# SCHEME AND SOLUTION



## Internal Assessment Test I NOVEMBER 2024

Sub:	Artific	ial Intell	igence				Sub Code:	BCS515B	Branch	: ISE		
Date:	08-11	-2024	Duration:   90 min's   Max Marks:   50   Sem/Sec:   V- A,B,C				•	OB				
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- Challenge: Defining what is "rational" or "optimal" can be complex, especially in dynamic and uncertain environments.
- b. Explain different properties of task environment. Fully Observable vs. Partially Observable
- **Fully Observable**: In a fully observable environment, the agent has complete information about the state of the environment at each point in time. This means the agent can access all relevant data needed for decision-making without uncertainty.
- **Partially Observable**: In a partially observable environment, the agent has limited information, often due to noisy or incomplete sensor data. This requires the agent to make decisions based on partial knowledge, possibly predicting or inferring missing details.

**Example**: Chess is fully observable (all pieces and positions are visible), while driving a car in fog is partially observable (limited visibility affects the agent's perception).

### 2. Deterministic vs. Stochastic

- **Deterministic**: In a deterministic environment, any action taken by the agent has a predictable outcome with no randomness involved. The agent's actions always result in expected results, making planning straightforward.
- **Stochastic**: In a stochastic environment, actions can lead to different outcomes due to randomness or uncertainty, which makes the environment unpredictable. Here, the agent might need strategies to handle variability and incorporate probabilities.

**Example**: Solving a mathematical puzzle is deterministic, while dealing with traffic is stochastic due to the unpredictable behavior of other drivers.

#### 3. Episodic vs. Sequential

- **Episodic**: In an episodic environment, the agent's actions are divided into separate, independent episodes. Each episode does not depend on the previous one, so the agent doesn't need to consider the past when making decisions.
- Sequential: In a sequential environment, current actions affect future ones. The agent must consider the sequence of actions and the resulting state after each action.

**Example**: Image recognition tasks are episodic (each image can be classified independently), while chess is sequential since each move impacts future moves and the overall game outcome.

#### 4. Static vs. Dynamic

- **Static**: A static environment remains unchanged while the agent is deliberating. The agent doesn't need to worry about the environment changing between decision-making steps.
- Dynamic: In a dynamic environment, the environment can change while the agent is choosing or executing an action, requiring it to react quickly or adapt its strategy continuously.

**Example**: Solving a crossword puzzle is static (no changes occur as you think about the answers), while driving is dynamic, as road conditions and other vehicles change in real-time.

#### 5. Discrete vs. Continuous

- **Discrete**: In a discrete environment, there are a finite number of distinct states, actions, or time intervals. The agent's actions and decisions are taken in fixed steps or increments.
- Continuous: A continuous environment has a vast or infinite number of possible states or time intervals, requiring the agent to operate in a fluid, real-time manner.

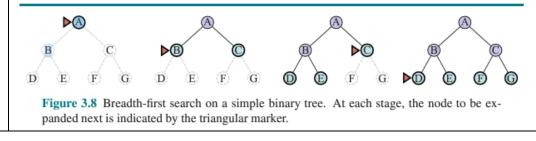
•	<ul> <li>Example: A turn-based game like chess is discrete, while controlling a robotic arm (which moves in smooth, continuous motions) is continuous.</li> <li>6. Single Agent vs. Multiagent</li> <li>Single Agent: In a single-agent environment, the agent is the only one acting within the environment, so it doesn't need to consider the influence of other agents' actions.</li> <li>Multiagent: In a multiagent environment, multiple agents interact, either as competitors or collaborators. Agents need to strategize based on other agents' actions, often involving cooperation, competition, or negotiation.</li> <li>Example: Playing a solo puzzle game is single-agent, while a multiplayer game like soccer is multiagent, as each player's actions affect others.</li> </ul>			
		5+5	CO1	L1
•	<ul> <li>all times. All relevant aspects of the environment's current state are available for the agent to perceive.</li> <li>Example: Chess, where the agent can see the entire board and all pieces.</li> <li>Partially Observable: The agent has limited information about the environment, often due to noisy sensors or restricted visibility. The agent must make decisions based on incomplete data.</li> <li>Example: Driving a car in fog, where visibility is reduced, and some information (like distant obstacles) is not immediately perceivable.</li> <li>ii) Single Agent vs. Multiagent</li> <li>Single Agent: The environment contains only one agent, so it doesn't need to consider the actions or decisions of other agents.</li> </ul>			

