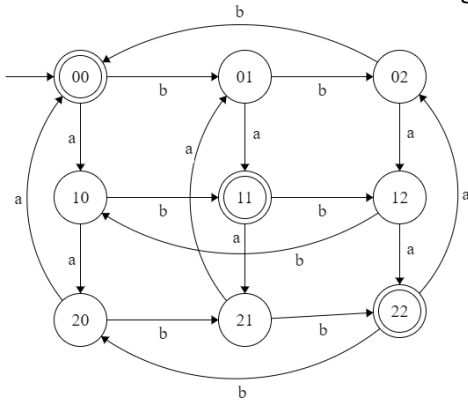
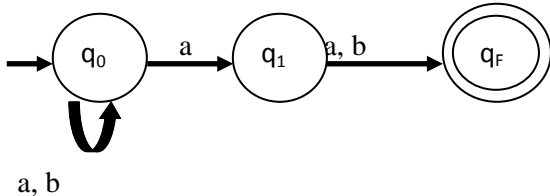


IAT 1 Question and Solutions

- 1 (a) Explain with example: i) Power of an alphabet ii)String reversal iii) Substring iv) Language [4]
- (b) If 'w' and 'x' are strings, Show that $(wx)^R = x^R w^R$ and give an example. [6]
- 2 (a) Write the formal definition of DFSM with a proper example. [4]
- (b) Design DFSM to accept the following languages. [6]
- i) $L = \{w \in \{a, b\}^* \mid \forall x, y \in \{a, b\}^*, (|w| \% 3 = 0)\}$
- ii) $L = \{w \in \Sigma^* \mid w \neq \epsilon \text{ and } e \text{ first and last character of } w \text{ are the same}\}$
- 3 (a) Design a NDFSM with $\Sigma = \{0, 1\}$ that accepts all string in which the third symbol from the right end is always 0. [4]
- (b) Let L1 and L2 be the two languages such that $L = \{001, 10, 111\}$ and $M = \{\epsilon, 001\}$, Find LUM, L.M, L-M, M-L, L^* when $L = \{0, 11\}$, L^* when $L = \{0, 1\}$. [6]
- 4 (a) Write a note on applications of theory of computation and compiler design. [4]
- (b) Obtain the transition table for the following DFSM: [6]



- 5 (a) Describe NDFSM using formal notation. Convert the following NDFSM to its equivalent DFSM (using subset construction method). [10]



- 6 (a) Obtain a DFSM for the following language- $L_1 = \{\text{starts with 'a' and ends with b}\}$ and $L_2 = \{\text{starts with 'b' and ends with 'a'}\}$. Then $L = L_1 \cup L_2$. [10]

CO 1	L1
CO 1	L1
CO 2	L3
CO 2	L3
CO 2	L3
CO 3	L2
CO 1	L1
CO 2	L3
CO2	L3
CO2	L3

Solution:

(a) Power of an alphabet :- denoted by

Σ^n where for some integers 'n',
the set of strings of length 'n' with
symbols from Σ .

$\Sigma^n = \{w \mid w \text{ is a string over } \Sigma \text{ and } |w|=n\}$

$\Sigma^0 =$ strings of length 0 = $\{\lambda\}$.

$\Sigma^1 = \{a, b\}$.

$\Sigma^2 = \{a, b\}, \{a, b\} = \{ab, ba, aa, bb\}$

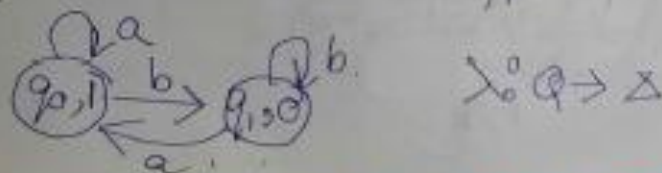
Finite State M/c with output

Moore Machine :-

$M = (Q, \Sigma, \delta, q_0, \Delta, \lambda)$

finite set of states input alphabet output alphabet output function

$q_0 \xrightarrow{a} q_0 \xrightarrow{b} q_1 \rightarrow \text{o/p}$



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	a	b
00	10	01
01	11	02
10	20	11
11	21	12
02	12	00
12	22	10
20	00	21
21	01	22
22	02	20

