

be heterogeneous in nature because a variety of resources, such as clusters and even networked PCs, can be used to build it.

 The physical infrastructure is managed by the core middleware, the objectives of which are to provide an appropriate runtime environment for applications and to best utilize resources. At the bottom of the stack, virtualization technologies are used to guarantee runtime environment customization, application isolation, sandboxing, and quality of service. Hardware virtualization is most commonly used at this level. Hypervisors manage the pool of resources and expose the distributed infrastructure as a collection of virtual machines. By using virtual machine technology it is possible to finely partition the hardware resources such as CPU and memory and to virtualize specific devices, thus meeting the requirements of users and applications. This solution is generally paired with storage and network virtualization strategies, which allow the infrastructure to be completely virtualized and controlled.

 Infrastructure management is the key function of core middleware, which supports capabilities such as negotiation of the quality of service, admission control, execution management and monitoring, accounting, and billing.

The combination of cloud hosting platforms and resources is generally classified as an Infrastructure-as-a-Service (IaaS) solution. We can organize the different examples of IaaS into two categories: Some of them provide both the management layer and the physical infrastructure; others provide only the management layer (IaaS (M)).

 In this second case, the management layer is often integrated with other IaaS solutions that provide physical infrastructure and adds value to them.

IaaS solutions are suitable for designing the system infrastructure but provide limited services to build applications. Such service is provided by cloud programming environments and tools, which form a new layer for offering users a development platform for applications.

 The range of tools include Web-based interfaces, command-line tools, and frameworks for concurrent and distributed programming. In this scenario, users develop their applications specifically for the cloud by using the API exposed at the user-level middleware. For this reason, this approach is also known as Platform-as-a-Service (PaaS) because the service offered to the user is a development platform rather than an infrastructure.

 The top layer of the reference model depicted in Figure 4.1 contains services delivered at the application level. These are mostly referred to as Software-as -a-Service (SaaS). In most cases these are Web-based applications that rely on the cloud to provide service to end users. The horsepower of the cloud provided by IaaS and PaaS solutions allows independent software vendors to deliver their application services over the Internet.

 Table 4.1 summarizes the characteristics of the three major categories used to classify cloud computing solutions. In the following section, we briefly discuss these characteristics along with some references to practical implementations.

1. Execution Virtualization

Execution virtualization includes all techniques that aim to emulate an execution environment that is separate from the one hosting the virtualization layer. All these techniques concentrate their interest on providing support for the execution of programs, whether these are the operating system, a binary specification of a program compiled against an abstract machine model, or an application.

- 1. Machine reference model
- 2. Hardware-level virtualization
	- a. Hypervisors
	- b. Hardware virtualization techniques
	- c. Operating system-level virtualization
- 3. Programming language-level virtualization
- 4. Application-level virtualization

1) Machine Reference Model

Modern computing systems can be expressed in terms of the reference model described in Figure 3.4. At the bottom layer, the model for the hardware is expressed in terms of the Instruction Set

Architecture (ISA), which defines the instruction set for the processor, registers, memory, and interrupt management. ISA is the interface between hardware and software, and it is important to the operating system (OS) developer (System ISA) and developers of applications that directly manage the underlying hardware (User ISA). The application binary interface (ABI) separates the operating system layer from the applications and libraries, which are managed by the OS. ABI covers details such as low-level data types, alignment, and call conventions and defines a format for executable programs.

The highest level of abstraction is represented by the application programming interface (API), which interfaces applications to libraries and/or the underlying operating system.

abstract execution environment in terms of computer hardware on top of which a guest operating system can be run. In this model, the guest is represented by the operating system, the host by the physical computer hardware, the virtual machine by its emulation, and the virtual machine manager by the hypervisor

The evolution of distributed computing technologies, 1950s-2010s.

(b) Briefly explain each of these technologies with suitable diagrams.