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	<ul> <li>Example: A robot vacuum cleaner in a home, where it is the only decision-making agent.</li> <li>Multiagent: The environment includes multiple agents, which may cooperate, environment in complex ways. The environment is complex ways.</li> </ul>			
	compete, or interact in complex ways. The agent must consider other agents'			
	actions and may need strategies to account for these interactions.			
	<ul> <li>Example: A game of soccer, where players (agents) must anticipate and respond to opponents and teammates.</li> </ul>			
	iii) Deterministic vs. Stochastic			
	• <b>Deterministic</b> : The outcome of any action is predictable and has no uncertainty.			
	Each action leads to a single, expected result, making it easier to plan.			
	<ul> <li>Example: Solving a mathematical problem, where each step has a known, predictable outcome.</li> </ul>			
	Stochastic: Actions lead to uncertain outcomes due to randomness or			
	unpredictability. The agent needs to handle variability and might use probabilities			
	to predict possible results.			
	• <b>Example</b> : Weather forecasting, where various factors introduce			
	uncertainty, leading to probabilistic predictions.			
	iv) Static vs. Dynamic			
	• Static: The environment remains unchanged while the agent is deliberating. It			
	doesn't evolve or alter independently of the agent's actions, making it simpler for planning.			
	• <b>Example</b> : A crossword puzzle, where the puzzle does not change as the agent thinks about solutions.			
	• Dynamic: The environment can change while the agent is in the process of			
	making or executing a decision, often requiring the agent to adapt in real-time.			
	$\circ$ Example: A self-driving car navigating through traffic, where the			
	environment (other vehicles, pedestrians) changes continuously.			
3	Provide a step-by-step illustration of how breadth-first search[BFS] works with example	10	<b>CO2</b>	L2
	and pseudo code			
	Breadth-First Search (BFS) is a search algorithm used for traversing or searching tree or			
	graph data structures. It explores all nodes at the present depth level before moving on to			
	nodes at the next depth level. BFS is typically implemented using a queue to keep track of			
	nodes to visit next.			
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When all actions have the same cost, an appropriate strategy is **breadth-first search**, in which the root node is expanded first, then all the successors of the root node are expanded next, then *their* successors, and so on. This is a systematic search strategy that is therefore complete even on infinite state spaces. We could implement breadth-first search as a call to BEST-FIRST-SEARCH where the evaluation function f(n) is the depth of the node—that is, the number of actions it takes to reach the node.

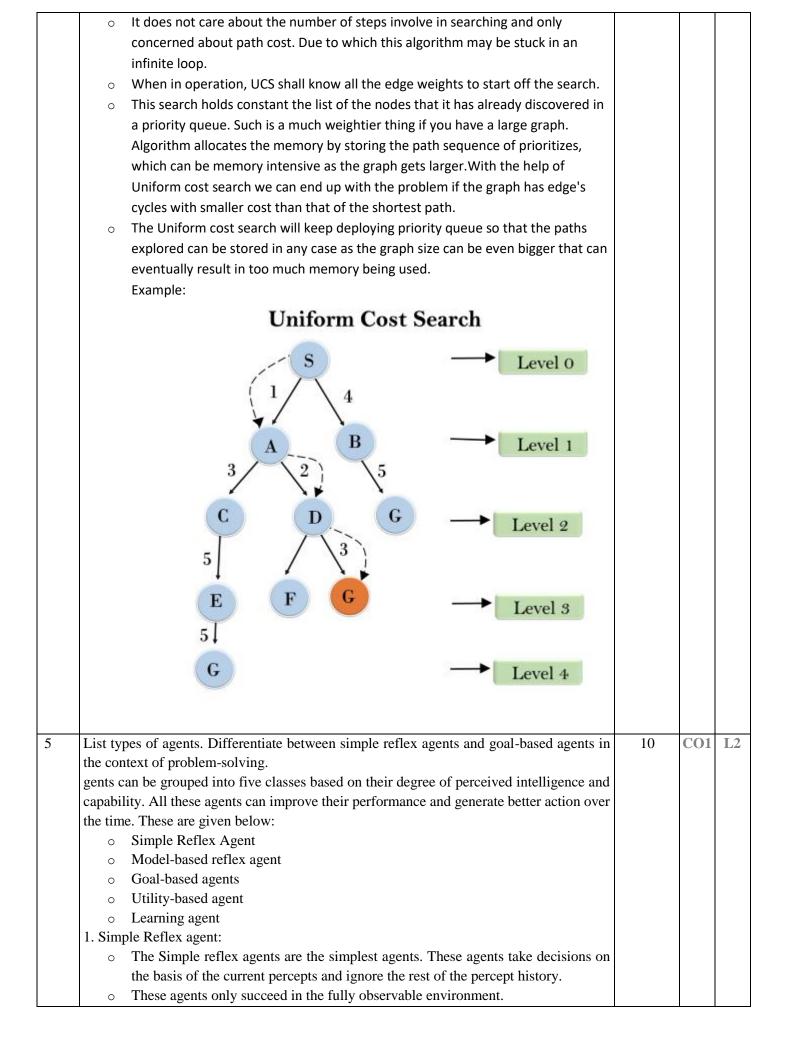
However, we can get additional efficiency with a couple of tricks. A first-in-first-out queue will be faster than a priority queue, and will give us the correct order of nodes: new nodes (which are always deeper than their parents) go to the back of the queue, and old nodes, which are shallower than the new nodes, get expanded first. In addition, *reached* can be a set of states rather than a mapping from states to nodes, because once we've reached a state, we can never find a better path to the state. That also means we can do an **early goal test**, checking whether a node is a solution as soon as it is *generated*, rather than the **late goal test** that best-first search uses, waiting until a node is popped off the queue. Figure 3.8 shows the progress of a breadth-first search on a binary tree, and Figure 3.9 shows the algorithm with the early-goal efficiency enhancements.

