USN					



Internal Assessment Test 2 – December 2021

Sub:	Computer Netv	works and Sec	curity		•	Sub Code:	18CS52	Bra	nch:	ISE	•	
Date:	16/12/2021	Duration:	90 min's	Max Marks:	50	Sem/Sec:	V A, B & C				OB	BE
			MA	RKS	CO	RBT						
1	Build a Relial	ble Data Tra	ansfer Proto	ocol to satisfy	follo	wing condi	tions		[]	10]	CO2	L3
	Stop a	nd wait pro	tocol.									
	 Packet 	 Packet contains sequence number, data and checksum. 										
	Need 1	 Need to use both ACK and NACK. 										
	Need to handle duplicates.											
2	Design a FSM	I for the pro	otocol whic	h retransmits	ever	v packet fro	m the last AC	CK	[:	10]	CO2	L2
	after timeout	-			•	. 1						
3a)	Draw the Con	nection Ori	iented Proto	ocol segment s	truc	ture and exp	plain.		[[7]	CO2	L2
							•			21	002	T 1
3b)	Explain about	-	lved in com	imunication vi	a Vı	rtual Circui	t Network		L	[3]	CO3	L1
	(VCN) in nety											
	How does cor			and connection	on te	rmination o	ccur in		[[7]	CO2	L1
	connection or	iented proto	ocol?									
4b)	Differentiate l	between AC	CK and SYI	N in connection	n or	iented comi	munication w	ith	[[3]	CO2	L2
	example.											

P.T.O.





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	_	<u>Ar</u>	swer any FI	VE FULL Ques	tions				MA	RKS	CO	RBT
1	Build a Relial	ole Data Tra	ansfer Proto	ocol to satisfy	follo	wing condi	itions		[3	10]	CO2	L3
	Stop and wait protocol.											
	 Packet contains sequence number, data and checksum. 											
	Need to use both ACK and NACK.											
	Need to handle duplicates.											
2	Design a FSM for the protocol which retransmits every packet from the last ACK									10]	CO2	L2
	after timeout event.											
3a)	Draw the Con	nection Or	iented Proto	ocol segment s	struc	ture and ex	plain.		[[7]	CO2	L2
							•		г	21	CO2	L.1
3b)	Explain about		lved in com	imunication vi	a Vi	rtual Circui	t Network		L	[3]	CO3	LI
	(VCN) in nety										000	Y 4
4a)	How does cor			and connection	on te	rmination o	occur in		L	[7]	CO2	L1
	connection or											
4b)	Differentiate l	between A(CK and SY	N in connection	n or	iented com	munication w	ith	[[3]	CO2	L2
	example.											

5	Design co	ommunication flow for reliable data transfer in following situations	[10]	CO2	L3
	(i)	Retransmission due to lost ACK.			
	(ii)	Sending multiple packets back-to-back.			
	(iii)	Avoiding retransmission by cumulative ACK.			
	(iv)	Fast retransmit.			
	(v)	Slow start.			
6 (a)	Draw an	[7]	CO3	L1	
6(b)	List and c	[3]	CO3	L1	

Faculty Signature CCI Signature HOD Signature

5	Design co (vi) (vii) (viii)	mmunication flow for reliable data transfer in following situations Retransmission due to lost ACK Sending multiple packets back-to-back Avoiding retransmission by cumulative ACK	[10]	CO2	L3
	(ix) (x)	Fast retransmit Slow start			
6 (a)	Draw an	d explain IPv4 datagram format	[7]	CO3	L1
6(b)	List and d	escribe 3 types of switching fabric	[3]	CO3	L1

Faculty Signature CCI Signature HOD Signature



Scheme of Evaluation Internal Assessment Test 2 – December 2021

Sub:		Computer N	Code:	18CS52					
Date:	16/12/2021	Duration:	90mins	Max Marks:	50	Sem:	V	Branch:	ISE

