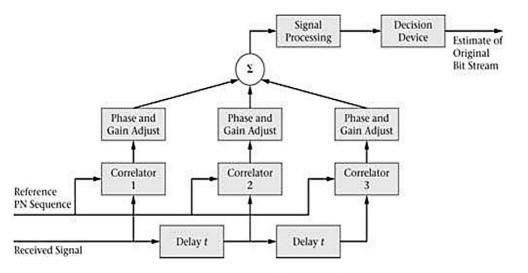


Figure 7.10: RAKE receiver block diagram

• <u>Each correlator detects a time-shifted version</u> of the original CDMA transmission, and each finger of the RAKE correlates to a portion of the signal, which is delayed by at least one chip in time from the other fingers.

- Assume M correlators are used in a CDMA receiver to <u>capture M strongest</u> <u>multipath components</u>.
- A weighting network is used to provide a <u>linear combination</u> of the correlator output for bit decision.
- Correlator 1 is synchronized to the strongest multipath m1.
- Multipath component m2 arrived t1 later than m1 but has low correlation with m1