Internal Assessment Test 2 - May 2025 Scheme of Evlauation And Cloud Computing and Security Sub Code: BIS613D Branch: ISE atte: 24/05/2025 Duration: 90 min Max Marks: 50 Sem/Sec: VI/A, B & C OBE Answer any FIVE FULL Ouestions Answer any Full		1						CMR	METETUTE O	S TECHNOL	CMI	RI	
the: Cloud Computing and Security Sub Code: BIS613D Branch: ISE Answer any FIVE FULL Questions Answer any FIVE FULL Questions Scheme: Explanation of each design challenges of cloud. Scheme: Explanation of each design challenge 2 marks Solutions: Challenge 1—Service Availability and Data Lock-in Problem The management of a cloud service by a single company is often the source of single points of failure. To achieve HA, one can consider using multiple cloud providers. Even if a company has multiple data centers located in different geographic regions, it may have common software infrastructure and accounting systems. Therefore, using multiple cloud providers may provide more protection from failures. Another availability obstacle is distributed denial of service (DDS) attacks. Criminals threaten to cut off the incomes of Sans providers by making their services unavailable. Some utility computing services offer SaaS providers the opportunity to defend against DDoS attacks by using quick scale-ups. Challenge 2—Data Privacy and Security Concerns Current cloud offerings are essentially public (rather than private) networks, exposing the system to more attacks. Many obstacles can be overcome immediately with well-understood technologies such as encrypted storage, virtual LANs, and network middleboxes (e.g., firewalls, packet filters). For example, you could encrypt your data before placing it in a cloud. Many nations have laws requiring SaaS providers to keep customer data and copyrighted material within national boundaries. Traditional network attacks include buffer overflows, DoS attacks, spyware, malware, rooklis, Troian horses, and worms. In a cloud environment, newer attacks may result from hypervisor mal-Challenge 3—Unpredictable Performance and Bottlenecks Multiple VMs. can share CPUs and main memory in cloud compating, but I/O sharing is problematic. For example, to run 75 EC2 instances with the STREAM benchmark requires a mean bandwidth of 1,355 MB/second. However, for each of the 75 EC2 instances				Internal			·	5					
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1 Describe six architectural design challenges of cloud. Scheme: Explanation of each design challenge 2 marks Solutions: Challenge 1—Service Availability and Data Lock-in Problem The management of a cloud service by a single company is often the source of single points of failure. To achieve HA, one can consider using multiple cloud providers. Even if a company has multiple data centers located in different geographic regions, it may have common software infrastructure and accounting systems. Therefore, using multiple cloud providers may provide more protection from failures. Another availability obstacle is distributed denial of service (DDoS) attacks. Griminals threaten to cut off the incomes of SaaS providers by making their services unavailable. Some utility computing services offer SaaS providers the opportunity to defend against DDoS attacks by using quick scale-ups. Challenge 2—Data Privacy and Security Concerns Current cloud offerings are essentially public (rather than private) networks, exposing the system to more attacks. Many obstacles can be overcome immediately with well-understood technologies such as encrypted storage, virtual LANs, and network middleboxes (e.g., firewalls, packet filters). For example, you could encrypt your data before placing it in a cloud. Many nations have laws requiring SaaS providers to keep customer data and copyrighted material within national boundaries. Traditional network attacks include buffer overflows, DoS attacks, spyware, malware, rootkits, Trojan horses, and womes. In a cloud environment, newer attacks may result from hypervisor mal-Challenge 3—Unpredictable Performance and Bottlenecks Multiple VMs can share CPUs and main memory in cloud computing, but I/O sharing is problematic. For example, to run 75 EC2 instances with the STREAM benchmark requires a mean bandwidth of 1,355 MB/second. This demands the design of clinical distributed SANs. Data centers and the programmers' expectations in terms of scalability, data domainty, and HA. Data consistence checking	te:										ļ		
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Scheme: Explanation of each cloud model + Diagram – **8+2 marks** Solutions:

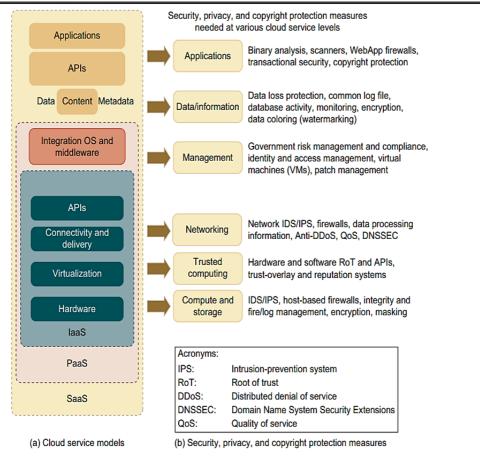


FIGURE 4.31

Cloud service models on the left (a) and corresponding security measures on the right (b); the laaS is at the innermost level, PaaS is at the middle level, and SaaS is at the outermost level, including all hardware, software, datasets, and networking resources.

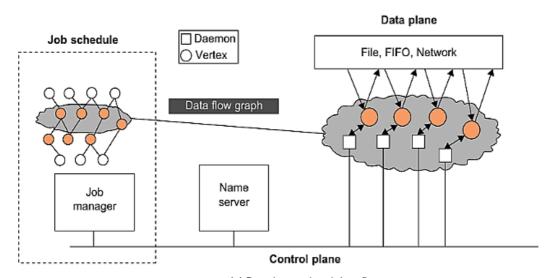
Three basic cloud security enforcements are expected. First, facility security in data centers demands on-site security year round. Biometric readers, CCTV (close-circuit TV), motion detection, and man traps are often deployed. Also, network security demands fault-tolerant external firewalls, intrusion detection systems (IDSes), and third-party vulnerability assessment. Finally, platform security demands SSL and data decryption, strict password policies, and system trust certification. Figure 4.31 shows the mapping of cloud models, where special security measures are deployed at various cloud operating levels.

Servers in the cloud can be physical machines or VMs. User interfaces are applied to request services. The provisioning tool carves out the systems from the cloud to satisfy the requested service. A security-aware cloud architecture demands security enforcement. Malware-based attacks such as network worms, viruses, and DDoS attacks exploit system vulnerabilities. These attacks compromise system functionality or provide intruders unauthorized access to critical information.

L2

Scheme: Definition+ Explanation + Components + Diagram – **2+3+3+2 marks** Solutions:

- **Dryad** is more flexible than MapReduce as the data flow of its applications is not dictated/predetermined and can be easily defined by users.
- To achieve such flexibility, a Dryad program or job is defined by a directed acyclic graph (DAG) where vertices are computation engines and edges are communication channels between vertices.
- •Given a DAG, Dryad assigns the computational vertices to the underlying computation engines (cluster nodes) and controls the data flow through edges (communication between cluster nodes).
- Data partitioning, scheduling, mapping, synchronization, communication, and fault tolerance are major implementation details hidden by Dryad to facilitate its programming environment.



(a) Dryad control and data flow

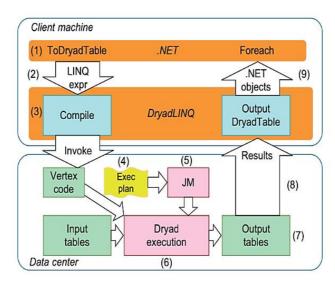


FIGURE 6.14

LINQ-expression execution in DryadLINQ.

DryadLINQ is built on top of Microsoft's Dryad execution framework (see http://research.microsoft .com/en-us/projects/DryadLINQ/). Dryad can perform acyclic task scheduling and run on large-scale servers. The goal of DryadLINQ is to make large-scale, distributed cluster computing available to ordinary programmers. Actually, DryadLINQ, as the name implies, combines two important components: the Dryad distributed execution engine and .NET Language Integrated Query (LINQ). LINQ is particularly for users familiar with a database programming model. Figure 6.14 shows the flow of execution with DryadLINQ. The execution is divided into nine steps as follows:

- 1. A .NET user application runs, and creates a DryadLINQ expression object. Because of LINQ's deferred evaluation, the actual execution of the expression has not occurred.
- 2. The application calls ToDryadTable triggering a data-parallel execution. The expression object is handed to DryadLINQ.
- 3. DryadLINQ compiles the LINQ expression into a distributed Dryad execution plan. The expression is decomposed into subexpressions, each to be run in a separate Dryad vertex. Code and static data for the remote Dryad vertices are generated, followed by the serialization code for the required data types.
- 4. DryadLINQ invokes a custom Dryad job manager which is used to manage and monitor the execution flow of the corresponding task.
- 5. The job manager creates the job graph using the plan created in step 3. It schedules and spawns the vertices as resources become available.

