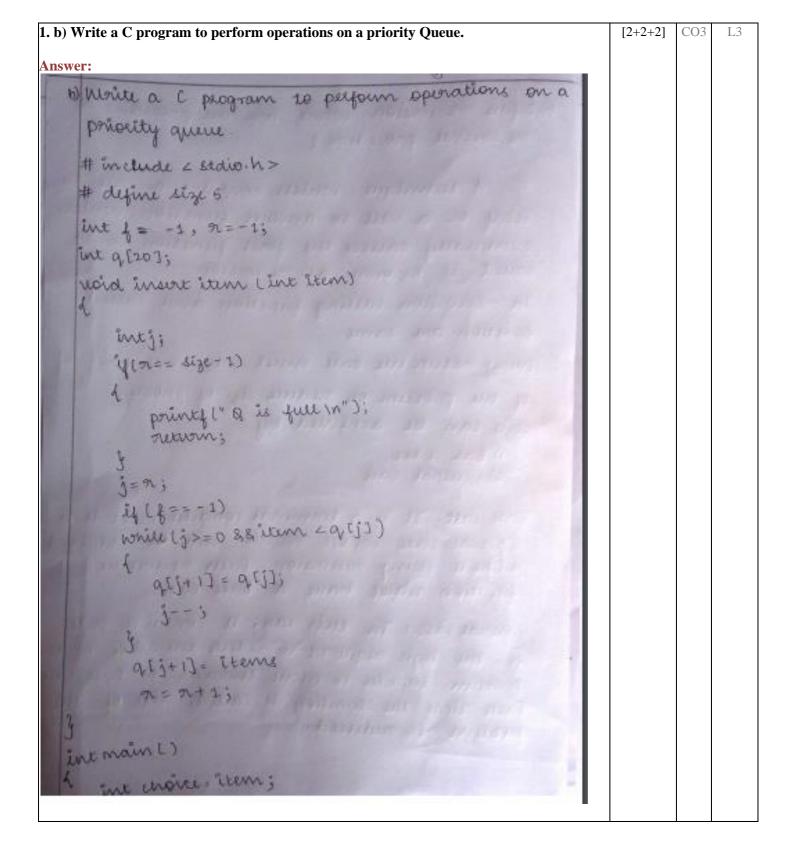


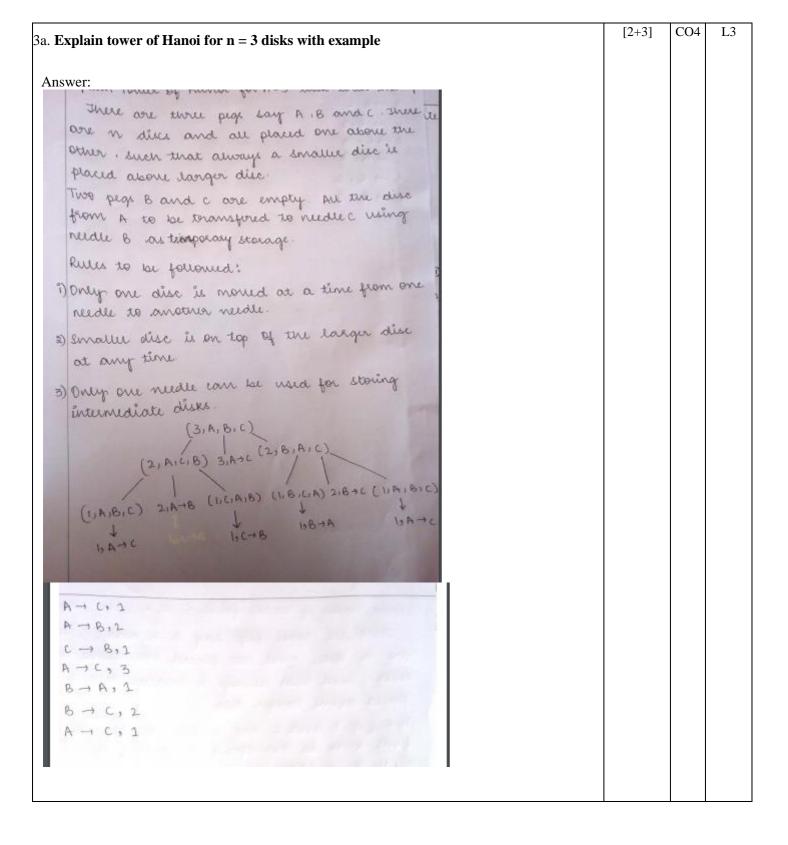
Sub: Data Structures and Applications  Date: 15/10/2019 Duration: 90 mins Max Marks: 50 Sem/ Sec: 3rd/A,B,C  Answer any FIVE FULL Questions  MARKS CO RB T  1. a) Define recursion. What are the properties of recursive procedure?  A turnique which contains a call to itself as a call to another function, which itself as a call to another function, which itself as a call to another function to be launced in known as recursion.  In recursion, calling function and called function are same Europe recursive call must salve one past of the problem of recursive stays of problem.  This can be achieved by:  () Base case 3) General case 4 Department call the absolute of the problem of	Inter	nal Asses	sment Test 2	2 – October 2	2019 Scheme an	d Sol	ution		ACCRED	ITED WITH A+ GR	ADE BY NAAC		
Answer any FIVE FULL Questions  Answer any FIVE FULL Questions  A technique, which contains a late to itsure as a late to itsure as a late to another function to be lated in known as necessary.  In necession, calling function and lated function are same function are same for the problem of reduced by:  () Base case 1) General case  Base case It is a terminal condition. It is a	Internal Assessment Test 2 – October 2019 Scheme and So Sub: Data Structures and Applications						Sub	18CS32	Branch:	ISE	ISE		
A technique, which contains a call to itself or a call to another function, which contains a call to itself or a call to another function, which eventually causes the first function to be called it known as newstron.  In necession, calling function and called function are same.  Every necessive call must solve one part of of the problem or reduce size of problem.  This can be achieved by:  () Base case  1) General case  Base case. It is a terminal condition. It is a	Date: 15	/10/2019	Duration:	90 mins	Max Marks:	50		3 <sup>rd</sup> /A,B,C			OBE		
