



t ū t  $\Omega$ r

 $\ddagger$ 'n m  $\Omega$ u  $\mathbf{r}$ 

A dash in the upper string indicates insertion. A substitution occurs when the two alignment symbols do not match (shown in bold). We can associate a weight or cost with each operation. The Levensthein distance between two sequences is obtained by assigning a unit cost to each operation. Another possible alignment for this sequences is:

t  $\mathbf{u}$  $\mathbf t$  $\overline{O}$  $\mathbf{r}$ t  $\mathbf u$  ${\bf m}$  $\Omega$  $\mathbf u$  $\mathbf{r}$ 

which has a cost of 3. We already have a better alignment than this one.

The problem of finding minimum edit distance seems quite simple but in fact is not so. A choice that seems good initially might lead to problems later. Dynamic programming algorithms can be quite useful for finding minimum edit distance between two sequences. Dynamic programming refers to a class of algorithms that apply a table-driven approach to solve problems by combining solutions to sub-problems. The dynamic programming algorithm for minimum edit distance is implemented by creating an edit distance matrix. This matrix has one row for each symbol in the source string and one column for each matrix in the target string. The  $(i, j)$ th cell in this matrix represents the distance between the first i character of the source and the first  $j$  character of the target string. Each cell can be computed as a simple function of its surrounding cells. Thus, by starting at the beginning of the matrix, it is possible to fill each entry iteratively. The value in each cell is computed in terms of three possible paths.

$$
dist[i, j] = \begin{cases} dist[i-1, j] + insert\_cost, \\ dist[i-1, j-1] + subst\_cos t[source_i, targ et_j] \\ dist[i, j-1] + delete\_cost \end{cases}
$$



Figure 3.13 Minimum edit distance algorithm









rules are used to expand the tree, which gives us two partial trees at the second level search, as shown in Figure 4.4. The third level is generated by expanding the non-terminal at the bottom of the search tree in the previous ply. Due to space constraints, only the expansion corresponding to the left-most non-terminals has been shown in the figure. The subsequent steps in the parse are left, as an exercise, to the readers. The correct parse tree shown in Figure 4.4 is obtained by expanding the fifth parse tree of the third level.

## $4.4.2$ **Bottom-up Parsing**

A bottom-up parser starts with the words in the input sentence and attempts to construct a parse tree in an upward direction towards the root. At each step, the parser looks for rules in the grammar where the right hand side matches some of the portions in the parse tree constructed so far, and reduces it using the left hand side of the production. The parse is considered successful if the parser reduces the tree to the start symbol of the grammar. Figure 4.5 shows some steps carried out by the bottomup parser for sentence  $(4.7)$ .



**CI** 

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**HOD**