

# CBCS SCHEME

22MCA11

USN 1 C R Q 3 m c l 0 1

## First Semester MCA Degree Examination, Dec.2023/Jan.2024 Mathematical Foundation for Computer Applications

Time: 3 hrs.

Max. Marks: 100

*Note: 1. Answer any FIVE full questions, choosing ONE full question from each module.  
2. M : Marks , L: Bloom's level , C: Course outcomes.*

		Module – 1	M	L	C
Q.1	a.	Define a set, an empty set and a singleton set, with an example for each.	6	L3	CO1
	b.	For any two sets A and B, prove that (i) $\overline{A \cap B} = \overline{A} \cup \overline{B}$ . (ii) $\overline{A \cup B} = \overline{A} \cap \overline{B}$ .	7	L3	CO1
	c.	Find the eigen values and eigen vectors of the matrix $\begin{pmatrix} 3 & 2 \\ -1 & 0 \end{pmatrix}$ .	7	L3	CO1
<b>OR</b>					
Q.2	a.	Let $U = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$ , $A = \{1, 2, 4, 6, 8\}$ , $B = \{2, 4, 5, 9\}$ . Compute the following: (i) $A \cup B$ (ii) $A \cap B$ (iii) $A - B$ (iv) $\overline{A}$	6	L3	CO1
	b.	A total of 1232 students have taken a course in Java, 879 in C and 114 in C++. Further, 103 have taken courses in both Java and C, 23 in both Java and C++ and 14 in both C and C++. If 2092 students have taken at least one of Java, C and C++, how many students have taken a course in all three subjects.	7	L3	CO1
	c.	State pigeon-hole principle. Show that if 50 books in a library contain a total of 27551 pages, one of the books must have atleast 552 pages.	7	L3	CO1
<b>Module – 2</b>					
Q.3	a.	Write the converse, inverse and contrapositive of, "If 2 is an even number then 7 is a prime number".	6	L2	CO3
	b.	Define tautology. Show that $[(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)] \rightarrow r$ is a tautology.	7	L2	CO3
	c.	If $A = \{1, 2, 3, 4, 5\}$ is the universe of discourse, determine the truth values of, (i) $\forall x \in A, (x + 2 < 10)$ (ii) $\exists x \in A, (x + 2 = 10)$ (iii) $\forall x \in A, (x^2 \leq 25)$	7	L2	CO3

		OR			
Q.4	a.	Show that for any three propositions p, q, r, $p \rightarrow (q \wedge r) \Leftrightarrow [(p \rightarrow q) \wedge (p \rightarrow r)]$ .	6	L2	CO3
	b.	Negate and simplify : (i) $\forall x, [p(x) \wedge \neg q(x)]$ . (ii) $\exists x, [(p(x) \vee q(x)) \rightarrow r(x)]$	7	L2	CO3
	c.	Show that the following argument is valid : "If today is Tuesday, I have a test in Mathematics or Economics. If my Economics professor is sick, I will not have a test in Economics. Today is Tuesday and my Economics professor is sick. Therefore, I have a test in Mathematics".	7	L2	CO3

## Module - 3

Q.5	a.	Let $A = \{1, 2, 3, 4, 6\}$ and R be a relation on A defined by 'aRb iff a is a multiple of b'. Write down R as a set of ordered pairs. Write matrix and digraph of R.	6	L3	CO1
	b.	Consider $A = \{1, 2, 3, \dots, 11, 12\}$ . The relation R on A is defined as $(x, y) \in R$ iff $x - y$ is a multiple of 5. Verify that R is an equivalence relation.	7	L3	CO1
	c.	Let $A = \{2, 3, 4, 6, 8, 12, 24\}$ and R denote the partial order of divisibility 'xRy iff x divides y'. Find R and draw its Hasse diagram.	7	L3	CO1

## OR

Q.6	a.	Let $A = \{1, 2, 3\}$ and $B = \{1, 2, 3, 4\}$ . Relations R and S from A to B are given by, $M_R = \begin{pmatrix} 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{pmatrix}, M_S = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 \end{pmatrix}$ Find (i) $R \cup S$ (ii) $R \cap S$ (iii) $\bar{S}$	6	L3	CO1
	b.	Let $A = \{1, 2, 3, 4\}$ & $R = \{(1, 1), (1, 2), (2, 3), (3, 4)\}$ , $S = \{(3, 1), (4, 4), (2, 4), (1, 4)\}$ be relations on A. Determine the relations $R \circ S$ , $S \circ R$ , $R^2$ , $S^2$ and draw their diagrams.	7	L3	CO1
	c.	Consider the POSET whose Hasse diagram Fig. Q6 (c) is as given below. Consider $B = \{3, 4, 5\}$ . Find (i) Maximal elements (ii) Minimal elements (iii) Greatest elements (iv) Least elements (v) Upper bound of B (vi) Lower Bounds of B (if they exist)	7	L2	CO4

