



		In	ternal Asse	ssment Test 1	Ansv	wer Key– S	ep. 2020			
Sub:	System Softwar	re		Sub Code:	18MCA34	Branch:	MC	A		
Date:	07/09/2020	Duration:	90 min's	50	Sem	]	III			

## Q1 ) Describe SIC/XE Machine Architecture.

<ol><li>Memory</li></ol>	1)	1	L)	Λ	Λe	n	10	ry	1
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Memory consists of 8-bit bytes.
3 consecutive bytes form a word (24 bits).
All the address in SIC/XE are byte addresses.
Words are addressed by the location of their lowest numbered byte.
Maximum memory available on a SIC/XE system is 1 megabyte (220 bytes).
This increase leads to a change in instruction formats and addressing modes

#### 2) Registers

Five registers of SIC machine remains same in SIC/XE. The additional registers provided by SIC/XE are as follows.

Mnemonic	Number	Use
В	3	Base register; used for addressing.
S	4	General working register – no special use.
Т	5	General working register – no special use.
F	6	Floating-point accumulator (48 bits).

## 3) Instruction Formats

- SIC/XE has larger memory hence instruction format of standard SIC version is no longer suitable.
- SIC/XE provide two possible options; using relative addressing (Format 3) and extend the address field to 20 bit (Format 4).
- In addition SIC/XE provides some instructions that do not reference memory at all. (Format 1 and Format 2).
- The new set of instruction format is as follows. Flag bit e is used to distinguish between format 3 and format 4. (e=0 means format 3, e=1 means format 4)
- 1. Format 1 (1 byte)

8

ор

Example RSUB (return to subroutine)

### opcode

0100	1100
4	С

# 2. Format 2 (2 bytes)

8	4	4
ор	r1	r2

Example COMPR A, S (Compare the contents of register A & S)

Opc	ode	Α	S
1010	0000	0000	0100
Α	0	0	4

# 3. Format 3 (3 bytes)

· ·	_	_	_	_	1	_	12
ор	n	i	Х	b	р	е	disp

# Example LDA #3(Load 3 to Accumlator A)

000000	0	1	0	0	0	0	0000 0000 0011
0	n	i	Х	b	g	e	0 0 3

# 4. Format 4 (4 bytes)

6	1	1	1	1	1	1	20
ор	n	i	х	b	р	е	address

# Example JSUB RDREC(Jump to the address, 1036)

0100 10	1	1	0	0	0	1		0000 0001 0000 0011 0110
			n	i	Х	b	р	e

## 4) Addressing Modes

Two new relative addressing modes are available for use with instructions assembled using Format 3

Mode	Indication	arget address
		alculation
Base Relative	b=1, p=0	TA = (B) + disp (
		0≤ disp ≤ 4095)
Program-counter	b=0, p=1	TA = (PC)+disp (-
relative		2048 ≤ disp ≤
		2047)

b represents for base relative addressing where as p represents program counter relative addressing. If both the bits b and p are 0 then target address is taken form the address field of the instruction (i.e displacement)

SIC/XE also support addressing modes that are assembled using Format 4.

Mode	Indication	Target address calculation
Direct	b=0, p=0, x=0	TA = disp
Indexed	x=1	TA = (x) + disp
Immediate	i=1, n=0	TA = operand value
Indirect	i=0, n=1	TA = address of
		operand value
simple	i=1, n=1 i=0, n=0	TA = location of the
		operand value

### 6) Instruction Set

- SIC/XE provides all of the instructions that are available on the standard version.
- In addition we have, Instructions to load and store the new registers LDB, STB, etc,
- Floating-point arithmetic operations, ADDF, SUBF, MULF, DIVF,
- Register move instruction: RMO,
- Register-to-register arithmetic operations, ADDR, SUBR, MULR, DIVR and,
- Supervisor call instruction : SVC

### 7) Input and Output

- There are I/O channels that can be used to perform input and output while the CPU is executing other instructions.
- Allows overlap of computing and I/O, resulting in more efficient system operation.
- The instructions SIO, TIO, and HIO are used to start, test and halt the operation of I/O channels.

