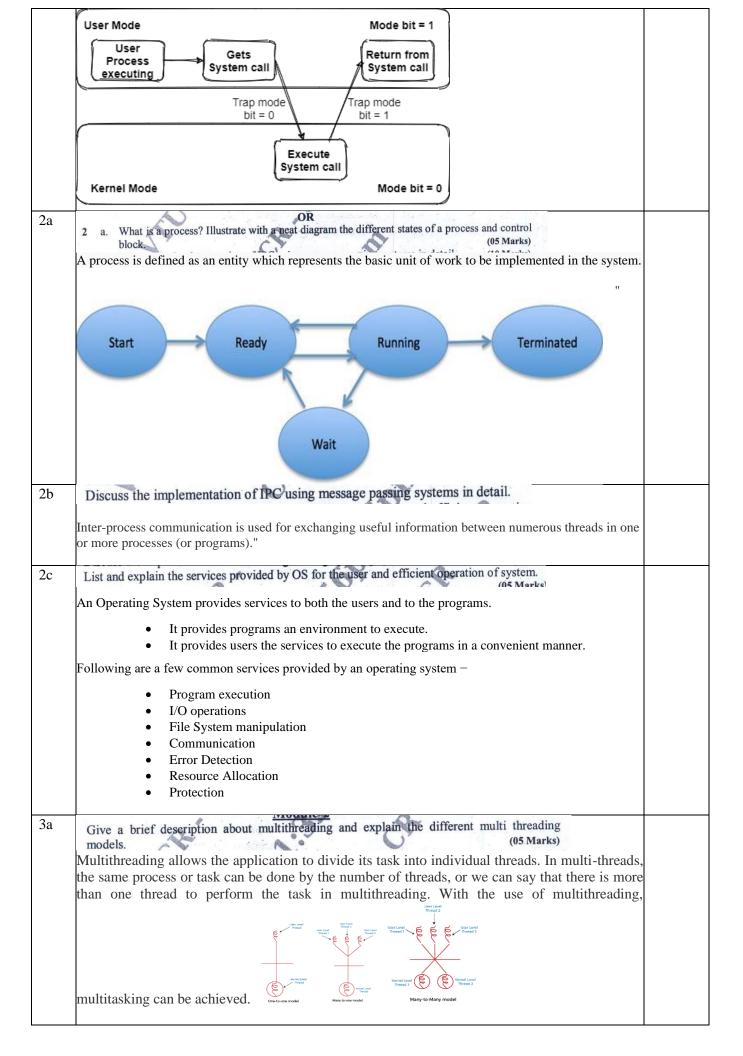
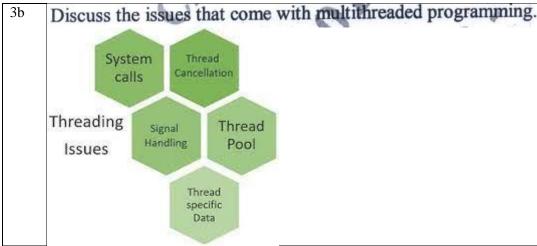
Semester End Examination – July/August 2022 Scheme and Solution Faculty-Divya Singh/Rashmi D

	Operating	g System				Sub Code:	18CS43	Branch:	ISE		
ate:	14/10/202	2 Duration:	3hrs N	Max Marks:	100	Sem/Sec:	IV A, B & C	2			
	0	h between –						MA	ARKS		
	MULTIPROGRAMMING AND MULTITASKING										
	Parameter	Multiprog	ramming	м	lultitasking						
	Definition	It lets multiple programs use	the CPU at the same tin	ne. execution of numer	efers to the simul rous programme th the certain tin	s, processes,					
	Objective	It's useful for cutting down or throughput as m		once, significantly							
	Mechanism	The context switch	ng method is used.	Based on a ti	me-sharing mech	nanism.					
	Time	Multiprogramming demands complete		to Multitasking allows	you to complete time.	tasks in less					
	Execution	In a multi-programmed syste completes its execution or cha temporarily suspends that proce execute from the pro	nges to an I/O task, the system is a second of the system of the system is a second of the syste	stem system by assigning							
	CPU Switching	The CPU shifts between proces	ses swiftly in this environr	ment. The CPU shifts between	ween the process single-user envir	ses of several onment.					
	CPU required	Only one CPU is needed in Mu	tiprogramming to run the t		-						
		1	Multiprogramming v	rs. Multitasking							
	-	ter. The applic	_	_		ther applica					
	request addition interface	ts for services in, users can ice, such as a co	through a interact dire	defined appectly with t	plication the ope	ne operating n program crating syst	interface (<u>A</u> em through	<u>API</u>). In			
	reques additio	n, users can ce, such as a co	through a interact dire mmand-line	defined appectly with t	plication the ope CLI) or	ne operating n program trating syst a graphical	interface (<u>A</u> em through	<u>API</u>). In			





3c Explain CPU scheduling criteria

The scheduling criterion is responsible for helping in the design of the good scheduler. These criteria are as follows –

CPU Utilization

The scheduling algorithm should be designed in such a way that the usage of the CPU should be as efficient as possible.