A technique, which contains a call to itself or a call to another function, which exercises the first function to be called in known as remotion.  In recursion, calling function and called function are same.  Every recursive call must solve one part of the problem of the problem of reduce size of problem.  This can be achieved by:  1) Base case  2) General case.  Base case. It is a terminal condition. It is a		Answ	er any FIVE F	ULL Questions					MARK	S	CO	RB T	
itself or a rate to another function, which eventually caused the first function to be latted in known as remotion.  In recursion, ratting function and called function are same.  Every recursion call must salve one past of of the problem or reduce size of problem.  This can be achieved by:  () Base case  2) General case  Base case: It is a terminal condition. It is a	1. a) Define	recursion	. What are	the properti	ies of recursive	e proc	edure?		[2+2		CO3	L1	
special case whose solution can be obtained nuthout using neurosion Each recursive function must have a base case.  general case: In this case, it contains code for the logic required to reduce the sixe of the problem, impade to more towards the Base case.	itself erent caucal In one function of the Inis & 1) B 2) G Base & special neither function of the Inis & I	or wall	cause the same ase with a case in the same	terminate this c	notice of first of previous functions use size by:  inal constitution is the base of the base of the preduce of	gun on a use of car	ion I be o	t is a lotained ne					

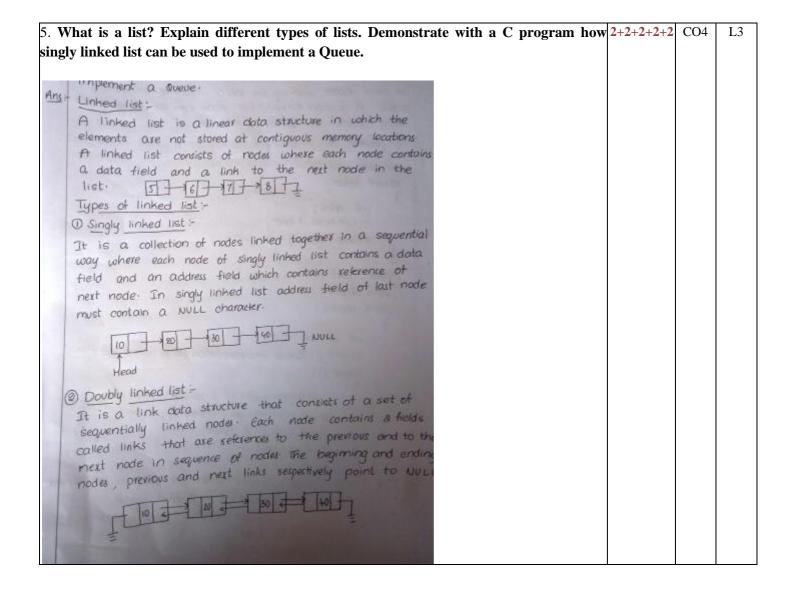


	(3)		
400	(133)		
4	print (1: insert, 2: delete, 3: display 4: exit (n);		
	Stornf (" %d", & choice);		
	Switch (choice)		
	case 1: pointly ("Enter the "items to be insisted in");		
	Scanf (" 1.d", & stem):		
	insert_item (item);		
	loneak;		
	case 2: delete front ()		
	break;		
	Case 3: olisplay()		
	break;		
	default: printf l'enualid choice");		
