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18EC641

Sixth Semester B.E. Degree Examination, July/August 2022 Operating System

Time: 3 hrs.

Max. Marks: 100

Note: Answer any FIVE full questions, choosing ONE full question from each module.

Module-1

- 1 a. Define operating system. Explain the goals of an operating system. (10 Marks)
 b. What are Computational structures? Explain the operation of an operating system. (10 Marks)

OR

- 2 a. Explain:
 i) Partition based
 ii) Pool based resource allocation strategies with a neat diagram. (10 Marks)
 b. Briefly explain the different classes of operating system, specifying the primary concern and key concepts used. (10 Marks)

Module-2

- 3 a. Explain the fundamental state transition for a process with state transition diagram. (10 Marks)
 b. For the following set of process perform FCFS and SRN scheduling. Calculate mean turn around time and mean weighted turn around.

Process	P ₁	P ₂	P ₃	P ₄	P ₅
Arrival time	0	2	3	5	9
Service time	3	3	2	5	3

(10 Marks)

OR

- 4 a. Define threads. Explain: i) User – level threads ii) Kernel level threads. (10 Marks)
 b. For the following set of process perform RR and LCN scheduling. (10 Marks)

Process	P ₁	P ₂	P ₃	P ₄	P ₅
Arrival time	0	2	3	5	9
Service time	3	3	2	5	3

Module-3

- 5 a. Explain contiguous and non contiguous memory allocation. (10 Marks)
 b. Explain: i) Paging ii) Segmentation. (10 Marks)

OR

- 6 a. Explain the important concepts in the operation of demand paging with diagram. (10 Marks)
 b. For the following page reference string apply FIFO and LRU page replacement policies to find number of page faults. Use alloc = 4.
 Page reference string: 5, 4, 3, 2, 1, 4, 3, 5, 4, 3, 2, 1, 5
 Reference time string: t₁, t₂, t₃, t₄, t₅, t₆, t₇, t₈, t₉, t₁₀, t₁₁, t₁₂, t₁₃.

(10 Marks)

Important Note : 1. On completing your answers, compulsorily draw diagonal cross lines on the remaining blank pages.
 2. Any revealing of identification, appeal to evaluator and /or equations written eg. 42+8 = 50, will be treated as malpractice.

Modified

Module-4

- 7 a. Explain the different file operations performed by processes. (10 Marks)
b. Explain the interface between file system and IOCS. (10 Marks)

OR

- 8 a. Explain the working of linked allocation of disk space with a figure. (10 Marks)
b. Explain implementing file access with neat figure. (10 Marks)

Module-5

- 9 a. Explain the primary issues in implementing message passing. (10 Marks)
b. Explain the working of blocking and non-blocking delivery protocols. (10 Marks)

OR

- 10 a. Define dead lock. Explain the condition of dead lock in resource allocation. (10 Marks)
b. Explain the three fundamental approaches used in dead lock handling. (10 Marks)

18 ECE 641
Sir, regarding Modification of Scheme and Solutions of ECE/ETE board

"Manjunatha P" <manjup.jnnce@gmail.com>

August 23, 2022 10:36 AM

To: boe@vtu.ac.in

Dear sir,

I am herewith sending the details of Comments from BoE of ECE Board for the following subjects towards Scheme and solution. Kindly download the attachment.

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Subject: Operating Systems Code: 18EC641 The following are the corrections required.

Q 9a. The answer for Q 9a is

Issues in message passing

Issue	Important aspects
Naming of processes	Processes participating in a message transfer are either explicitly indicated in <i>send</i> and <i>receive</i> statements, or they are deduced by the kernel in some other manner.
Method for transferring messages	Whether a sender process is blocked until the message sent by it is delivered, the order in which messages are delivered to the receiver process, and handling of exceptional conditions in message passing.
Kernel responsibilities	Buffering of messages pending delivery to recipient processes. Blocking and activation of processes.

generality of message passing. For example, requiring a sender process to know the identity of a receiver process may limit the scope of message passing to processes in the same application program. Relaxing this requirement would extend message passing to processes of different application programs and even to processes executing in different computer systems. Similarly, providing FIFO message delivery may be rather restrictive; processes may wish to receive messages in some order of their own choice.