4 a. Explain the interactions among VM managers for cloud creation and management. [05]

Scheme: Explanation + Diagram – 3+2 marks Solutions:

3 L2

ndividual users Other clouds Cloud consumers Need raw Need to outsource Need resources on which to instantiate infrastructure excess workloads services (web, databases, and so on) for their users (a) Cloud interfaces (Amazon EC2WS Nimbus WSRF, ElasticHosts REST) Cloud toolkits currently do not use virtual infrastructure managers and, instead, manage VMs themselves Cloud Eucalyptus directly, without providing the full set management of features of VI managers public clouds (b) **OpenNebula** VI management and others (c) **KVM VMware** VM managers

FIGURE 4.4

(d)

Cloud ecosystem for building private clouds: (a) Consumers demand a flexible platform; (b) Cloud manager provides virtualized resources over an laaS platform; (c) VI manager allocates VMs; (d) VM managers handle

An ecosystem was suggested by Sotomayor, et al. [39] (Figure 4.4) for building private clouds. They suggested four levels of ecosystem development in a private cloud. At the user end, consumers demand a flexible platform. At the cloud management level, the cloud manager provides virtualized resources over an IaaS platform. At the virtual infrastructure (VI) management level, the manager allocates VMs over multiple server clusters. Finally, at the VM management level, the VM managers handle VMs installed on individual host machines. An ecosystem of cloud tools attempts to span both cloud management and VI management. Integrating these two layers is complicated by the lack of open and standard interfaces between them.

An increasing number of startup companies are now basing their IT strategies on cloud resources, spending little or no capital to manage their own IT infrastructures. We desire a flexible and open architecture that enables organizations to build private/hybrid clouds. VI management is aimed at this goal. Example VI tools include oVirt (https://fedorahosted.org/ovirt/), vSphere/4

b. Illustrate the infrastructure needed to virtualized servers in data centers.

Scheme: Explanation + Diagram – 3+2 marks

Solutions:

[05]

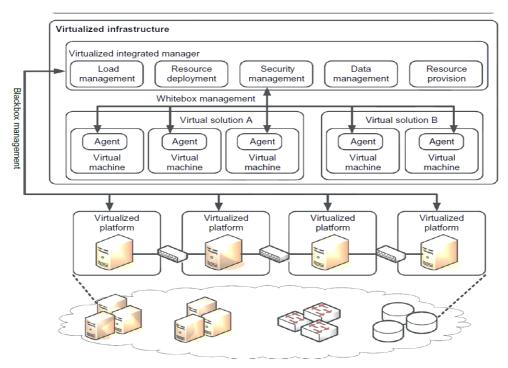


FIGURE 4.17

Virtualized servers, storage, and network for cloud platform construction.

- One very distinguishing feature of cloud computing infrastructure is the use of system virtualization and the modification to provisioning tools.
- Virtualization of servers on a shared cluster can consolidate web services. As the VMs
 are the containers of cloud services, the provisioning tools will first find the
 corresponding physical machines and deploy the VMs to those nodes before scheduling
 the service to run on the virtual nodes.
- In addition, in cloud computing, virtualization also means the resources and fundamental infrastructure are virtualized.
- The user will not care about the computing resources that are used for providing the services.
- Cloud users do not need to know and have no way to discover physical resources that are involved while processing a service request.
- Also, application developers do not care about some infrastructure issues such as scalability and fault tolerance (i.e., they are virtualized). Application developers focus on service logic.