	getien 0;		
	*		
7			
	The second section of the second seco		
4			
is a Owens/I :	recor Output 2 List different terms of Output Write C implementation for [2+2-	+3+3] CO1	ł
_	near Queue? List different types of Queue. Write C implementation for $Q(0)$ operation.		

```
A Quene is a special data structure, where
elements are inserted from one end and
elements are deleted from the other end.
The end from the elements are inscrited is call
-ed greage and and where the elements agre
delited he could from end. Quere is also called
as first in first out data structure.
Types of Guene:
1. Linear Queue
2. likellar Quere.
3 Double Ended Quine
4 Priority Quene
program for insertal)
 word insert [ int item)
      4 (n = = lige - 1)
         printy (" Quere ourflow");
         rement;
     4(===-1)
       91=91+23
       gen] = Item;
Program for delete &():
  word delete ()
     4(f==-1 && 91==-1)
        printf ("Queue underflow (n");
         mendure;
        printy ["The element deleted "ie 1. d', q[f]);
         f= f+ 13
```

f = -13 90 = -13





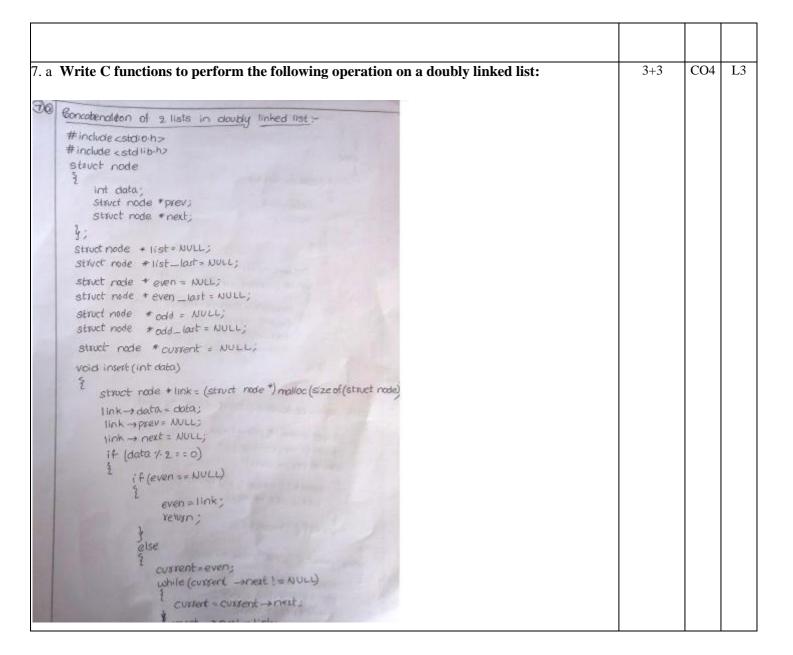
```
3 Circular linked list:-
It is a linked list in which the last node points to head
 of front node making the data structure to look like a
 Circle A circular linked list node can be implemented using singly or doubly linked list
 C program to implement Queve using singly linked list:
 #Include <stdio-h>
struct Node
 struct Node
       int obta;
       Struct Node * next;
       * Front = NULL, * recor = NULL;
 Void insert (int);
Void delete();
    Void main()
         int choice , value;
         printf (" queue implementation using linked list:\n");
         while (1)
              printf (" ir insert in a deletein 3 exitin");
              printf (" enter your choice:");
              scanf (" v.d", & choice);
              Switch (choice)
                  case 1: printf ("enter the value to be insented.")