System calls *send* and *receive* act as communication primitives. The number and meaning of their parameters are determined by the issues described in Table 10.1. We discuss various aspects of these primitives in the following Sections.

Q 6b For LRU policy - Number of page faults should be 8.

Other than this, as per the scrutiny from BoE members, there are no correction
Hence the same may be considered for the further process

Dr. Manjunatha. P
Chairman BoE for ECE
Professor & Dean Academics JNN College of Engineering
Shimoga

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Scheme & Solutions

Signature of Scrutinizer

Subject Title : Operating System

Subject Code : 18EC641

Question Number	Solution	Marks Allocated
1a	<p>An operating system is system software that manages computer hardware, software resources, and provides common services for computer programs. It is a software that controls the operation of computer and directs the processing of programs.</p> <p>Goal of operating system The key goal of an operating system is to provide a combination of efficient use and user convenience.</p> <p>Efficient use: Efficient use of resources can be obtained by monitoring the use of resources and performing corrective actions when necessary.</p> <p>However, a computer system contains several resources like memory, I/O devices and the CPU. So a comprehensive method of monitoring efficiency may lead to high overhead. Consequently, operating system employ simple and easy to apply but known-to-be-suboptimal strategies for ensuring good efficiency, e.g. they either focus on efficiency of a few important resources like the CPU and memory, or handle users programs in a manner that guarantees high efficiency.</p>	(3)

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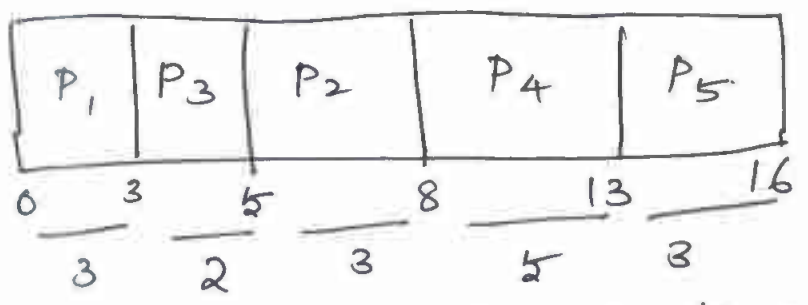
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Question Number	Solution	Marks Allocated
	<p>Used Convenience: User Convenience has several facets</p> <p>Examples</p> <p>(1) Necessity Ability, to execute programs, use the file system.</p> <p>(2) Good service speedy response to computational requests</p> <p>(3) user friendly or Easy-to-use commands, graphical-user interface (GUI)</p> <p>(4) New programming model Concurrent programming</p> <p>(5) Features for experts Means to Set up complex computational structures</p> <p>(6) Web-Oriented Services Means to set up web-enabled servers.</p>	<p>(3+7 = 10m)</p>
1b	<p>Computational structures: A Computational Structure is a configuration of one or more Programs that work towards a common goal. Some typical Computational Structures are</p> <ul style="list-style-type: none"> * A single Program * A sequence of single programs * Co-executing Programs. 	<p>(2m)</p>
	<p>operation of an OS An operating system implements computational requirements of its users with the help of resources of the computer system.</p> <p>Key Concepts of an operating system</p> <p>(1) Programs: Initiation and termination of programs. Providing convenient methods so that several programs can work towards a common goal.</p> <p>Resources: Ensuring availability of resources in the system and allocating them to programs.</p> <p>Scheduling: Deciding when, and for how long, to</p>	<p>(8m)</p>