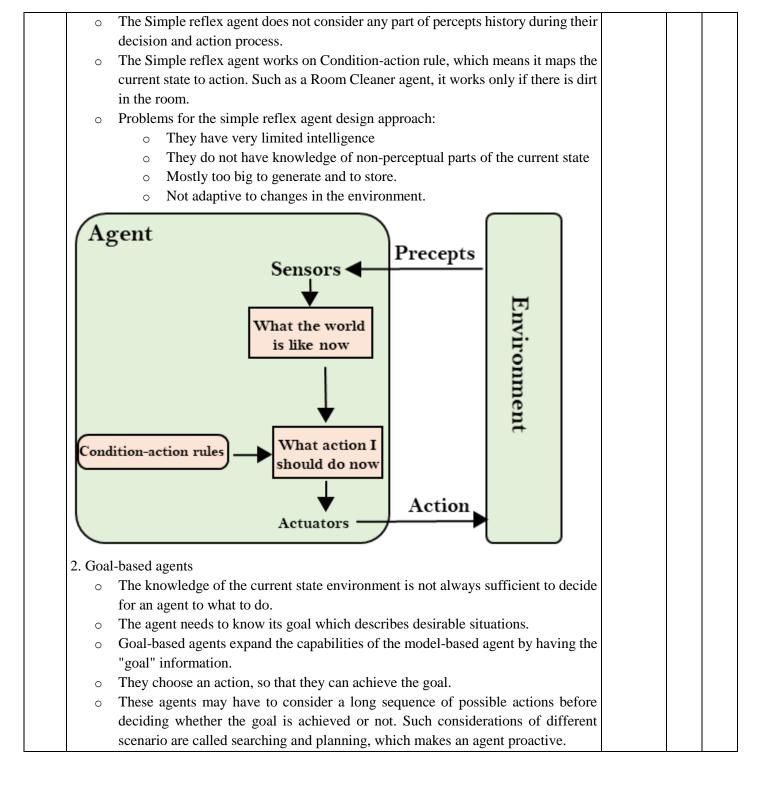
Breadth-first search always finds a solution with a minimal number of actions, because when it is generating nodes at depth d, it has already generated all the nodes at depth d - 1, so if one of them were a solution, it would have been found. That means it is cost-optimal



<ul> <li>protection intervent each intervent in terms of time and space, imagine searching a uniform tree where every state has b successors. The root of the search tree generates b nodes, each of which generates b more nodes, for a total of b<sup>2</sup> at the second level. Each of these generates b more nodes, yielding b<sup>3</sup> nodes at the third level, and so on. Now suppose that the solution is at depth d. Then the total number of nodes generated is         <ul> <li>1+b+b<sup>2</sup>+b<sup>3</sup>++b<sup>d</sup> = O(b<sup>d</sup>)</li> </ul> </li> <li>All the nodes remain in memory, so both time and space complexity are O(b<sup>d</sup>). Exponential bounds like that are scary. As a typical real-world example, consider a problem with branching factor b = 10, processing speed 1 million nodes/second, and memory requirements of 1 Kbyte/node. A search to depth d = 10 would take less than 3 hours, but would require 10 treabytes of memory. The memory requirements are a bigger problem for breadth-first search than the execution time. But time is still an important factor. At depth d = 14, even with infinite memory, the search would take 3.5 years. In general, exponential-complexity search problem. Solt. TEST(node. STATE) then return SolUTION(node) frontier - a FIFO queue with node as the only element explored - an empty set loog do</li></ul>	lems where all actions have the same cost, but not for problems that don't have that	
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<pre>node ← a node with STATE = problem.INITIAL-STATE, PATH-COST = 0 if problem.GOAL-TEST(node.STATE) then return SOLUTION(node) frontier ← a FIFO queue with node as the only element explored ← an empty set loop do     if EMPTY?(frontier) then return failure     node ← POP(frontier) /* chooses the shallowest node in frontier */     add node.STATE to explored     for each action in problem.ACTIONS(node.STATE) do         child ← CHILD-NODE(problem, node, action)     if child.STATE is not in explored or frontier then         if problem.GOAL-TEST(child.STATE) then return SOLUTION(child)         frontier ← INSERT(child, frontier) </pre>	s cannot be solved by uninformed search for any but the smallest instances.	
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Figure 3.11 Breadth-first search on a graph.		
	$frontier \leftarrow \text{INSERT}(child, frontier)$	
a. Explain the components and architecture of a problem-solving agent. 4+6 CO2	re 3.11 Breadth-first search on a graph.	
		)2