Throughput

It can be defined as the number of processes executed by the CPU in a given amount of time. It is used to find the efficiency of a CPU.

Response Time

The Response time is the time taken to start the job when the job enters the queues so that the scheduler should be able to minimize the response time.

Response time = Time at which the process gets the CPU for the first time - Arrival time

Turnaround time

Turnaround time is the total amount of time spent by the process from coming in the ready state for the first time to its completion.

Turnaround time = Burst time + Waiting time

or

Turnaround time = Exit time - Arrival time

Waiting time

The Waiting time is nothing but where there are many jobs that are competing for the execution so the Waiting time should be minimized.

Waiting time = Turnaround time - Burst time

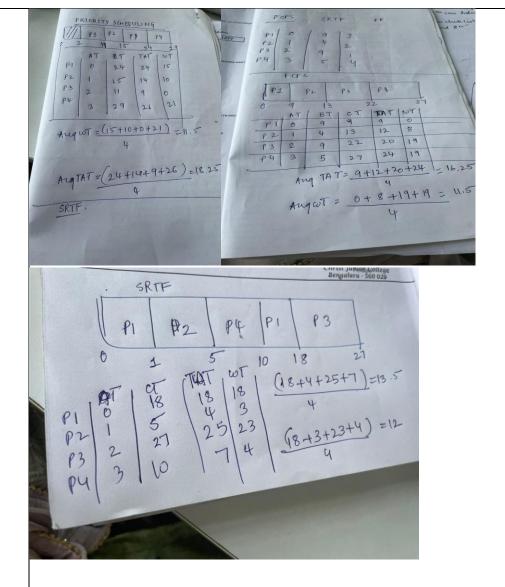
Fairness

For schedulers, there should be fairness in making sure that the processes get a fair share of chances to be executed.

4 a. Calculate the average waiting time and the average turnaround time by drawing the Gantt chart using FCFS, SRTF, RR (q = 2ms) and priority algorithms. Lower priority number represents higher priority.

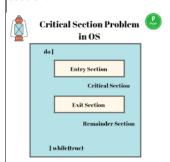
√P1	ocess	Arrival Time	Burst Time	Priority
	Pı	0_0	9	3
	P ₂	1 1	4	2
17	P ₃	2	9	1
	P4 4	3	5	4

(12 Marks)



b. What is critical section problem? What are the requirements for the solution to critical section problem? Explain Peterson's solution. (08 Marks)

The critical section problem is **used to design a protocol followed by a group of processes**, so that when one process has entered its critical section, no other process is allowed to execute in its critical section.