Note: <u>Answer Any five full questions.</u>

_	stion #	Description	Marks Dis	stribution	Max Marks
		Build a Reliable Data Transfer Protocol to satisfy following conditions • Stop and wait protocol. • Packet contains sequence number, data and checksum. • Need to use both ACK and NACK. Need to handle duplicates.			
1		Diagram Explanation		10M	10M
			6M+4M		
2		Design a FSM for the protocol which retransmits every packet from the last ACK after timeout event.		10M	10M
		FSM Diagram Protocol Explanation	6M 4M		
	a)	Draw the Connection Oriented Protocol segment structure and explain.			
3		Packet Diagram Explain each field	4M 3M	7M	10M

	b)	Explain about steps involved in communication via Virtual Circuit Network (VCN) in network layer.			
3		Listing Steps Explanation	1M 2M	3M	
4	a)	Explanation	2.5M 1M 2.5M	7M	10M
4	b)	Differentiate between ACK and SYN in connection oriented communication with example. Explanation with Example	1.5M+ 1.5M	3M	
5	a)	Design communication flow for reliable data transfer in following situations (i) Retransmission due to lost ACK. (ii) Sending multiple packets back-to-back. (iii) Avoiding retransmission by cumulative ACK. (iv) Fast retransmit.		10M	
		Liveral are a fine	4M 6M		10M

	a)	Draw and explain IPv4 datagram format.			
		Explanation Packet format	3M		
6		r acket format	4M	7M	
					10M
	b)	List and describe 3 types of switching fabric.			TOW
6		Listing 3 types	1M 2M	3M	
		Explanation	21V1		



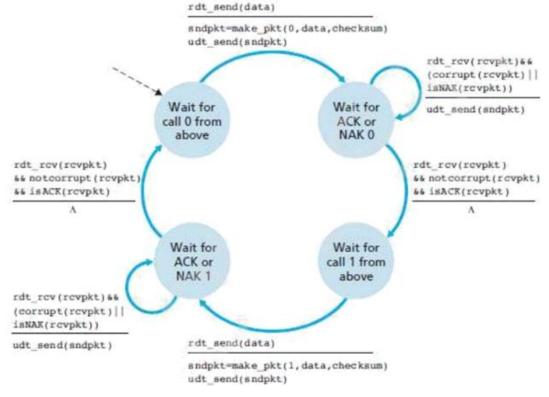
Scheme Of Evaluation Internal Assessment Test 2 – December 2021

Sub:		Computer N	Code:	18CS52					
Date:	16/12/2021	Duration:	90mins	Max Marks:	50	Sem:	V	Branch:	ISE

Note: Answer Any full five questions

Q. 1 Build a Reliable Data Transfer Protocol to satisfy following conditions

- Stop and wait protocol.
- Packet contains sequence number, data and checksum.
- Need to use both ACK and NACK.
- Need to handle duplicates.



Problem with rdt2.0: If an ACK or NAK is Corrupted, the sender cannot know whether the data is received or not.

Solution: The sender resends the current data packet when it receives garbled ACK or NAK packet.

Problem: This approach introduces duplicate packets into the channel.

Solution: Add sequence-number field to the data packet.

The receiver has to only check the sequence-number to determine whether the

Received packet is a retransmission or not. For a stop-and-wait protocol, a 1-bit sequence-number

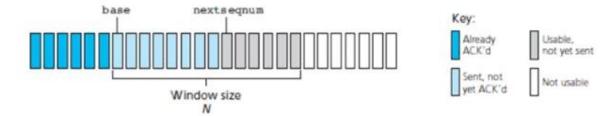
will be sufficient.

- A 1-bit sequence-number allows the receiver to know whether the sender is sending previously transmitted packet {0} or new packet (1).
- We call this protocol as rdt2.1

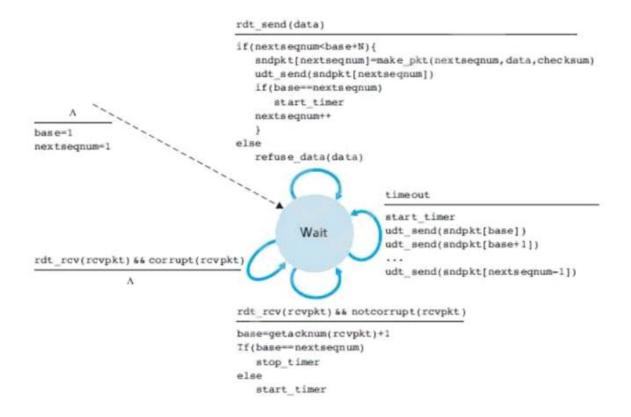
Q.2 Design a FSM for the protocol which retransmits every packet from the last ACK after timeout event.

It is GoBack-N protocol.