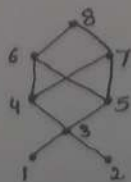


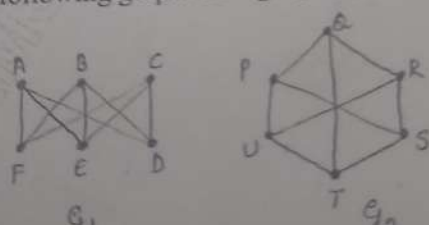
Fig. Q6 (c)

Module - 4				10	L3	CO2																		
Q.7	a.	Find the value of K such that the following distribution represents a finite probability distribution. Also find its mean and variance.	<table border="1"> <tr> <td>x</td> <td>0</td> <td>1</td> <td>2</td> <td>3</td> <td>4</td> <td>5</td> <td>6</td> <td>7</td> </tr> <tr> <td>P(x)</td> <td>0</td> <td>k</td> <td>2k</td> <td>2k</td> <td>3k</td> <td>k<sup>2</sup></td> <td>2k<sup>2</sup></td> <td>7k<sup>2</sup>+k</td> </tr> </table>	x	0	1	2	3	4	5	6	7	P(x)	0	k	2k	2k	3k	k <sup>2</sup>	2k <sup>2</sup>	7k <sup>2</sup> +k			
x	0	1	2	3	4	5	6	7																
P(x)	0	k	2k	2k	3k	k <sup>2</sup>	2k <sup>2</sup>	7k <sup>2</sup> +k																
	b.	When a coin is tossed 4 times, find the probability of getting, (i) exactly one head (ii) atmost 3 heads (iii) atleast two heads.		10	L3	CO2																		

OR

Q.8	a.	Find K such that $f(x) = \begin{cases} kx^2, & -3 \leq x \leq 3 \\ 0, & \text{elsewhere} \end{cases}$ is a probability density function. Also find (i) $P(1 \leq x \leq 2)$ (ii) $P(x > 1)$	6	L3	CO2
	b.	In a certain town, the duration of a shower is exponentially distributed with mean 5 minutes. What is the probability that the shower will last for, (i) 10 minutes or more (ii) less than 10 minutes (iii) between 10 and 12 minutes.	7	L3	CO2
	c.	The marks of 1000 students in an exam follows a normal distribution with mean 70 and standard deviation 5. Find the number of students whose marks will be, (i) less than 65 (ii) more than 75 (iii) 65 to 75. Given that $\phi(1) = 0.3413$ , where $z = \frac{x - \mu}{\sigma}$ .	7	L3	CO2

Module - 5

Q.9	a.	Define the following with an example for each : (i) Simple graph (ii) Regular graph (iii) Spanning subgraph	6	L2	CO4
	b.	Check whether the following graphs in Fig. Q9 (b) are isomorphic. 	7	L2	CO4
	c.	Explain Konigsberg Bridge Problem.	7	L2	CO4

OR

Q.10	a.	Define Euler circuit, Hamilton cycle and Planar graphs with an example for each.	6	L2	CO4
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OR

... three propositions p, q, r,

6 L2 CO3

b. Give graph coloring of the graph F.

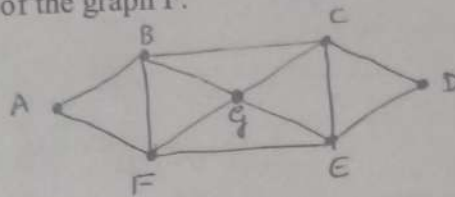


Fig. Q10 (b)

7 L2 CO4

c. Using the Dijkstra's algorithm. Obtain the shortest path from vertex 1 to each of the other vertices in the weighted, directed network shown below in Fig. Q10 (c)

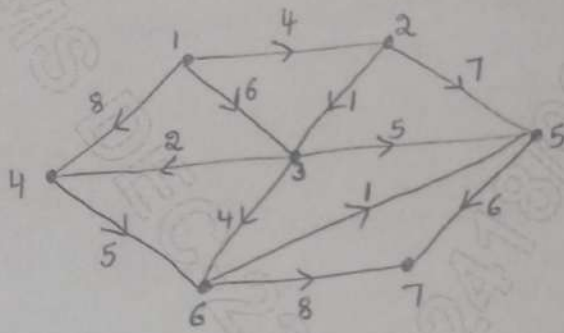


Fig. Q10 (c)

7 L2 CO4

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1a) ~~Q~~ A set is a collection of well defined objects.

$$\text{Eg: } \{2, 4, 6, 8, \dots\}$$

~~Q~~ Empty set: A set that contains no elements is called an empty set.

Eg: set of all prime numbers from 8 to 10.

singleton set: A set containing a single element is called a singleton set.

Eg: set of all multiples of 3 from 7 to 10.