Question Number	Solution	Marks Allocated																								
	<p>devote the CPU, to a program.</p> <p>Protection: protect data, and programs against interference from other users and their programs.</p> <p>Common tasks performed by operating systems.</p> <ol style="list-style-type: none"> 1. Maintain a list of authorized users. 2. Construct a list of all resources in the system. 3. Initiate execution of programs. 4. Maintain resource usage information by programs and current status of all programs. 5. Maintain current status of all resources and allocate resources to programs when requested. 6. Perform scheduling for protection. 7. Maintain information for protection. 8. Handle requests made by users and their programs. 	<p>(2+8) = (10m)</p>																								
2a	<p>Partition based: The OS decides a priori what resources should be allocated to a user program. This approach is called static Allocation because the allocation is made before the execution of a program begins.</p>	(5m)																								
	<p>Pool Based: The OS maintains a common pool of resources and allocates from this pool whenever a program requests a resource. This approach is called dynamic Allocation because allocation takes place during execution of a program.</p>	(5m) (5+5) = 10m																								
2b	<table border="0"> <tr> <td>OS class</td> <td>Period</td> <td>Prime concern</td> <td>Key concepts</td> </tr> <tr> <td>Batch processing</td> <td>1960's</td> <td>CPU utilization</td> <td>Spooling, Command processors</td> </tr> <tr> <td>Multi programming</td> <td>1970's</td> <td>Resource utilization</td> <td>Program priorities, prompts</td> </tr> <tr> <td>Timesharing</td> <td>1970s</td> <td>Good response time.</td> <td>Time slice, round-robin scheduling.</td> </tr> <tr> <td>Real time</td> <td>1980s.</td> <td>Meet the deadline</td> <td>Real time scheduling</td> </tr> <tr> <td>Distributed</td> <td>1990's</td> <td>Resource sharing</td> <td>Transparency, distributed control</td> </tr> </table>	OS class	Period	Prime concern	Key concepts	Batch processing	1960's	CPU utilization	Spooling, Command processors	Multi programming	1970's	Resource utilization	Program priorities, prompts	Timesharing	1970s	Good response time.	Time slice, round-robin scheduling.	Real time	1980s.	Meet the deadline	Real time scheduling	Distributed	1990's	Resource sharing	Transparency, distributed control	<p>Each (2m) 5 x 2 = (10m)</p>
OS class	Period	Prime concern	Key concepts																							
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Question Number	Solution	Marks Allocated
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shortest request next



(5+5)
(10m)

Turn Around time = Completion time - Arrival time

- $P_1 \rightarrow t_a = 3 - 0 = 3 \text{ sec}$
- $P_3 \rightarrow t_a = 5 - 3 = 2 \text{ sec}$
- $P_2 \rightarrow t_a = 8 - 2 = 6 \text{ sec}$
- $P_4 \rightarrow t_a = 13 - 5 = 8 \text{ sec}$
- $P_5 \rightarrow t_a = 16 - 9 = 7 \text{ sec}$

Average turn Around time = $\frac{3+2+6+8+7}{5} = \frac{26}{5} = 5.2 \text{ sec}$