	<ul> <li>A problem-solving agent is a <i>goal-based agent</i> and use <i>atomic representations</i>.</li> <li>In atomic representations, states of the world are considered as wholes, with no internal structure vi the problem solving algorithms.</li> </ul>		
	<ul> <li>Intelligent agents are supposed to maximize their performance measure. Achieving this is som simplified if the agent can adopt a goal and aim at satisfying it.</li> </ul>		
	• Problem formulation is the process of deciding what actions and states to consider, given a g		
	• The process of looking for a sequence of actions that reaches the goal is called <b>search</b> .		
	• A <i>search algorithm</i> takes a problem as input and returns a <b>solution</b> in the form of an action sequence.		
	• Once a <i>solution</i> is found, the carrying actions it recommends is called the <b>execution phase</b> .		
	<ul> <li>A problem-solving agent has three phases:</li> <li>problem formulation, searching solution and executing actions in the solution.</li> </ul>		
	<ul> <li>A problem can be defined by five components:</li> <li>initial state, actions, transition model, goal test, path cost.</li> </ul>		
	<b>INITIAL STATE:</b> The <b>initial state</b> that the agent starts in.		
	ACTIONS: A description of the possible actions available to the agent.		
	• Given a particular state s, ACTIONS(s) returns the set of actions that can be executed in s.		
	• Each of these actions is <b>applicable</b> in s.		
	TRANSITION MODEL: A description of what each action does is known as the transition mod		
	• A function RESULT(s,a) that returns the state that results from doing action a in state s.		
	• The term <b>successor</b> to refer to any state reachable from a given state by a single		
	• The <b>state space</b> of the problem is the set of all states reachable from the <i>initial state</i> by any sequ of actions.		
	<ul> <li>The state space forms a graph in which the nodes are states and the links between nodes are act</li> </ul>		
1	• A <b>path</b> in the state space is a sequence of states connected by a sequence of actions.		
b.	Identify situations where uniform cost search performs better than other algorithms and explain with example.		
	Uniform-cost search is a searching algorithm used for traversing a weighted tree		
	or graph. This algorithm comes into play when a different cost is available for		
	each edge. The primary goal of the uniform-cost search is to find a path to the		
	goal node which has the lowest cumulative cost. Uniform-cost search expands		
	nodes according to their path costs form the root node. It can be used to solve		
	any graph/tree where the optimal cost is in demand. A uniform-cost search		
	algorithm is implemented by the priority queue. It gives maximum priority to		
	the lowest cumulative cost. Uniform cost search is equivalent to BFS algorithm if		
	the path cost of all edges is the same.		
	Advantages:		
0	Uniform cost search is optimal because at every state the path with the least cost is chosen.		
0	It is an efficient when the edge weights are small, as it explores the paths in an		
	order that ensures that the shortest path is found early.		
0	It's a fundamental search method that is not overly complex, making it accessible for many users.		
~	-		
0	It is a type of comprehensive algorithm that will find a solution if one exists. This		
	means the algorithm is complete, ensuring it can locate a solution whenever a		
	viable one is available. The algorithm covers all the necessary steps to arrive at a		
	resolution.		
	Disadvantages:		





	Sta How the world ev What my action Goa Agent	What the world is like now ms do What it will be like if I do action A	Precepts Fivironment			
6	<ul> <li>Evaluate the performance of BFS and DFS for different performance measures. The performance of tree search strategies is often evaluated based on several key criteria</li> <li>Completeness</li> <li>Optimality</li> <li>Time Complexity</li> <li>Space Complexity</li> <li>These criteria help determine the effectiveness of each search strategy based on the nature of the problem and the properties of the graph or tree being explored.</li> <li>Criteira BFS DFS</li> </ul>					L3
	Completeness	Completeness Complete for finite graphs Incomplete for infinite grap				
	Optimality	Optimal for unweighted graphs	Not optimal			
	Time Complexity	O(b^d)	O(b^m)			
	<ul> <li>Space Complexity</li> <li>b = branchin</li> <li>d = depth of</li> <li>m = maximut</li> </ul>	O(b * m)				