1b. ~~Q~~  $\overline{A \cap B} = \{x / x \in \overline{A \cap B}\}$

$$= \{x / x \in \overline{A} \text{ and } x \in \overline{B}\}$$

$$= \{x / x \notin A \text{ and } x \notin B\}$$

$$= \{x / x \notin A \cup B\}$$

$$= \{x / x \in \overline{A \cup B}\} = \overline{A \cup B} \quad \therefore \boxed{\overline{A \cap B} = \overline{A \cup B}}$$

$$\overline{A \cup B} = \{x / x \in \overline{A} \text{ or } x \in \overline{B}\}$$

$$= \{x / x \notin A \text{ or } x \notin B\}$$

$$= \{x / x \notin A \cap B\}$$

$$= \{x / x \in \overline{A \cap B}\} \quad \therefore \boxed{\overline{A \cup B} = \overline{A \cap B}}$$

1c)  $A = \begin{pmatrix} 3 & 2 \\ -1 & 0 \end{pmatrix}$

characteristic eq<sup>n</sup> of A =  $|A - \lambda I| = 0$

$$\begin{vmatrix} 3-\lambda & 2 \\ -1 & -\lambda \end{vmatrix} = 0 \quad \Rightarrow \quad -\lambda(3-\lambda) + 2 = 0$$

$$\lambda^2 - 3\lambda + 2 = 0$$

$$\lambda^2 - 2\lambda - \lambda + 2 = 0$$

$$\lambda(\lambda - 2) - 1(\lambda - 2) = 0$$

$$(\lambda - 2)(\lambda - 1) = 0$$

$$\lambda = 1, 2$$

To find Eigen vectors consider

$$(A - \lambda I)X = 0$$

$$\begin{pmatrix} 3-\lambda & 2 \\ -1 & -\lambda \end{pmatrix} \begin{pmatrix} x \\ y \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

$$\Rightarrow \left. \begin{array}{l} (3-\lambda)x + 2y = 0 \\ \cancel{(3-\lambda)x + 2y = 0} \\ -x - \lambda y = 0 \end{array} \right\} \text{--- (1)}$$

Case (i) Put  $\lambda = 1$  in (1)

$$2x + 2y = 0$$

$$\Rightarrow x = -y$$

$$-x - y = 0$$

$(1, -1)$  is the

eigen vector

Case (ii) Put  $\lambda = 2$  in (1)

$$x + 2y = 0$$

$$\Rightarrow x = -2y$$

$$-x - 2y = 0$$

$$\frac{x}{-2} = \frac{y}{1}$$

$(-2, 1)$  is the eigen vector.

2(a)  $U = \{1, 2, 3, 4, 5, 6, 7, 8, 9\}$

$$A = \{1, 2, 4, 6, 8\}$$

$$B = \{2, 4, 5, 9\}$$

(i)  $A \cup B = \{1, 2, 4, 5, 6, 8, 9\}$

(ii)  $A \cap B = \{2, 4\}$

(iii)  $A - B = \{1, 6, 8\}$

(iv)  $\bar{A} = U - A = \{3, 5, 7, 9\}$

(b) Let  $A, B, C$  be the set of all students who have taken a course in

Java, C, C++ respectively.

Given  $|A| = 1232, |B| = 879, |C| = 114, |A \cap B| = 103, |A \cap C| = 23,$

$|B \cap C| = 14, |A \cup B \cup C| = 2092, |A \cap B \cap C| = ?$

From the principle of Inclusion-Exclusion,

$$|A \cup B \cup C| = |A| + |B| + |C| - |A \cap B| - |A \cap C| - |B \cap C| + |A \cap B \cap C|$$

$$2092 = 1232 + 879 + 114 - 103 - 23 - 14 + |A \cap B \cap C|$$

$$\Rightarrow |A \cap B \cap C| = 7$$

$\therefore$  7 students have taken all three courses.

(c) Let us treat 27551 pages as pigeons & 50 books as pigeon holes.

From Pigeon hole principle, at least one pigeon hole must contain

$p+1$  or more pigeons in it. Here,  $m = 27551$ ,  $n = 50$

$$p = \left\lfloor \frac{m-1}{n} \right\rfloor = \left\lfloor \frac{27551-1}{50} \right\rfloor = 551$$

$$p+1 = 552.$$

Hence, at least one book must contain 552 pages in it.

3(a) Let  $p$ : 2 is an even no

$q$ : 7 is a prime no

Given statement is  $p \rightarrow q$

Converse:  $q \rightarrow p$

ie., If 7 is a prime no then 2 is an even no.

Inverse:  $\neg p \rightarrow \neg q$

ie., If 2 is not an even no then 7 is not a prime no.

Contrapositive:  $\neg q \rightarrow \neg p$

ie., If 7 is not a prime no then 2 is not an even no.

3(b) If a compound statement is always true regardless of truth values of its components then it is called a tautology.

p	q	r	$p \vee q$	$p \rightarrow r$	$q \rightarrow r$	$(p \vee q) \wedge (p \rightarrow r) \wedge (q \rightarrow r)$	$x \rightarrow r$
0	0	0	0	1	1	0	1
0	0	1	0	1	1	0	1
0	1	0	1	1	0	0	1
0	1	1	1	1	1	1	1
1	0	0	1	0	1	0	1
1	0	1	1	1	1	0	1
1	1	0	1	0	0	0	1
1	1	1	1	1	1	1	1

Since all the entries of the last column are 1's, it is a tautology.

3(c)  $A = \{1, 2, 3, 4, 5\}$

(i)  $\forall x \in A, (x+2 < 10)$

$1+2 < 10, 2+2 < 10, 3+2 < 10, 4+2 < 10, 5+2 < 10$

Hence the statement is true.