Weighted Average turn

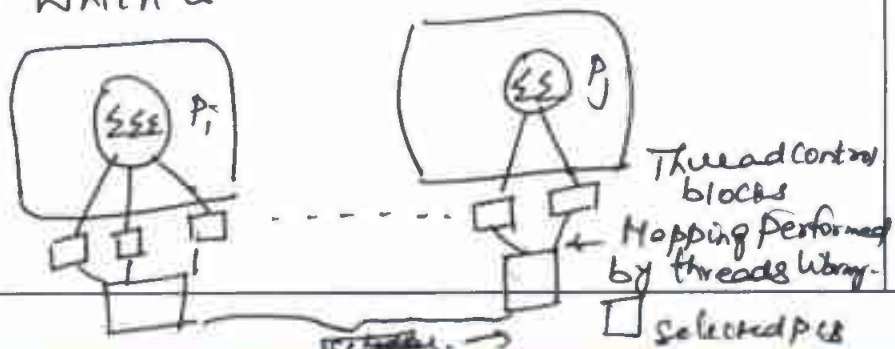
Around time (W) =

- $P_1 = 3/3 = 1$; $P_3 = 2/2 = 1$; $P_2 = 6/3 = 2$;
- $P_4 = 8/5 = 1.60$; $P_5 = 7/3 = 2.33$

Average weighted turn Around time = $\frac{1+1+2+1+1.60+2.33}{5} = 1.59$

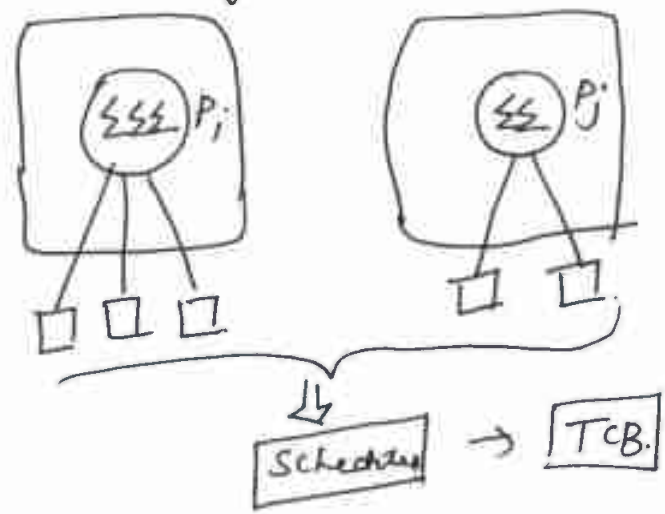
Q1a Threads: A thread is a program execution that uses the resources of a process. (2m)

User level Thread: Are implemented by a thread library which is linked to the code of a process (4m)



Question Number	Solution	Marks Allocated
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Kernel-Level Thread: A kernel level thread is implemented by kernel.



(4m)
(2+4+4)
=(10m)

4b Round Robin scheduling

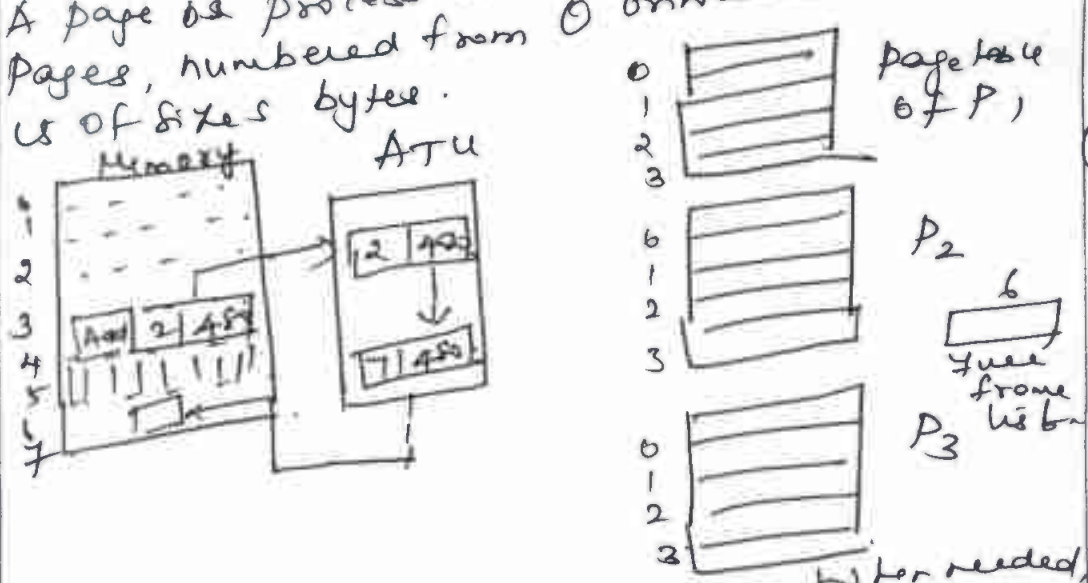
P	AT	BT	W	CT	$\frac{CT}{BT-AT}$
P ₁	0	3	$\frac{4}{3} = 1.333$	4	$4 - 0 = 4$
P ₂	2	3	$\frac{7}{3} = 2.33$	9	$9 - 2 = 7$
P ₃	3	2	$\frac{5}{2} = 2.50$	8	$8 - 2 = 5$
P ₄	5	5	$\frac{11}{5} = 2.20$	16	$16 - 5 = 11$
P ₅	9	3	$\frac{6}{3} = 2$	15	$15 - 9 = 6$

(5m)

LCN scheduling

P	c	ba	waiting time	waited time around time
P ₁	9	0-9=9	9-3=6	$\frac{9}{3} = 3$
P ₂	12	12-2=10	10-3=7	$\frac{10}{3} = 3.33$
P ₃	7	7-3=4	4-2=2	$\frac{4}{2} = 2$
P ₄	16	16-5=11	11-5=6	$\frac{11}{5} = 2.20$
P ₅	14	14-9=5	5-3=2	$\frac{5}{3} = 1.67$