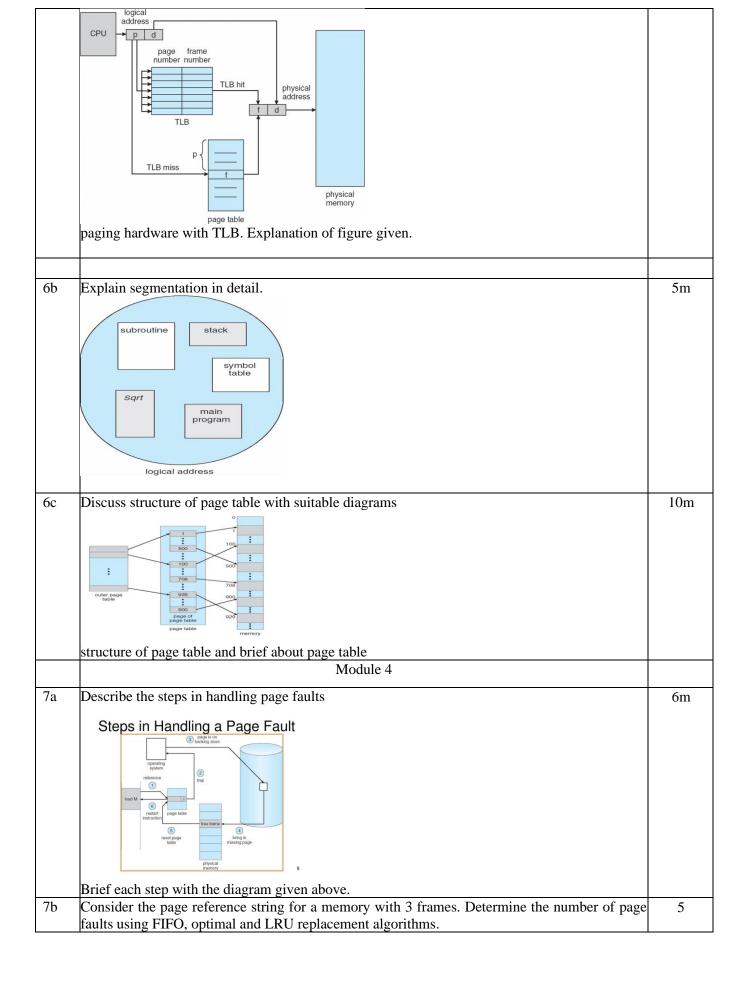
Peterson's algorithm (or Peterson's solution) is a concurrent programming algorithm for mutual exclusion that allows two or more processes to share a single-use resource without conflict, using only shared memory for communication. It was formulated by Gary L. Peterson in 1981.

```
turn = 0;
                          flag[0] = false;
flag[1] = false;
        while(true){
                                      while(true){
         flag[0] = true;
                                       flag[1] - true;
         turn=1;
                                       turn = 0:
                                        while (flag[0] and turn ==0) do no-op;
         while (flag[1] and turn == 1) do no-op;
          CS
                                        CS
                                       flag[1] = false;
         flag[0] = false;
         remainder section;
                                       remainder section:
5a
       What is a deadlock? What are the necessary conditions for the deadlock to occur? (
      Deadlock is a situation where two or more processes are waiting for each other. For example, let us
      assume, we have two processes P1 and P2. Now, process P1 is holding the resource R1 and is waiting
      for the resource R2. At the same time, the process P2 is having the resource R2 and is waiting for the
      resource R1. So, the process P1 is waiting for process P2 to release its resource and at the same time,
      the process P2 is waiting for process P1 to release its resource. And no one is releasing any resource.
      So, both are waiting for each other to release the resource. This leads to infinite waiting and no work is
      done here. This is called Deadlock.
      Necessary conditions-
                  1. Mutual exclusiom
                  2. No pre-emption
                  3. Hold and wait
                  4. Circular wait
        How to prevent the occurrence of deadlock, explain in detail.
5<sub>b</sub>
      Banker's algorithms
       Consider the following snapshot of a system:
5c
                        Process
                                   Allocation
                                                     Max
                                                                  Available
                                         CD
                                                    BC
                                                                   В
                                                                       C
                                  AB
                                                A
                                                            D
                                                    2 1
                                                            2
                                  2
                                      0
                                         0
                                             1
                                  3
                                                            2
                                         2
                                                    2
                                                        5
                                     1
                                             1
                                                 5
                                  2
                                     1
                                         0
                                             3
                                                    3
                                                        1
                                                            6
                                  1
                                     3
                                         1
                                                    4
                                                        2
                                                            4
                                     4 3
                           PA
      Answer the following using Banker's algorithm.
            Is the system in safe state? If so, give the safe sequence.
            If process P<sub>2</sub> requests (0, 1, 1, 3) resources can it be granted immediately? (
      ii)
                      The Banker's Algorithm
       Handles multiple instances for resource types.
       n = \text{number of processes}; m = \text{number of resource types};
       Data Structures:

    Available: vector [1..m], Available[]] - the number of

                instances currently available for resource j.
               Max matrix[1_n, 1_m], Max[i, j] - the maximum number of
                instances of resource j that process i can request at any one
               Allocation: matrix[1..n, 1..m] - process i currently holds
                Allocation[ i, j ] instances of resource j.
               Need: matrix[1..n, 1..m] - process i may need
               Need[i, j] instances of resource j.
                     Need[i,j] = Max[i,j] - Allocation[i,j]
      air Amir
                                      Fall 00 / Lecture 4
      Explain paging hardware with TLB
6a.
                                                                                                           5<sub>m</sub>
```