(ii)  $\exists x \in A, (x+2 = 10)$

False  $x+2 \neq 10$  for any  $x \in A$ .

(iii)  $\forall x \in A, (x^2 \leq 25)$

$1^2 \leq 25, 2^2 \leq 25, 3^2 \leq 25, 4^2 \leq 25, 5^2 \leq 25$

True.

4(a) LHS =  $p \rightarrow (q \wedge r)$

$\Leftrightarrow \neg p \vee (q \wedge r)$

$\Leftrightarrow (\neg p \vee q) \wedge (\neg p \vee r)$

$\Leftrightarrow (p \rightarrow q) \wedge (p \rightarrow r)$

= RHS.

$p \rightarrow q \Leftrightarrow \neg p \vee q$



4(b) (i) Negation  $\neg [\forall x, \{p(x) \wedge \neg q(x)\}]$

$$\Leftrightarrow \exists x, \neg \{p(x) \wedge \neg q(x)\}$$

$$\Leftrightarrow \exists x, \{\neg p(x) \vee q(x)\}$$

(ii) Negn:  $\neg \{ \exists x, [\{p(x) \vee q(x)\} \rightarrow r(x)] \}$

$$\Leftrightarrow \forall x, \neg [\{p(x) \vee q(x)\} \rightarrow r(x)]$$

$$\Leftrightarrow \forall x, \{p(x) \vee q(x)\} \wedge \neg r(x)$$

4(c) Let  $p$ : Today is Tuesday.

$q$ : I have test in Maths

$r$ : I have a test in Economics.

$s$ : My Economics prof is sick

~~$t$ : I will have a test in Economics~~

Given  $p \rightarrow (q \vee r)$

$$s \rightarrow \neg r$$

$$p \wedge s$$

$$\therefore q$$

$\Rightarrow$

$$p \rightarrow (q \vee r)$$

$$s \rightarrow \neg r$$

$$p$$

$$s$$

$$\therefore q$$

$$\Rightarrow \frac{(q \vee r)}{\neg r}$$

$$\therefore q$$

Modus ponens

$\Rightarrow$

$$\neg q \rightarrow r$$

$$\neg r$$

$$\therefore q$$

$$\Rightarrow \frac{\neg(\neg q)}{\therefore q}$$

Modus Tollens

$\Rightarrow$

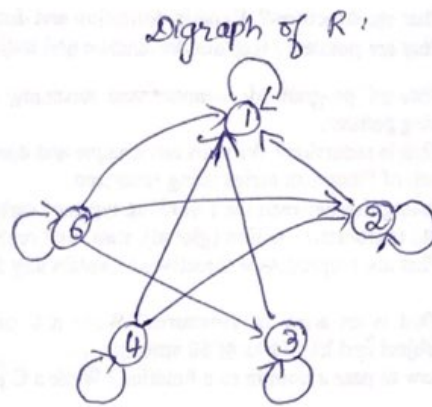
$$\therefore q$$

It is a valid argument.

5 (a)  $A = \{1, 2, 3, 4, 6\}$  & 'aRb iff a is a multiple of b'.

$R = \{(1,1), (2,1), (2,2), (4,1), (4,2), (4,4), (6,1), (6,2), (6,3), (6,6)\}$

$$M(R) = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 1 \end{bmatrix}$$



5(b)  $A = \{1, 2, 3, \dots, 11, 12\}$

$R = \{(1,1), (1,6), (1,11), (2,2), (2,7), (2,12), (3,3), (3,8), (4,4), (4,9), (5,5), (5,10), (6,1), (6,6), (6,11), (7,2), (7,7), (7,12), (8,3), (8,8), (9,4), (9,9), (10,5), (10,10), (11,1), (11,6), (11,11), (12,2), (12,7), (12,12)\}$

Reflexive: Wkt  $x-x=0$  is a multiple of 5,  $\forall x \in A$   
 $\Rightarrow (x,x) \in R$ ,  $\forall x \in A$   
 $\Rightarrow R$  is reflexive

Symmetric: Suppose  $(x,y) \in R$   
 $\Rightarrow x-y$  is a multiple of 5.  
 $\Rightarrow -(x-y)$  is a multiple of 5.  
 $\Rightarrow (y-x)$  is a multiple of 5.  
 $\Rightarrow (y,x) \in R$   
 $\therefore R$  is symmetric.

Anti-Symmetric: Suppose  $(x,y) \in R$  &  $(y,z) \in R$   
 $\Rightarrow x-y$  is a multiple of 5 &  $y-z$  is a multiple of 5

$$\Rightarrow x-y=5k_1 \quad \& \quad y-z=5k_2 \quad ; \quad k_1, k_2 \in \mathbb{Z}$$

$$\Rightarrow (x-y) + (y-z) = x-z = 5(k_1+k_2) \quad \text{which is a multiple of } 5.$$

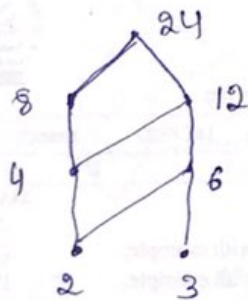
$$\Rightarrow (x, z) \in R$$

$\therefore R$  is transitive.