Clouds are essentially large distributed computing facilities that make available their services to third parties on demand. As a reference, we consider the characterization of a distributed system proposed by Tanenbaum et al.:

"A distributed system is a collection of independent computers that appears to its users as a single coherent system."

Three major milestones have led to cloud computing: mainframe computing, cluster computing, and grid computing.

1. Mainframes: These were the first examples of large computational facilities leveraging multiple processing units. Mainframes were powerful, highly reliable computers specialized.

for large data movement and massive input/output (I/O) operations. They were mostly used by large organizations for bulk data processing tasks such as online transactions, enterprise resource planning, and other operations involving the processing of significant amounts of data.

2. Clusters: Cluster computing started as a low-cost alternative to the use of mainframes and supercomputers. The technology advancement that created faster and more powerful mainframes and supercomputers eventually generated an increased availability of cheap commodity machines as a side effect. These machines could then be connected by a high-bandwidth network and controlled by specific software tools that manage them as a single system. Starting in the 1980s.

Cluster technology contributed considerably to the evolution of tools and frameworks for distributed computing, including Condor, Parallel Virtual Machine (PVM), and Message Passing Interface (MPI).

3. Grid Computing: Grid computing appeared in the early 1990s as an evolution of cluster computing. In an analogy to the power grid, grid computing proposed a new approach to access large computational power, huge storage facilities, and a variety of services.

layer. The guest represents the system component that interacts with the virtualization layer rather than with the host, as would normally happen. The host represents the original environment where the guest is supposed to be managed.

(b) Discuss the reference model of full virtualization with a neat diagram.

3.6.2.1 Full virtualization and binary translation

VMware is well known for the capability to virtualize x86 architectures, which runs unmodified on top of their hypervisors. With the new generation of hardware architectures and the introduction of hardware-assisted virtualization (Intel VT-x and AMD V) in 2006, full virtualization is made possible with hardware support, but before that date, the use of *dynamic binary translation* was the only solution that allowed running x86 guest operating systems unmodified in a virtualized environment.

As discussed before, x86 architecture design does not satisfy the first theorem of virtualization, since the set of sensitive instructions is not a subset of the privileged instructions. This causes a different behavior when such instructions are not executed in Ring 0, which is the normal case in a virtualization scenario where the guest OS is run in Ring 1. Generally, a trap is generated and the way it is managed differentiates the solutions in which virtualization is implemented for x86 hardware. In the case of dynamic binary translation, the trap triggers the translation of the offending instructions into an equivalent set of instructions that achieves the same goal without generating exceptions. Moreover, to improve performance, the equivalent set of instruction is cached so that translation is no longer necessary for further occurrences of the same instructions. Figure 3.12 gives an idea of the process.

This approach has both advantages and disadvantages. The major advantage is that guests can run unmodified in a virtualized environment, which is a crucial feature for operating systems for which source code is not available. This is the case, for example, of operating systems in the Windows family. Binary translation is a more portable solution for full virtualization. On the other hand, translating instructions at runtime introduces an additional overhead that is not present in other approaches (paravirtualization or hardware-assisted virtualization). Even though such disadvantage exists, binary translation is applied to only a subset of the instruction set, whereas the others are managed through direct execution on the underlying hardware. This somehow reduces the impact on performance of binary translation.

CPU virtualization is only a component of a fully virtualized hardware environment. VMware achieves full virtualization by providing virtual representation of memory and I/O devices. Memory virtualization constitutes another challenge of virtualized environments and can deeply impact performance without the appropriate hardware support. The main reason is the presence of a *memory* management unit (MMU), which needs to be emulated as part of the virtual hardware. Especially in the case of *hosted hypervisors* (Type II), where the virtual MMU and the host-OS MMU are traversed sequentially before getting to the physical memory page, the impact on performance can be significant. To avoid nested translation, the *translation look-aside buffer (TLB)* in the virtual MMU directly maps physical pages, and the performance slowdown only occurs in case of a TLB miss.