Q2 a) Define system software? List and explain any four assembler directives with examples.

System Software consists of a variety of programs that support the operation of a computer. It makes possible for the user to focus on an application or other problem to be solved, without needing to know the details of how the machine works internally.

They are usually related to the architecture of the machine on which they are to run.

Example: Assembler, Compiler, text editor, loader and linkers etc.

#### **Assembler Directives**

In addition to the mnemonic machine instructions assembler uses following assembler directives. These statements are not translated into machine instructions. Instead they provide instructions to assembler itself.

1) START

START specify the name and starting address of the program.

Example: START 1000

2) END

Indicate the end of the source program and (optionally) specify the first executable instruction in the program.

Example: END FIRST

3) BYTE

Generate character or hexadecimal constant, occupying as many bytes as
needed to represent the constant.
Example: BYTE X'F1'
4) WORD
Generate one-word integer constant
Example: THREE WORD 3
5) RESB
Reserve the indicate number of bytes for a data area.
Example: BUFFER RESB 4096 6) RESW
Reserve the indicate number of words for a data area.
Example: LENGTH RESW 1
Q2 b) Write general description of pass 1 and pass 2 for two pass assembler
Pass 1 (define symbols)
☐ Assign addresses to all statements in the program
☐ Save the addresses assigned to all labels for use in Pass 2
☐ Perform some processing of assembler directives, (including those for address assignment, such as
BYTE and RESW
Pass 2 (assemble instructions and generate object program)
☐ Assemble instructions (translate opcodes and look up addresses)
☐ Generate data values defined by BYTE, WORD etc.
$\square$ Perform processing of assembler directives not done during Pass 1
$\square$ Write the object program and the assembly listing

Q3) Write pass 1 Algorithm for pass two assembler

```
Assembler Pass 1:
begin
  read first input line
     if OPCODE ='START' then
           save #[OPERAND] as starting address
           initialize LOCCTR to starting address
           write line to intermediate file
           read next input line
           end {if START}
     else
           initialize LOCCTR to 0
     while OCODE != 'END' do
          begin
                if this is not a comment line then
                      if there is a symbol in the LABEL field then
                           search SYMTAB for LABEL
                           if found then
                            set error flag (duplicate symbol)
                             insert (LABEL,LOCCTR) into SYMTAB
                       end {if symbol}
                      search OPTAB for OPCODE
                      if found then
                       add 3 (instruction length) to LOCCTR
                      else if OPCODE='WORD' then
                           add 3 to LOCCTR
                      else if OPCODE = 'RESW' then
                           add 3 * #[OPERAND] to LOCCTR
                      else if OPCODE = 'RESB' then
                           add #[OPERAND] to LOCCTR
                      else if OPCODE = 'BYTE' then
                       begin
                           find length of constant in bytes
                           add length to LOCCTR
                       end {if BYTE}
                      else
                        set error flag (invalid operation code)
                end {if not a comment}
           write line to intermediate file
           read next input line
      end {while not END}
 write last line to intermediate file
save (LOCCTR - starting address) as program length
end {Pass 1}
```

Q4) Write pass 2 Algorithm for pass two assembler

```
Assembler Pass2:
begin
 read first input line (from intermediate file)
 if OPCODE ='START' then
  begin
      write listing line
      read next input line
  end (if START)
 write Header record to object program
 initialize first Text record
 while OPCODE != 'END' do
   begin
       if this is not a comment line then
      begin
             search OPTAB for OPCODE
             if found then
                  begin
                   if there is a symbol in OPERAND field then
                        begin
                          search SYMTAB for OPERAND
                         if found then
                              store symbol value as operand address
                         else
                              begin
                                store 0 as operand address
                                set error flag (undefined symbol)
                            end
                        end (if symbol)
                    else
                        store 0 as operand address
                   assemble the object code instruction
        end (if opcode found)
            else if OPCODE ='BYTE' or 'WORD' then
                  convert constant to object code
       if object code will not fit into the current Text record then
              begin
                  write Text record to object program
                  initialize new Text record
              end
       add object code to Text record
            end (if not comment)
    write listing line
    read next input line
  end(while not END)
 write last Text record to object program
 write End record to object program
 write last listing line
end(Pass 2)
```