Hence  $R$  is an equivalence relation.

$$5(c) \quad A = \{2, 3, 4, 6, 8, 12, 24\}$$

$$R = \{(2,2), (2,4), (2,6), (2,8), (2,12), (2,24), (3,3), (3,6), (3,12), (3,24), (4,4), (4,8), (4,12), (4,24), (6,6), (6,12), (6,24), (8,8), (8,24), (12,12), (12,24), (24,24)\}$$



$$6(a) \quad R = \{(1,1), (1,3), (2,4), (3,1), (3,2), (3,3)\}$$

$$S = \{(1,1), (1,2), (1,3), (1,4), (2,4), (3,2), (3,4)\}$$

$$(i) \quad R \cup S = \{(1,1), (1,2), (1,3), (1,4), (2,4), (3,1), (3,2), (3,3), (3,4)\}$$

$$(ii) \quad R \cap S = \{(1,1), (1,3), (2,4), (3,2)\}$$

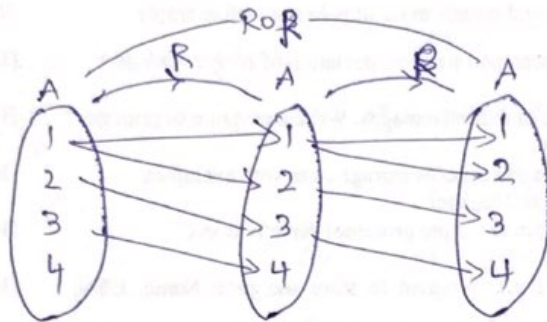
$$(iii) \quad \bar{S} = (A \times B) - S = \{(2,1), (2,2), (2,3), (3,1), (3,3)\}$$

6(b)  $A = \{1, 2, 3, 4\}$

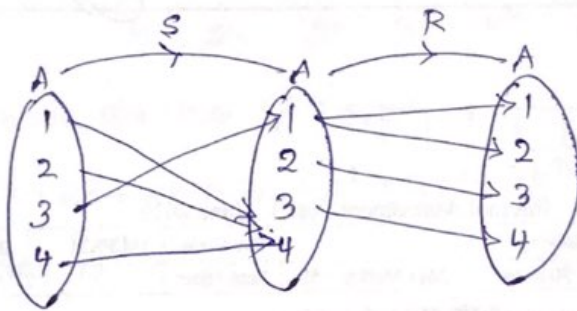
$R = \{(1,1), (1,2), (2,3), (3,4)\}$

$S = \{(3,1), (4,4), (2,4), (1,4)\}$

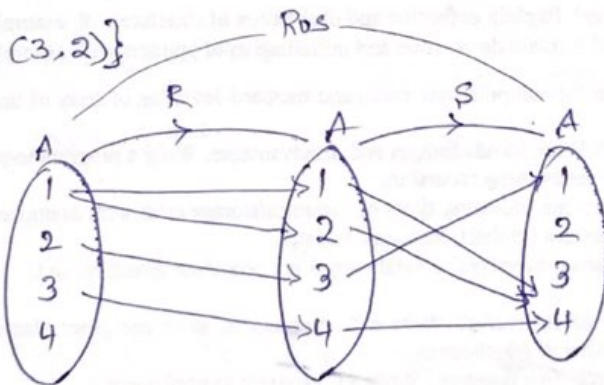
$R \circ R$



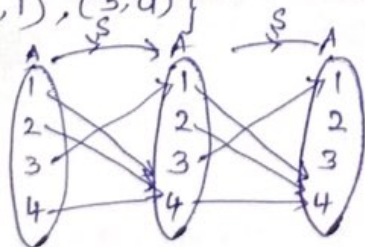
$R^2 = R \circ R = \{(1,1), (1,3), (1,2), (2,4)\}$



$S \circ R = \{(3,1), (3,2)\}$

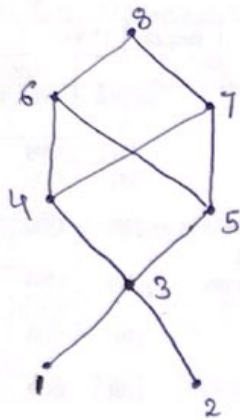


$R \circ S = \{(1,4), (2,1), (3,4)\}$



$S^2 = S \circ S = \{(1,4), (2,4), (3,4), (4,4)\}$

6(c)



$$B = \{3, 4, 5\}$$

(i) Maximal Elements =  $\{8\}$ (ii) Minimal Elements =  $\{1, 2\}$ (iii) Greatest element =  $\{8\}$ 

(iv) Least element = Doesn't exist

(v) Upper Bound of  $B = \{6, 7, 8\}$ (vi) Lower Bound of  $B = \{1, 2, 3\}$ 

7(a)