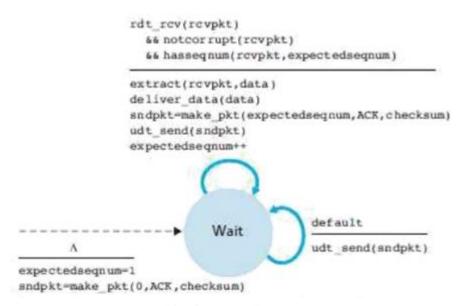
- The sender is allowed to transmit multiple packets without waiting for an acknowledgement
- \bullet But, the sender is constrained to have at most N unacknowledged packets in the Where N -window-size which refers maximum no. of unacknowledged packets in the pipeline
- GBN protocol is called a sliding-window protocol.



Sender's view of sequence-numbers in Go-Back-N

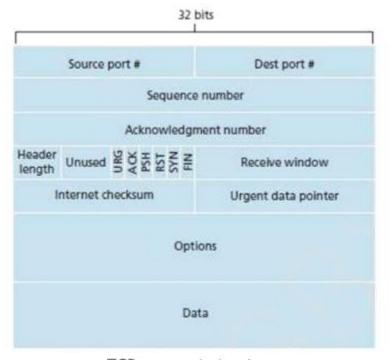


Extended FSM description of GBN sender



Extended FSM description of GBN receiver

Q.3 a) Draw the Connection Oriented Protocol segment structure and explain..



TCP segment structure

- The segment consists of header-fields and a data-field.
- The data-field contains a chunk of-data.
- When TCP sends a large file, it breaks the file into Chunks of size MSS.
- The fields of TCP segment are as follows:
- 1) **Source and Destination Port Numbers:** These fields are used for multiplexing/demu1tiplexing data from/to upper-layer applications.
- 2) **Sequence Number & Acknowledgment Number:** These fields are used by sender & receiver in implementing a reliable data-transfer-service.
- 3) **Header Length:** This field specifies the length of the TCP header.
- 4) Flag: This field contains 6 bits.
- i} ACK This bit indicates that value of acknowledgment field is valid.
- ii) RST, SYN & FIN: These bits are used for connection setup and teardown.
- iii) **PSH:** This bit Indicates the sender has invoked the push operation.
- iv) URG: This bit indicates the segment contains urgent-data.
- 5) Receive Window: This field defines receiver's window size This field is used for low control.
- 6) **Checksum:** This field is used for error-detection.
- 7) urgent Data Pointer: This field indicates the location of the last byte of the urgent data.
- 8) **Options:** This field is used when a sender & receiver negotiate the OSS for use in high-speed networks.

Q.3 b) Explain about steps involved in communication via Virtual Circuit Network (VCN) in network layer.

- A VC consists of
- 1) A path between the source and destination.
- 2) VC number: This is one number for each link along the path.
- 3) Entries in the forwarding-table in each router.
- A packet belonging to a virtual-circuit will carry a VC number in its header.
- At intervening router, the VC number of traversing packet is replaced with a new VC number.
- The new VC number is obtained from the forwarding-table.

Virtual—circuit setup

VC setup is simplified by permitting a different VC number at each link along the path.

• Disadvantage:

The routers must maintain connection srate information for the ongoing connections.

• Three phases in a virtual-circuit

1) VC Setup

- During the setup phase, the sending transport-layer
- contacts the network-layer
- specifies the receiver's address and waits for the network to set-up the VC.

The network-layer determines the path between sender and receiver.

The network-layer also determines rhe VC number for each link along the path. Finally, the network-layer adds an entry in the forwarding-table in each router.

• During VC setup, the network-layer may also reserve resources.

2) Data Transfer

Once the VC has been established, packets can begin to flow along the VC.

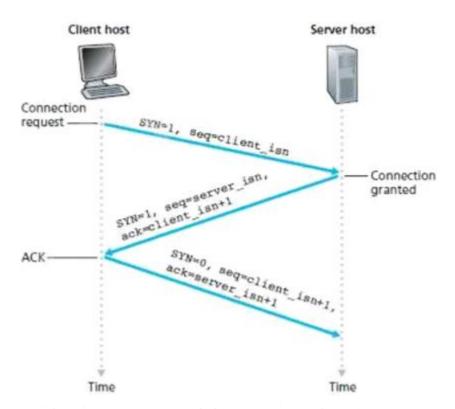
3) VC Teardown

This is initiated when the sender/receiver wants to terminate the VC. The network-layer

—. informs the other end-system of the call termination and

removes the appropriate entries in the forwarding-table in each router.