Q5 a) Write a program for SIC and SIC/XE machine to perform ALPHA = 11+BETA \* 4

LDA BETA

	MUL	FOUR
	ADD	ELEVEN
	STA	ALPHA
ELEVEN	WORD	11
FOUR	WORD	4
BETA	RESW	1
ALPHA	RESW	1
	LDA	#4
	MUL	BETA
	ADD	#11
	STA	ALPHA
BETA	RESW	1
ALPHA	RESW	1

Q6 a) Give the target address generated for following machine instruction.

if (B)=006000 (PC)=003000 (x)=000090 i)75101000 ii)032026 iii)03C300 iv)0310C303

i)75101000

000110   1   0   0   0   1   0   0000 0000 0000	00	0110	1	0	0	0	1	0	0000 0000 0000
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TA= operand value

TA = 01000

ii) 032026

000000	1	1	0	0	1	0	0000 0010 0110

TA = disp + (PC)

TA = 026 + 003000

TA=3026

iii) 03C300

000000	1	1	1	1	0	0	0011 0000 0000
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$$TA = disp + (x) + (b)$$
  
 $TA = 6390$ 

#### iv) 0310C303

000000   1   1   0   0   0   1   0000 1100 0011 0000 0011
---

TA = C303

Q7 a ) List and describe data structures used by two-pass assembler

### 1) OPTAB:

② It is used to lookup mnemonic operation codes and translates them to their machine language equivalents.

In more complex assemblers the table also contains information about instruction format and length

② In pass 1 the OPTAB is used to look up and validate the operation code in the source program.

② In pass 2, it is used to translate the operation codes to machine language.

In simple SIC machine this process can be performed in either in pass1 or in pass2.

But for machine like SIC/XE that has instructions of different lengths, we must search OPTAB in the first pass to find the instruction length for incrementing LOCCTR.

In pass 2 we take the information from OPTAB to tell us which instruction format to use in assembling the instruction, and any peculiarities of the object code instruction.

② OPTAB is usually organized as a hash table, with mnemonic operation code as the key.

The hash table organization is particularly appropriate, since it provides fast retrieval with a minimum of searching.

② Most of the cases the OPTAB is a static table- that is, entries are not normally added to or deleted from it. In such cases it is possible to design a special hashing function or other data structure to give optimum performance for the particular set of keys being stored.

#### 2) SYMTAB:

It is table includes the name and value for each label in the source program, together with flags to indicate the error conditions (e.g., if a symbol is defined in two different places).

② During Pass 1: labels are entered into the symbol table along with their assigned address value as they are encountered. All the symbols address value should get resolved at the pass 1.

② During Pass 2: Symbols used as operands are looked up the symbol table to obtain the address value to be inserted in the assembled instructions.

② SYMTAB is usually organized as a hash table for efficiency of insertion and retrieval. Since entries are rarely deleted, efficiency of deletion is the important criteria for optimization.

#### 3) LOCCTR:

② Apart from the SYMTAB and OPTAB, this is another important variable which helps in the assignment of the addresses.

② LOCCTR is initialized to the beginning address mentioned in the START statement of the program.

② After each statement is processed, the length of the assembled instruction is added to the LOCCTR to make it point to the next instruction.

② Whenever a label is encountered in an instruction the LOCCTR value gives the address to be associated with that label.

Q10 a) Write object program format. (Header, Text, End Record) for the following code

	QUIZ	START	0	
0000	FIRST	LDA	FIVE	000015
0003		STA	ALPHA	0C001B
0006		LDA	TWO	000018
0009		STA	BETA	0C001E
000C		LDA	ALPHA	00001B
000F		ADD	BETA	18001E
0012		STA	RESULT	0C0021
0015	FIVE	WORD	5	
0018	TWO	WORD	2	
001B	ALPHA	RESW	1	
001E	BETA	RESW	1	
0021	RESULT	RESW	1	
		END	FIRST	

HQUIZ 000000000023 T00000120000150C001B0000180C001E0C0021 E000000