$x$	0	1	2	3	4	5	6	7
$p(x)$	0	$k$	$2k$	$2k$	$3k$	$k^2$	$2k^2$	$7k^2+k$

We must have  $p(x) \geq 0$  &  $\sum p(x) = 1$ 

$$\Rightarrow 0 + k + 2k + 2k + 3k + k^2 + 2k^2 + 7k^2 + k = 1$$

$$\Rightarrow 10k^2 + 9k - 1 = 0$$

$$\Rightarrow k = -1 \text{ \& } k = 1/10$$

$$\Rightarrow 10k^2 + 10k - k - 1 = 0$$

If  $k = -1$  then the condition

$$\Rightarrow 10k(k+1) - 1(k+1) = 0$$

 $p(x) \geq 0$  fails.

$$\Rightarrow (k+1)(10k-1) = 0$$

Hence  $k = 1/10$ 

$x$	0	1	2	3	4	5	6	7
$p(x)$	0	$1/10$	$1/5$	$1/5$	$3/10$	$1/100$	$1/50$	$17/100$

$$\begin{aligned} \text{Mean, } \mu &= \sum x p(x) = 0(0) + 1(1/10) + 2(1/5) + 3(1/5) + 4(3/10) + 5(1/100) \\ &\quad + 6(1/50) + 7(17/100) = \frac{183}{50} = 3.66 \end{aligned}$$

$$\begin{aligned} \text{Variance, } V &= \sum (x - \mu)^2 \cdot p(x) = (1 - 3.66)^2 \cdot 1/10 + (2 - 3.66)^2 \cdot 1/5 + \\ &\quad (3 - 3.66)^2 \cdot 1/5 + (4 - 3.66)^2 \cdot 3/10 + (5 - 3.66)^2 \cdot 1/100 + (6 - 3.66)^2 \cdot 1/50 + (7 - 3.66)^2 \cdot 17/100 \end{aligned}$$

$$7(b) \quad p = P(\text{head}) = \frac{1}{2} = 0.5$$

$$q = 1 - p = 0.5$$

$$n = 4$$

Wkt probability of  $x$  successes out of  $n$  trials is given by

$$P(x) = {}^n C_x p^x q^{n-x}$$

$$(i) \quad P(1 \text{ head}) = P(x=1) = {}^4 C_1 (0.5)^1 (0.5)^3 = 0.25$$

$$(ii) \quad P(\text{atmost } 3 \text{ heads}) = P(x=0) + P(x=1) + P(x=2) + P(x=3)$$

$$= {}^4 C_0 (0.5)^0 (0.5)^4 + {}^4 C_1 (0.5)^1 (0.5)^3 + {}^4 C_2 (0.5)^2 (0.5)^2 + {}^4 C_3 (0.5)^3 (0.5)^1$$

$$= 0.9375$$

$$(iii) \quad P(\text{atleast } 2 \text{ heads}) = 1 - [P(x=0) + P(x=1)]$$

$$= 1 - [(0.5)^4 + {}^4 C_1 (0.5)(0.5)^3]$$

$$= 0.6875$$

$$8(a) \quad f(x) = \begin{cases} kx^2, & -3 \leq x \leq 3 \\ 0, & \text{elsewhere} \end{cases}$$

$$\text{Wkt } f(x) \geq 0 \quad k > 0 \quad \& \quad \int_{-\infty}^{\infty} f(x) dx = 1$$

$$\Rightarrow \int_{-3}^3 kx^2 dx = 1 \quad \Rightarrow k \left[ \frac{x^3}{3} \right]_{-3}^3 = 1 \quad \Rightarrow \frac{k}{3} [3^3 - (-3)^3] = 1$$

$$\Rightarrow \frac{k}{3} [27 + 27] = 1 \quad \Rightarrow k = \frac{1}{18}$$

$$(i) \quad P(1 \leq x \leq 2) = \int_1^2 \frac{1}{18} x^2 \cdot dx = \frac{1}{18} \left[ \frac{x^3}{3} \right]_1^2 = \frac{7}{54}$$

$$(ii) \quad P(x > 1) = \int_1^3 \frac{1}{18} x^2 \cdot dx = \frac{1}{18} \left[ \frac{x^3}{3} \right]_1^3 = \frac{13}{27}$$

8(b) The p.d.f of exp distribution is given by

$$f(x) = \lambda e^{-\lambda x}, \quad x > 0$$

$$\text{Mean, } \mu = \frac{1}{\lambda} = 5 \quad \Rightarrow \quad \lambda = \frac{1}{5}$$

$$\therefore f(x) = \frac{1}{5} e^{-x/5}$$

$$(i) P(x \geq 10) = \int_{10}^{\infty} \frac{1}{5} e^{-x/5} dx = \frac{1}{5} \left[ \frac{e^{-x/5}}{-1/5} \right]_{10}^{\infty} = \frac{1}{e^2} - 0 = 0.1353$$

$$(ii) P(x < 10) = \int_0^{10} \frac{1}{5} e^{-x/5} dx = - \left[ e^{-x/5} \right]_0^{10} = 1 - e^{-2} = 0.8647$$