                          scent (" 'r.d", & value).
                          insest (value);
                  break;
case a: delete();
break [];
                   case 3: exie(o);
default : printf (* invalid choice \n ");
```

```
void insest (int value)
    struct Node *newNode;
    newlode = (struct Node *) malloc (size of (struct Node));
    newNode → data = value;
    newNode → next = NULL;
    if (front = > NULL)
         front = rear = newNode;
     else
         reas - next = newNode;
          Year = newNode;
void delete()
     if (front == NULL)
          printf ("Queve is empty in");
     else
        struct wode + temp = front;
        front = front - next;
         printf (" deleted element : /d\n", temp -data);
         free (temp);
```

6а.	Implement the concept of multiple stack using arrays in C. Explain with suitable	[4+2+2+2]	CO4	L3	Ī
exa	ample pictorially.				
	capan was a same and a				
Ans	multiple stack - we can not be				
	When a stack is created using single array, we can not be able to store large amount of data, thus this problem is				
	sufficient array. This technique is called stack 8				
	3tock A -> 3 4 5 6 7 8 9				
	1234560000				
	Top Top				
	C - Program :-				
	#inclinde estato.h>				
	#include < mailochi				
	# define MAX 10 int stock[MAX], topA=-1, topB=MAX;				
	void push A (Int va)				
	if (top A = x top8-1)				
	if (top A = *top8-1)  printf ("\n overflow");				
19					
в	else				
10	$ \frac{1}{\text{stack}[top A]} = \text{val}; $				
	Stack Lot-3				
я					
	int pop +()				
	t int val;				
	if (top A = = -1)				

```
printf ("In underflow");
      Val = -999;
      val = stock [top A];
       topA -- ;
   return val;
void push B (int val)
   if (top 8 - 1 = = top A)
      printf ("In overflow");
    else
       top8 -= 1;
       stack[tops]=val;
 int pops ()
     int val;
     if (topB = = MAX)
        print+ (" in underflow")
         Val=- 999;
         val = stack(tope);
         topB++;
```

```
int main()
   int option, val;
   while (1)
      printf (" In 1. Push an element into stack A");
      printf ("In 2. Push an element into stock 8");
       printf ("In 3 Pop an element from stack A");
      printf (" In 4 Pop an element from stack B");
       printf ("In 5- Exit");
       printf (" in enter your choice");
       scanf (" 1/1 d", Loption);
       switch (option)
          case 1: printf ("In Enter the value to push on stack A:");
                   scanf (" /d", Rval);
                   push A (val);
                    break;
          Case a: printf (" In Enter the value to push on stack 8)
                    scanf ("1/d", (val);
                    push B(val)
                    break;
           case 3: FA (val ! = - 999)
                     prints ("In The value popped from start A
                               = 1/ d" , val),
                      break)
          case 4: 1f(val! = -999)
                      printf("In The value papped from stack B
                                = " d", val);
                      break;
      setusno;
```



```
link-> prev = current;
 else
     if (odd == NULL)
           odd = link;
           seturn;
       else
           current = odd;
          while (current - next! = NULL)
            i custert = custert > next;
g custert > next = link;
               odd-last=link;
               link → prev = current;
Void print-backward (struct node + head)
     struct node + ptr - head,
     printf("In [las](9");
     while (Pts! = NULL)
         printf ("1.d (=)", pts -> data);
         ptr = ptr - pser;
      printf ("Enead(\n");
 void printlist (struct node + head)
      struct node * ptr = head;
      printf ("In [head] >"); while (Ptx! = NULL)
```

```
Printf ("/d (=)", ptr -> data);
ptx = ptr -> next;
           grintf ("[last] \n");
yold combine()
             struct node * link;
             list = even;
link = list;
while (link = next! = wull)
                            link = link - next;
                  ink +next = odd;
