USN					



				Interna	al Assessment	Test	1 - June 202	.4	1		. =			
Sı	ıb:	Database 3	Manageme	nt Systems			Sub Code:	BCS403	Bra	nch:	AIM AIM	IL/CS IL	E-	
Da	te:	04/06/24	06/24Duration:90 minutesMax Marks:50Sem/Sec:		IV		0		)BE					
				SCH	EME AND S	OLUI	<b>FIONS</b>			MAF	RKS	CO R		
1	a	file system? Answer:- Explanation Self-descri Insulation	of 4 points- bing nature o between prog multiple view data and mul	<b>4X2=8M</b> of a database s grams and data ws of the data tiuser transact	a, and data abstition processing	raction		from traditi	onal	[1(	D]	1	L1	
2	а	your own attr Answer: - Selection of 4	ibutes and re 4 entities rela h proper not	lationships ated to Hospi ration of entit	ta base conside tal-2M ies, attributes,	Ū				[6	j]	1	L3	
	b	Which one i Answer: Difference I Which one Achieving independence	s harder to a <b>between Lo</b> g <b>is harder to</b> logical dat ce. Changes	chieve? Why gical and Ph achieve? W a independe at the phys	ysical data ir	ndepe ler co not	ndence-2M ompared to need to be	physical reflected a	data t the		-]	1	L2	
3	a	Answer:- Data model I A data mode database—pro Different cat 1.High-level 2.Low-level o	Definition -2 el—a collecti ovides the ne egories or conceptus or physical d	M on of concep cessary means al data model lata models		used s abstr	to describe t action		of a	[6	5]	1	L2	

## Internal Assessment Test 1 – June 2024

b	<ul> <li>Explain the importance of Cardinality ratio and participation Constraints in a relational model.</li> <li>Answer-</li> <li>Cardinality Ratio- Definition +Example –2M</li> <li>The cardinality ratio for a binary relationship specifies the maximum number of relationship instances that an entity can participate in.</li> <li>Participation Constraint- Definition+Example—2M</li> <li>The participation constraint specifies whether the existence of an entity depends on its being related to another entity via the relationship type</li> <li>This constraint specifies the minimum number of relationship instances that each entity</li> </ul>	[4]	1	L2
	can participate in and is sometimes called the minimum cardinality constraint There are two types of participation constraints—total and partial			
4 a	Briefly describe the steps involved in ER- to relational mapping algorithm. Answer: Mention the steps-2M Each step +Explanation -8M Step 1: For each regular (strong) entity type E in the ER schema, create a relation R that includes all the simple attributes of E. Step 2: For each weak entity type W in the ER schema with owner entity type E, create a relation R, and includeall simple attributes (or simple components of composite attributes) of W as attributes. In addition, include as foreign key attributes of R the primary key attribute(s) of the relation(s) that correspond to the owner entity type(s). Step 3: For each binary 1:1 relationship type R in the ER schema, identify the relations, say S, and include the primary key of T as a foreign key in S. Include all the simple attributes of R as attributes of S. Step 4: For each regular binary 1:N relationship type R identify the relation (N) relation S. the primary key of T as a foreign key of S. Simple attributes of R map to attributes of S. Step 5: For each binary M:N relationship type R, create a relation S. Include the primary keys of participant relations as foreign keys in S. Their combination will be the primary keys for S. Simple attributes of S. Step 6: For each multi-valued attribute A, create a new relation R. This relation will include an attribute corresponding to A, plus the primary key K of the parent relation (entity type or relationship type) as a foreign key in R. The primary key of R is the combination of A and K. Step 7: For each n-ary relationship type R, where n>2, create a new relation S to represent R. Include the primary keys of the relations participating in R as foreign keys in S. Simple attributes of R map to attributes of S.	[10]	2	L2

		The primary key of S is a combination of all the foreign keys that reference the participants			
		that have cardinality constraint $> 1$ . For a recursive relationship, we will need a new relation.			
		Explain the different schema-based constraints.			
		Answer:-			
		List the constraints-1M			
	а	Explanation of each type -1X5=5M	[6]	2	L2
		The schema-based constraints include domain constraints, key constraints, constraints on			
		NULLs, entity integrity constraints, and referential integrity constraints.			
		Consider the relations.			
		EMPLOYEE(emp_id,name)			
		ASSIGNED_TO(projectno,emp_id)			
5		PROJECT(projectno,project_name)			
		Express the following queries in Relational Algebra.			
		i)Get details of employees working on both P354 and P345 project numbers.			
	b	ii)Find the employee number of employee who work on project P678	[4]	2	L3
		Answers: Each query 2M each			
		i) $\Pi$ emp_id,name( $\partial$ project_no.=P354 (EMPLOYEE*ASSIGNED_TO)) $\cap$ $\Pi$			
		emp_id,name (ôproject_no.=P345 (EMPLOYEE*ASSIGNED_TO))			
		ii) Π emp_id(∂project_no.=P678 (EMPLOYEE*ASSIGNED_TO))			
		List the various operations of relational algebra and explain the purpose of each with			
		examples.			
6	а	Answer:-	[10]	2	L2
		Each operator=1X10=10M with examples.			
L	1				

Operation	Purpose	Notation
SELECT	Selects all tuples that satisfy the selection condition from a relation $R$ .	$\sigma_{}(R)$
PROJECT	Produces a new relation with only some of the attributes of <i>R</i> , and removes duplicate tuples.	$\pi_{\text{}}(R)$
THETA JOIN	Produces all combinations of tuples from $R_1$ and $R_2$ that satisfy the join condition.	$R_1 \bowtie_{<\text{join condition}>} R_2$
EQUUOIN	Produces all the combinations of tuples from $R_1$ and $R_2$ that satisfy a join condition with only equality comparisons.	$R_1 \bowtie_{<\text{join condition>}} R_2,$ OR $R_1 \bowtie_{(<\text{join attributes 1>}),}$ ( <ioin 2="" attributes="">) <math>R_2</math></ioin>
NATURAL JOIN	Same as EQUUOIN except that the join attributes of $R_2$ are not included in the resulting relation; if the join attributes have the same names, they do not have to be specified at all.	$R_1 *_{} R_2,$ OR $R_1 *_{),$ $() R_2$ OR $R_1 * R_2$
UNION	Produces a relation that includes all the tuples in $R_1$ or $R_2$ or both $R_1$ and $R_2$ ; $R_1$ and $R_2$ must be union compatible.	$R_1 \cup R_2$
INTERSECTION	Produces a relation that includes all the tuples in both $R_1$ and $R_2$ ; $R_1$ and $R_2$ must be union compatible.	$R_1 \cap R_2$
SET DIFFERENCE	Produces a relation that includes all the tuples in $R_1$ that are not in $R_2$ ; $R_1$ and $R_2$ must be union compatible.	$R_1 - R_2$
CARTESIAN PRODUCT	Produces a relation that has the attributes of $R_1$ and $R_2$ and includes as tuples all possible combinations of tuples from $R_1$ and $R_2$ .	$R_1 \times R_2$
DIVISION	Produces a relation $R(X)$ that includes all tuples $t[X]$	$R_1(Z) + R_2(\gamma)$

CI

-----

CCI

All the Best-----

HoD