$$(iii) P(10 < x < 12) = \int_{10}^{12} \frac{1}{5} e^{-x/5} dx =$$

8(c) Here Mean,  $\mu = 70$ ,  $\sigma = 5$

$$Z = \frac{x - \mu}{\sigma} = \frac{x - 70}{5}$$

$$(i) \text{ If } x = 65, \quad z = -1$$

$$P(z < -1) = P(z > 1) = P(z \geq 0) - P(0 \leq z \leq 1)$$

$$= 0.5 - \phi(1) = 0.1587$$

$$(ii) \text{ If } x = 75, \quad z = 1$$

$$P(z > 1) = 0.1587$$

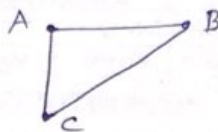
$$(iii) P(-1 < z < 1) = 2P(0 < z < 1) = 2\phi(1) = 2(0.3413) = 0.6826$$

$$\therefore \text{No. of students scoring marks less than 65} = 0.1587 \times 1000 \\ = 158.7 \approx 159$$

$$\text{a } \text{more than 75} = 0.1587 \times 1000 \\ \approx 159$$

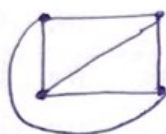
$$\text{a } \text{blw 65 to 75} = 0.6826 \times 1000 \approx 683$$

9(a) (i) Simple graph: A graph that has no loops or no multiple edges is called a simple graph.



(ii) Regular graph:

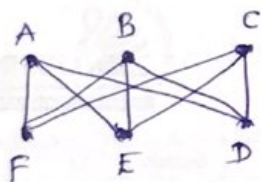
A graph in which the degree of every vertex is same, say  $k$ , is called a  $k$ -regular graph.



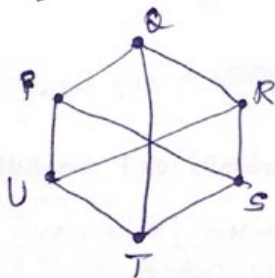
3-regular graph

(iii) A subgraph of  $G_1(V_1, E_1)$  of  $G_2(V, E)$  is said to be a spanning subgraph of  $G_2$  if  $V_1 = V$ .

9(b)



$G_1$



$G_2$

- Both  $G_1$  &  $G_2$  have 6 vertices.
- Both  $G_1$  &  $G_2$  have 9 edges.
- Both  $G_1$  &  $G_2$  are cubic graphs.

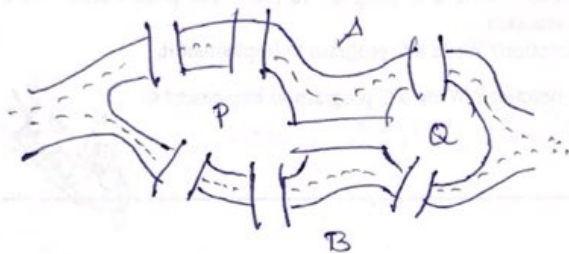
• Mapping b/w the vertices:  $A \leftrightarrow P, B \leftrightarrow T, C \leftrightarrow R, D \leftrightarrow S,$   
 $E \leftrightarrow Q, F \leftrightarrow U.$

• Mapping b/w the edges:  $\{A, F\} \leftrightarrow \{P, U\}, \{A, E\} \leftrightarrow \{P, Q\}, \{A, D\} \leftrightarrow \{P, S\}$   
 $\{B, F\} \leftrightarrow \{T, U\}, \{B, E\} \leftrightarrow \{T, Q\}, \{B, D\} \leftrightarrow \{T, S\}, \{C, F\} \leftrightarrow \{R, U\},$   
 $\{C, E\} \leftrightarrow \{R, Q\}, \{C, D\} \leftrightarrow \{R, S\}$

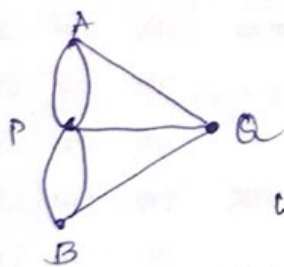
There exists an 1-1 correspondence b/w the vertices of  $G_1$  &  $G_2$ .  
Hence they are Isomorphic.



9(c) In 18<sup>th</sup> century, in a city named Königsberg, a river flowed named Pregel river which divides city into 4 parts - 2 banks of the river and 2 islands. These 4 land areas were connected to each other by 7 bridges. Citizens of the city posed a problem - "by starting at any of the four land areas, can we return to that area after crossing each bridge exactly once?". This problem is known as Königsberg bridge problem.



In 1736, Euler analyzed the problem with the help of a graph & gave the solution. Denote the land areas as A, B, P, Q as shown in the diagram. Construct a graph by treating the four land areas as four vertices & 7 bridges as edges connecting the vertices.



We note that in this graph

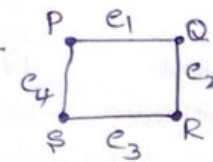
$$\deg(A) = \deg(B) = \deg(Q) = 3 \quad \& \quad \deg(P) = 5$$

which are not even.  $\therefore$  the graph doesn't have an

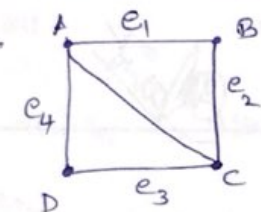
Euler circuit. Hence, it is not possible to walk over

each of the seven bridges exactly once & return to the starting point.