States = { ... }

$\Sigma = \{a, b\}$

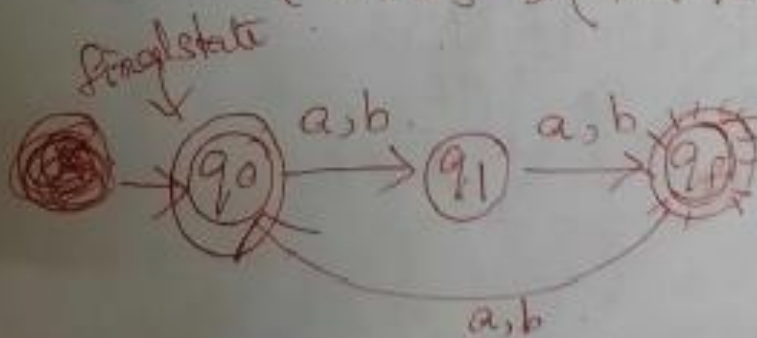
$F = \{00, 11, 22\}$

A: 01

B: 02

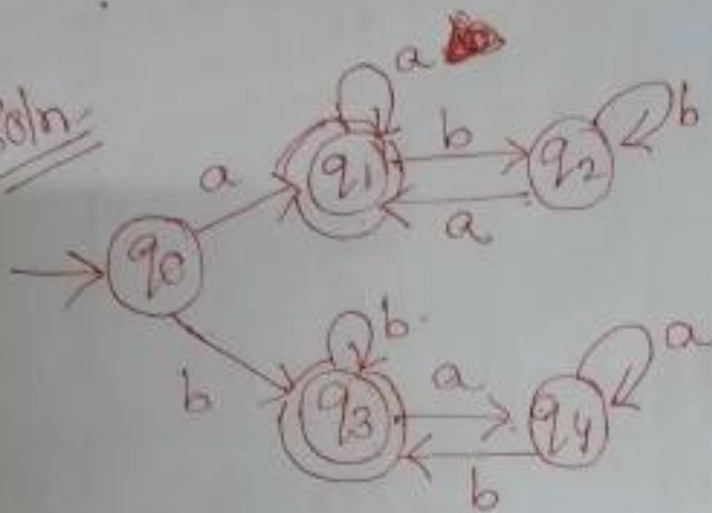
C: 00

(2b) (i) $L = \{w \in \{a, b\}^* \mid \forall x, y \in \{a, b\}^*, (|w| \% 3 = 0)\}$



2b(ii) $L = \{w \in \Sigma^* \mid w \neq \epsilon \text{ and the first and last char of } w \text{ are same}\}$

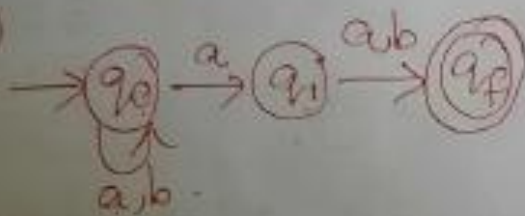
Soln



~~aaab~~
aabbaa

~~bbabab~~
bbabab
baab
bbbaaab

(5a)



(Final part of)

Soln

step 1

	a	b
q0	{q0, q1}	{q0}
q1	∅	∅
q2	∅	∅

NFA Transition table

Step 2 DFA Transition Table

	a	b	
$\rightarrow q_0$	$[q_0, q_1]$	$[q_0]$	$\delta([q_0, q_1], a)$
$[q_0, q_1]$	$[q_0, q_1, q_f]$	$[q_0, q_f]$	$= \delta(q_0, a) \cup \delta(q_1, a)$ $= [q_0, q_1] \cup [q_f]$
$[q_0, q_1, q_f]$	$[q_0, q_1, q_f]$	$[q_0, q_f]$	$= [q_0, q_1, q_f]$
$[q_0, q_f]$	$[q_0, q_1]$	$[q_0]$	$\delta([q_0, q_1], b)$ $= \delta(q_0, b) \cup \delta(q_1, b)$ $= [q_0] \cup [q_f]$ $= [q_0, q_f]$