odd + prev = link;
                   link _last = link;
  int main()
             for (1=1:31<=10;1++)

insert(1);

printf("Even:");

printf("Even(R)");

print = bockward (even_last);

printf("Odd ");

printf("Odd ");

printf("Odd (R).");

print=bockward (Odd_Art);

eombine();

printf("combined list : \r");

printf("Combined list (R). \r");
```

```
Reverse a doubly linked list:
     #include < stdio h >
     #include astalibih>
     struct Node
         int data;
         struct Node * next;
         struct Node + prev;
      Void reverse (struct Node ** head_ref)
         struct Node * temp = NULL;
         struct Node # custent = *head-ref;
         while (current 1 = NULL)
              temp = current - prevs
              current - prev = current - next;
              current - next = temp;
              current = current -> prev;
          (f (semp ( = NULL)
              * head_xef = temp -> prev;
      Yold push (struct Node ** head-let, int new_data)
          struct Node * new_node = (struct Node *) malloc (sizeof
           new_node - data = new_data;
           new_node -> prev = NULL;
           new_node - next = ( + head - ref),
            if (( * head_ ref) ! = NULL)
               ( head_ref) - prev = new_node;
               (thead_sef) = new_node,
7. b) Write a C function to evaluate a given polynomial (For eg: 5x^3 + 2x^2 + 9) using single
                                                                                                      [2+2]
                                                                                                               CO3
                                                                                                                      L3
linked list for a given value of X
void evaluate_poly()
    int x, sum = 0;
    printf("\nEnter the value of x: ");
    scanf("%d", &x);
    POLY temp;
    temp = first;
    while(temp != NULL)
         sum = sum + (temp->coeff*pow(x, temp->px));
         temp = temp->link;
    printf("\nThe sum is: %d", sum);
```

```
8.a What is the disadvantage of singly linked list? How is it overcome? Explain front
                                                                                                   2.5+1.5+3+ CO4
                                                                                                                     L3
insertion and end deletion on a doubly linked list.
     Disadvantages of singly linked list:
Anst
   More memory is required to store elements in linked list
    =) If we have to go to a particular element then we
       have to go through all those elements that come before
       that element
    >) We can't traverse it from last but only from beginning
    It can be overcome by using doubly linked list. As we
    have a links tright and left, we can go through a
     particular element without passing through all elements
     Code for front insestion :-
     struct Node
          int data;
          struct Node * rest;
          struct Node * prev;
     void push (stavet Node ** head - ref)
          struct Lode * new_node = (struct Node *) malloc(size of (state
           new=node + data = new_data;
           rew_node = next = (*head_set);
            new-node -> prev = NULL;
            if (( + head_ set) ! = NULL)
                (+ head_ref) -> prev = new_node
             (*nead_set) = new_node;
      Code for end deletion :-
      struct Node
         int data;
         struct Node * next;
         struct Node + psev;
      void pop ()
          struct node * ptr;
          if (head = = NULL)
              printf ("In underflow "):
           else if (head-next == NULY)
              head = NULL;
              free (head),
               print ("In Node deleted");
              pty=head;
              if (ptx -> next!=NULL)
                  pty = ptx -next;
               pty -> prev -> next = NULL)
               free (ptr);
               printf (" In Node deleted");
```