4. How does connection establishment and connection termination occur in connection oriented protocol?



TCP three-way handshake: segment exchange

Connection Setup & data Transfer

- To setup the connection, three segments are sent between the end hosts. The referred to as a three-way handshake.
- Suppose a client-process wants to initiate a connection wirh a server-process.
- Step 1: Client sends a connection-request segment to the Server The client first sends a connection-request segment to the server.
- The connection-request segment contains:
- 1) SYN bit is set to 1.
- 2) Znitial sequence -number (clients.isn).
- The SYN segment is encapsulated within an TCP datagram and sent to the server.
- Step 2: Server sends a connection-granted segment to the Client Then, the server
- extracts the SYN segment from the datagram allocates the buffers and variables to the connection and sends a connection-granted segment to the client.

The connection-granted segment contains:

1) SYN bit is set to 1.

- 2) Acknowledgment field is set to client_isn +1.
- 3) Initial sequence -number (server_isn).

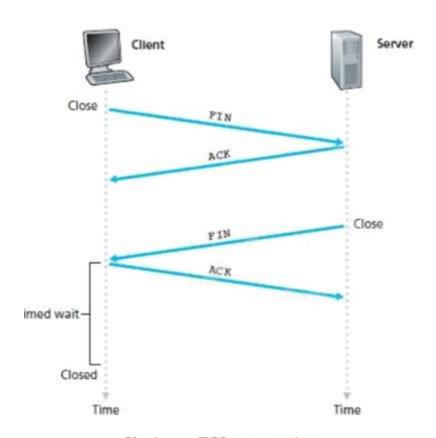
Step 3: Client sends an ACy segment to the Server

Finally, the client

allocates buffers and variables to tne connection and

sends an ACK segment to the server The ACK segment acxnowledges the server.

SYN bit is set to zero, since the connection is established.



Closing a TCP connection

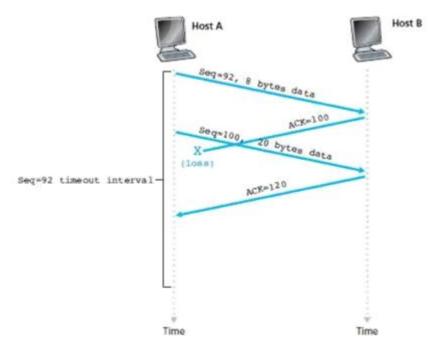
Connection Release

- Either of the two processes in a connection can end the connection.
- When a connection ends, the "resources'£ in the hosts are de -allocated.
- Suppose the client decides to close the connectiren.
- 1) The client-process issues a close command.

Then, the client sends a shutdown-segment to the server. x Th is segment has a FIN bit set to 1.

- 2) The server responds with an acknowledgment to the client.
- 3) The server then sends its own shutdown-segment x This segment has a FIN bit set to 1.
- 4} Finally, the client acknowledges the server's shutdown-segment.

4. b) Differentiate between ACK and SYN in connection oriented communication with example.



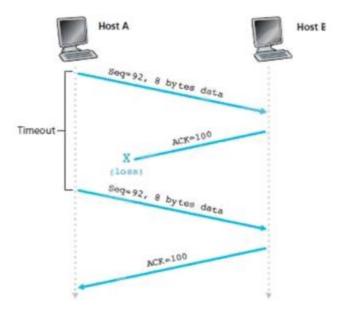
Flow dig to show ACK and SEQ

- ACK is the number of next expecting packet, in the above example, ACK= 100 is the packet number
 100 that is expected, next iteration the packet with SEQ number 100 is sent.
- SYN is the term used in order to ask end users to wait for reply from other end user before sending next packet.

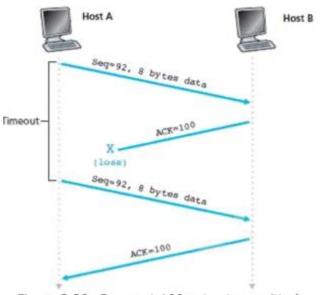
5. Design communication flow for reliable data transfer in following situations

- (i) Retransmission due to lost ACK.
- (ii) Sending multiple packets back-to-back.
- (iii) Avoiding retransmission by cumulative ACK.
- (iv) Fast retransmit.
- (v) Slow start.