10 a) Euler Circuit - If there is a circuit in  $G$  that contains all the edges of  $G$ , then the circuit is called an Euler circuit in the graph  $G$ .

Eg: -  Here,  $Pe_1Qe_2Re_3Se_4P$  is an Euler circuit.

Hamilton Cycle: In a connected graph  $G$ , if there is a cycle that contains all the vertices of  $G$ , then that cycle is called a Hamilton cycle in  $G$ .

Eg: -  Here,  $Ae_1Be_2Ce_3De_4A$  is a Hamilton cycle.

Planar Graphs: A graph which can be represented as one plane drawing in which edges meet only at vertices is called a planar graph.

Eg: 

10 b) Let's assign the colour  $C_1$  to  $g$ . Then let's assign a colour  $C_2$  to  $B$  &  $E$  which are adjacent to  $g$ .  $C$  &  $F$  are adjacent to  $B$  &  $E$ . So, let's assign a colour  $C_3$  to  $C$  &  $F$ . Vertices  $A$  &  $D$  are not adjacent to  $g$ . So, we can assign the colour  $C_1$  to them. We have done the colouring with only 3 colours.  $\therefore$  chromatic number is 3.

10 c) For the first iteration, consider the set  $P = \{1\}$  &  $t_j = a_{1j}$ ,  
 $t_2 = a_{12} = 4$ ,  $t_3 = a_{13} = 6$ ,  $t_4 = 8$ ,  $t_5 = \infty$ ,  $t_6 = \infty$ ,  $t_7 = \infty$ .

Step 1:  $t_j$  is minimum for  $j=2$  i.e.,  $t_2=4$ . Label the arc  $(1,2)$  as  $P_1$ .

Also we adjoin 2 to P.  $\Rightarrow P = \{1,2\}$

Step 2:  $P = \{1,2\}$ ,  $t_2=4$ . Choose new  $t_3 = \min\{t_3, t_2 + q_{23}\} = \min\{6, 5\} = 5$

new  $t_4 = \min\{t_4, t_2 + q_{24}\} = \min\{8, 4 + \infty\} = 8$

new  $t_5 = 11$ , new  $t_6 = \infty$ , new  $t_7 = \infty$ .

Second Iteration:  $P = \{1,2\}$ ,  $t_2=4$ ,  $t_3=5$ ,  $t_4=8$ ,  $t_5=11$ ,  $t_6=\infty$ ,  $t_7=\infty$   
} new  $t_j$ 's

Step 1: Among new  $t_j$ 's,  $\min t_j = 5$  for  $j=3$ .

Among the vertices from P, arc  $(2,3)$  has least weight.  $(2,3)$  is  $P_2$ .

$P = \{1,2,3\}$

Step 2: new  $t_4=7$ , new  $t_5=10$ , new  $t_6=9$ , new  $t_7=\infty$

Third itr:  $P = \{1,2,3\}$ ,  $t_2=4$ ,  $t_3=5$  New  $t_4=7$ ,  $t_5=10$ ,  $t_6=9$ ,  $t_7=\infty$

Step 1: Among new  $t_j$ 's  $t_4=7$  is min for  $j=4$ . Label  $(3,4)$  as  $P_3$  that has least weight among the arcs from vertices of P to vertex 4.  $P = \{1,2,3,4\}$ .

Step 2: new  $t_5=10$ , new  $t_6=9$ , new  $t_7=\infty$

Fourth itr:

Step 1:  $\min t_j$  is for  $j=6$ . Label  $(3,6)$  as  $P_4$ .  $P = \{1,2,3,4,6\}$

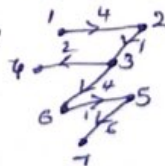
Step 2: new  $t_5=10$ ,  $t_7=\infty$

Fifth itr: Step 1: Among new  $t_j$ 's,  $t_5=10$  is min for  $j=5$ . Label  $(6,5)$  as  $P_5$ .

$P = \{1,2,3,4,6,5\}$

Final Step: Only one  $t_j$  is left.  $(5,7)$  has least weight. Label  $(5,7)$  as  $P_6$ .

$P = \{1,2,3,4,5,6,7\}$  which is the vertex set. Arcs labeled as  $P_1$  to  $P_6$  are  $(1,2)$ ,  $(2,3)$ ,  $(3,4)$ ,  $(3,6)$ ,  $(6,5)$ ,  $(5,7)$ . Shortest path algorithm is:



- (1) Path from 1 to 2:  $1 \rightarrow 2$ ; weight 4
- (2) — " — 1 to 3:  $1 \rightarrow 2 \rightarrow 3$ ; weight 5
- (3) — " — 1 to 5:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 6 \rightarrow 5$ ; weight 10
- (4) — " — 1 to 6:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 6$ ; weight 9
- (5) — " — 1 to 4:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 4$ ; weight 7
- (6) — " — 1 to 7:  $1 \rightarrow 2 \rightarrow 3 \rightarrow 6 \rightarrow 5 \rightarrow 7$ ; weight 16