(5m)

Question Number	Solution	Marks Allocated
5a	<p>Contiguous Memory Allocation : Each process is allocated a single contiguous area in memory. In early computer systems the memory allocation decision was made stoniscolly, i.e before the execution of a process began. Write the memory structure</p> <p>Non-Contiguous Memory Allocation : Several non-adjacent memory areas are allocated to a process. None of these areas is large enough to hold the complete process. Write the memory allocation structure.</p>	<p>(5+5) =10m.</p>
5b	<p>Paging : In paging each process consist of fixed size components called pages.</p> <p>Segmentation : A programmer identifies components called segments in a process. A segment is a logical entry in a program. eg a function, a data structure or an object.</p>	<p>(5+5) =10m</p>
6a	<p>Demand paging : A page of process is moved to consist of pages, numbered from 0 onwards. Each page is of size s bytes.</p>  <p>The diagram illustrates the demand paging mechanism. On the left, a 'Memory' block is divided into slots 0-7. An 'ATU' (Address Translation Unit) is shown with a table containing logical addresses 2 and 4, which are mapped to physical addresses 1 and 5. To the right, three 'Page Table' structures are shown for processes P1, P2, and P3. P1's page table has entries for logical pages 0, 1, 2, and 3. P2's page table has entries for logical pages 0, 1, 2, and 3. P3's page table has entries for logical pages 0, 1, 2, and 3. A note indicates that pages 4, 5, 6, and 7 are 'free from P3's table'. Below the page tables, it is noted that pages are loaded 'when needed'.</p> <p>A page is loaded in memory when a logical address is generated by process points to a page that is not present in memory.</p>	<p>Expt-8 Qn-2 (8+2) (10m)</p>

Question Number	Solution	Marks Allocated
<p>6b</p>	<p>FIFO + LRU Page replacement process.</p> <p>FIFO Alloc = 4</p> <p>Page fault : 10, Page hit : 3</p>	<p>(5m)</p>
<p>7a</p>	<p>LRU. (Alloc = 4)</p> <p>Page fault → 11, Page hit → 2</p> <p>Module - 4</p>	<p>(5m) (5+5=10m)</p>
<p>7a</p>	<p>(1) opening a file (2) Reading or writing (3) closing (4) file creation (5) file deletion (6) file renaming</p> <p>Explanation is omitted</p>	<p>2 marks Each (5*2=10m)</p>
<p>7b</p>	<p>Interface between the system & I/O & uses I/O mechanism to implement I/O operation. These mechanism uses I/O programming to handle device level details of I/O without interrupt handling</p> <p>Diagram - 2</p> <p>Expt - 8</p>	<p>(Expt-8 Dia-2 (8+2) 10m)</p>
<p>8a</p>	<p>Linked Allocation of disk space. A file is represented by a linked list of disk blocks</p> <p>Diagram 2</p> <p>Expt - 8</p>	<p>(8+2) 10m</p>

Question Number	Solution	Marks Allocated
8b	<p>Implementing File Access: Use of directory structure and file control block (FCB) in implementing a file access. The scheme discussed there uses an FCB to keep track of current state of a file processing activity.</p> <p>Diagram → 2</p> <p>Expt → 8</p>	<p>(8+2) 10m.</p>
9a	<p>Implementing Message Passing: When a process P1 sends a message to some process P2 using nonblocking send, the kernel builds on inter process message control block (IMCB) to store all information.</p> <p>Diagram → 2</p> <p>Expt - 8</p>	<p>(8+2) 10m</p>
9b	<p>Blocking and Non blocking protocols</p> <p>de Wreey</p> <p>Diagram → 3</p> <p>Expt - 7</p>	<p>(7+3)</p>
10a	<p>Deadlock: A set of processes is deadlocked if each of them waits for an event that can be caused only by process in the set.</p> <p>Conditions of deadlock in Resource Allocation</p> <p>Expt + Diagram → (5+2)</p>	<p>3+7 = 10m</p>
10b	<p>Fundamental Approaches used in Deadlock Handling</p> <p>Three fundamental Approaches used: (1) Deadlock detection & Resolution. Deadlock Prevention. Deadlock Avoidance.</p>	<p>(10m)</p>

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