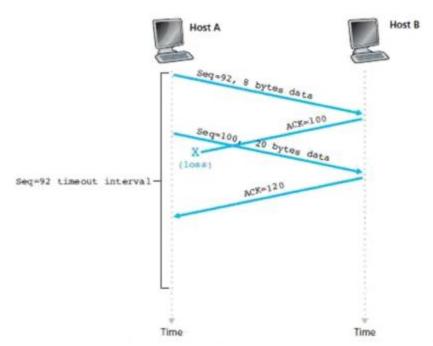
(i)



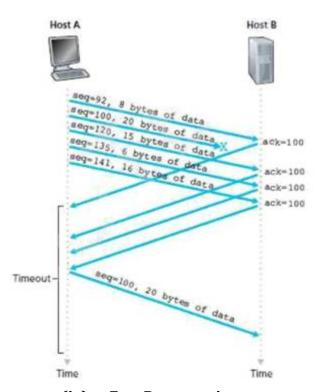
(i) Retransmission due to lost ACK.



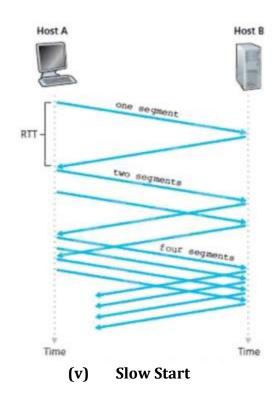
(ii) Sending multiple packets back-to-back.



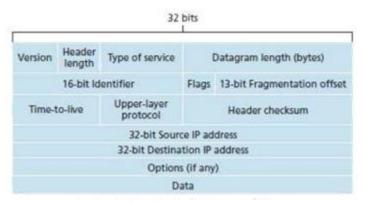
(iii) Avoiding retransmission by cumulative ACK.



(iv) Fast Retransmit



6 a) Draw and explain IPv4 datagram format.



IPv4 datagram format

Payload (or Data): This field contains the data to be delivered to the destination.

Header: Header Contains information essential to routing and delivery.

Version: This field specifies version of the 1Pv4 datagram, i.e. 4.

Header Length: This field specifies length of header. Without options field, header-ten gLh = 5 bytes.

Type of Service (TOS): This field specifies priority of packet based on parameters such as delay, throughput, reliability & cost.

datagram Length: This field specifies the total length of the datagram {header + data).

Maximum length = 6553d bytes.

Identifier, Flags, Fragmentation Offset

These fields are used for fragmentation and reassembly.

Fragmentation occurs when the size of the data9ram is larger than the MTU of the network.

i) **Identifier**: This field uniquely identifies a datagram packet.

ii) Flage: It is a 3-bit field. The first bit is not used.

The second bit D is called the do not fragment bit. The third bit M is called the more fragment bit.

- iii) fragmentation Offset: This field identifies location of a fragment in a datagram.
- **6) Time-To-Live (TTL):** This defines lifetime of the datagram (default value 64) in hops. Each router decrements TTL by 1 before forwarding. If TTL is zero, the datagram is discarded.
- 7) **Protocol:** This field specifies upper-layer protocol used to receive the datagram at the destination-nost. For example, TCP=6 and UDP=17.
- **8) Header Checksum:** This fiend is u sed to verify integrity of header. If the verification process fails, the packet is discarded.
- 9) Source IP Address & destination IP Address

These fields contain the addresses of source and destination respectively.

10) Options

This field allows the packet to request special features such as

- security level
- —• route to be taken by packet at each router.

6 b) List and describe 3 types of switching fabric.

Switching via Memory

- Switchin g b/w input-ports & output-ports is done under direct control of CPU i. e. routing processor.
- Input and output-ports work like a tradition al t/O devices in a computer.
- i) On arrival of a packet, the input port n otifies the routing-processor via an interrupt.
- ii) Then, the packet is copied from the input-port to processor-memory.
- iii) Finally, the routin g-processor extracts the destination-address from the header looks up the appropriate output-port in the forwarding-table and copies the packet into the output -port's buffers.
- Let memory-bandwidth = B packets per second.

Thus, the overall forwarding throughput must be less than B/2.

Switching via a Bus

- Switchin g b/yv input-p0rts & output-ports is done without intervention by the routing-processor.
- i) The input-port appends a switch-internal label (header) Lo the packet.
- The label indicates the local output-port to which the packet must be transferred. Then, the packet is received by all output-ports.
- But, only the port that matches the label will keep the packet.
- iii) Finally, the label is removed at the output-port.

Switching via an Interconnection Network

A crossbar switch is an interconnection nettwork.

- The network consists of 2N buses that connect N input-ports output-ports.
- Each vertical bus intersects each horizontal bus at a Crosspoint.
- The crosspoint can be opened or closed at any time by the switch-controller.
- 1) To move a packet from port A to port Y, the switch-controller closes the crosspoint at the intersection of buses A and Y.
- 2) Then, port A sends the packet onto its bus, which is picked up by bus Y.