Peterson's solution



	FIFO	FIRST IN FIR	ST OUT															
	PAGE DEMAND	6	2	3	4	2 1	1 5 6	2	1 2	3 7	6 3	4 1	2	7	6			
	f1 nory f2 f3	6	2	2	2	2 2	1 1 1 2 5 5	5	5 5	3 3 5 7	3 3 7 7	3 3 7 7	7	7	6			
	f4	fault fi	ault fau	ult fault	4 hit	4 4	4 4 4 4 fault	2 fault hit	2 2	2 2	2 2	2 1	1 ault hit	1 fault	1			
	REFERENC	E		10011		10011				The state of the s								
	STRING PAGE	20																
	PAGE HITS																	
	OPTIMAL PAGE REPL	ACEMENT																
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	MM f2 f3		2	2	2	2 2	2 2 2	2	2 2	2 7	7 7 3	7 7	7 2	7 2	7			
	f4	fault fa	ault fau	ult fault	4 hit	4 1 fault	1 1 1 fault hit	1 hit hit	1 1 hit fault	1 1 fault hit	1 1 hit faul	1 1	1 ault hit	1 hit	1			
		PAGE FAULT		10														
		PAGE HIT		10														
	LRU LEAST RECENTLY	USED	2	2	4 :	2 1	1 5 6	2	1 2	2 7	6 2	4 1	2	7	6			
	f1	6	6	6	6	6 1		1	1 1	1 1	6 6	6 6	2	2	2			
	f2 f3		2	2	3	2 2 3 3	2 2 2 3 5 5	2 5	2 2 5 5	2 2 3	2 2 3 3	4 4 3 3	4	4 7	6 7			
	f4	fault fa	ault fau	ult fault	4 hit	4 4	4 4 6 fault fault	6 hit hit	6 6 fault	6 7 fault fault	7 7 hit faul	7 1 fault fa	1 oult faul	1 fault	1			
		PAGE FAULT		15														
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	Contig																	
	Linked																	
	Indexe	d Al	loca	tion														
b	Discus	s the	var	ious	dire	ctor	v struct	ures v	vith rec	uired o	diagrai	ns.						10
					Directory	1	J			L	υ							
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	F1 F2	F3	7 [F4	FS	<u></u>	F6 F7	F8	F9									
	Files																	
	Directo	ory si	truct	tures	exp	alar	ntion											
		•						N	Iodule :	5								
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b	A drive																	10m
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												nove t	o sat	ısty	an a	pendi	ng	
	position, what is a total distance(in cylinders) that the Disk arm move to satisfy all a pending requests using FCFS, SSTF, LOOK AND C-LOOK Algorithms?																	
			-6 - 1	,						_							ļ	
	FCFS						just fol		na ordan			o oiro	,					

The FCFS schedule is:

143, 86, 1470, 913, 1774, 948, 1509, 1022, 1750, 130

The total distance is: 7,081

SSTF

The SSTF algorithm starts a cylinder 143 and from there successively selects the shortest request from its current location.

The SSTF schedule is:

143, 130, 86, 913, 948, 1022, 1470, 1509, 1750, 1774

The total distance is: 1,745

LOOK

The LOOK algorithm is just like the SCAN algorithm, except the disk head only goes as far as the last request in each direction.

The LOOK schedule is:

143, 913, 948, 1022, 1470, 1509, 1750, 1774, 130, 86

The total distance is: 3,319

C-LOOK

The C-LOOK algorithm is a circular version of the LOOK algorithm. It doesn't scan on the way back to the beginning of the disk, rather operates in a circular fashion.

The C-LOOK schedule is:

143, 913, 948, 1022, 1470, 1509, 1750, 1774, 86, 130

The total distance is: 3,363

10a Write a neat diagram, explain the components of a Linux system

system- management programs	user processes	user utility programs	compilers								
	system shared libraries										
	Linux kernel										
	loadable ker	nel modules									

Kernel. The kernel is responsible for maintaining all the important abstractions of the operating system, including such things as virtual memory and processes.

System libraries. The system libraries define a standard set of functions through which applications can interact with the kernel. These functions implement much of the operating-system functionality that does not need the full privileges of kernel code.

8m

	System utilities. The system utilities are programs that perform individual, specialized						
	management tasks. Some system utilities may be invoked just once to initialize and						
	configure some aspect of the system; othersknown as daemons in						
10b	Explain the different IPC mechanisms available in Linux.	6m					
	Synchronization and Signals						
	Passing of Data Among Processes						
10c	Discuss about scheduling in Linux	6m					
	Process Scheduling						
	Kernel Synchronization						
	Symmetric Multiprocessing						