$\delta([q_0, q_1, q_f], a)$

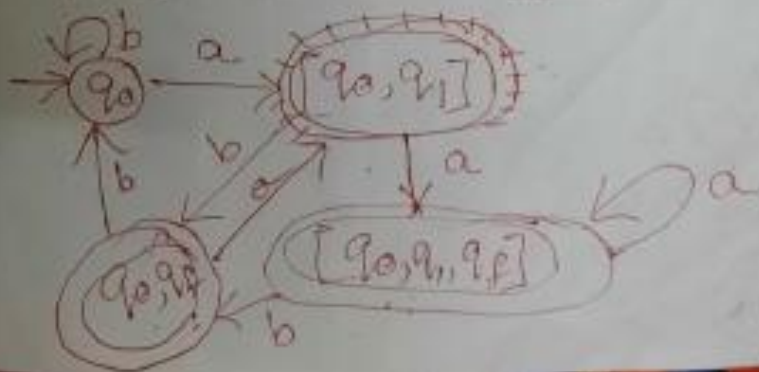
$= [q_0, q_1] \cup [q_f] \cup \emptyset$
 $= [q_0, q_1, q_f]$

$\delta([q_0, q_1, q_f], b)$
 $= [q_0] \cup [q_f] \cup \emptyset$
 $= [q_0, q_f]$

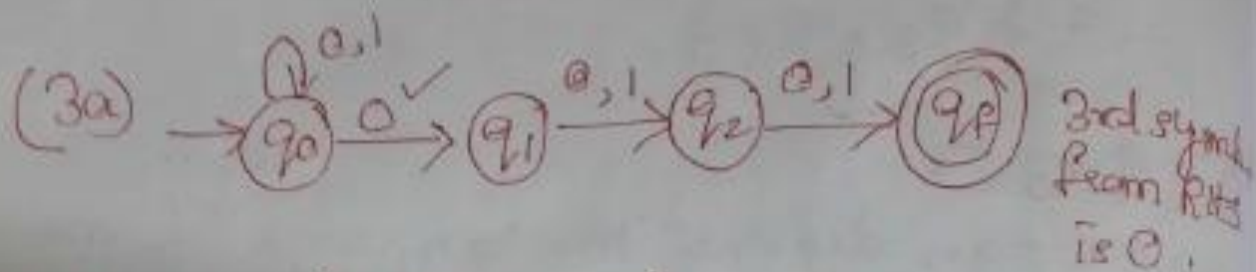
$\delta([q_0, q_f], a)$

$= [q_0, q_1] \cup \emptyset$
 $= [q_0, q_1]$

$\delta([q_0, q_f], b)$
 $= [q_0] \cup \emptyset$
 $= [q_0]$



IAT 1 Solution



(3b) $L = \{001, 10, 111\}$.

$H = \{\epsilon, 001\}$.

$L \cup H = \{\epsilon, 001, 10, 111\}$.

$L \cdot H = \{001, 10, 111, 001001, 10001, 111001\}$.

$L \cap H = \{001\}$.

$L - H = \{10, 111\}$.

$H - L = \{\epsilon\}$.

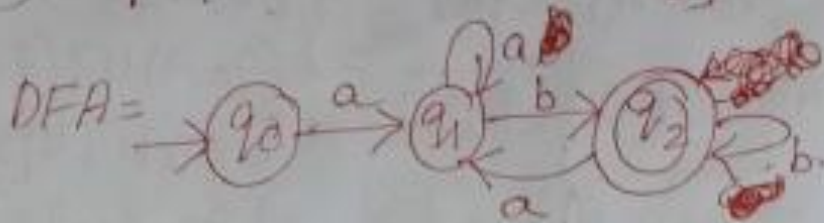
$L^* = 001, 11110, \epsilon$ (when $L = \{0, 1\}$)
 $= \{ \text{strings of 0 and 1 such that 1 come in pairs} \}$.

01011 X

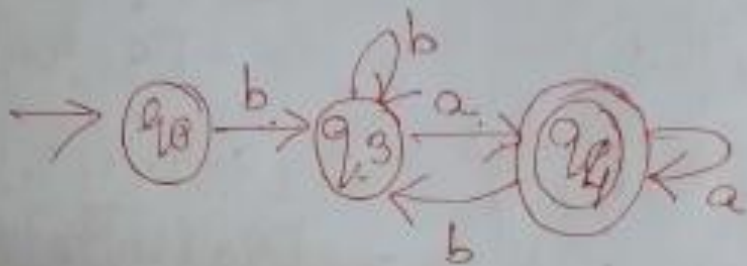
101 X

$L^* = \{ \text{all strings of 0's \& 1's} \}$.
 when $L = \{0, 1\}$.

(6a) $L_1 = \{ \text{starts with } a \text{ and ends with } b \}$.



$L_2 = \{ \text{starts with } b \text{ and ends with } a \}$.



$\therefore L = L_1 \cup L_2$ (Answer)