L2

[10]

5

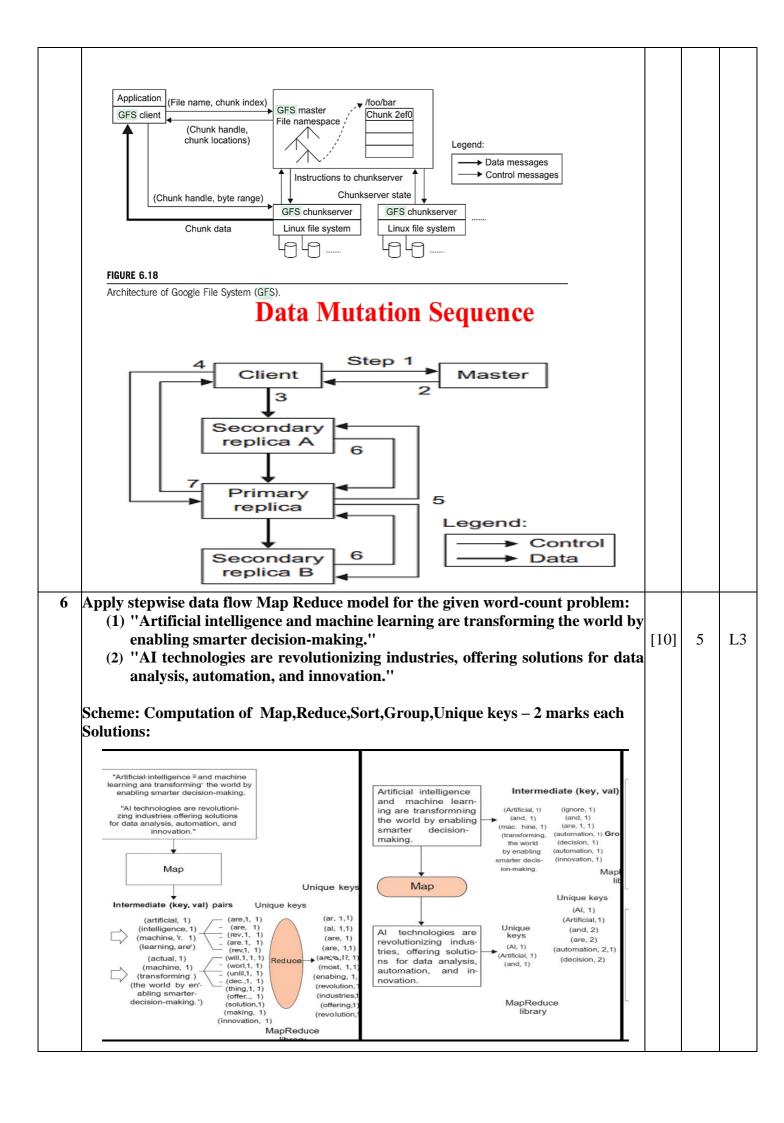
5 Describe architecture and data mutation sequences in GFS.

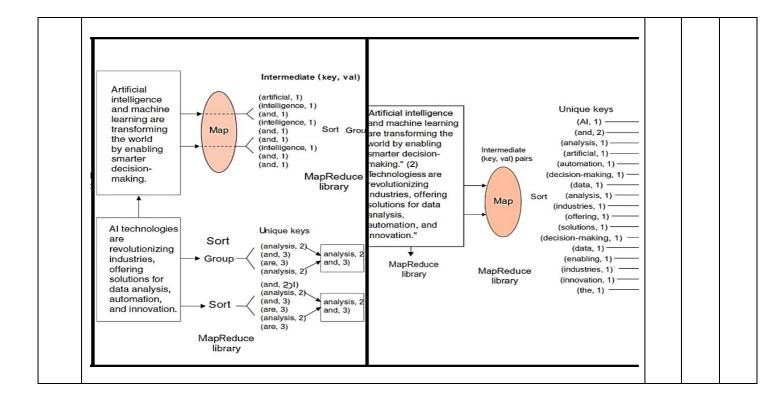
Scheme: Architecture (Explanation + Diagram) -3+2 marks and Mutation Sequence (Explanation + Diagram) -3+2 marks

Solutions:

GFS was built primarily as the fundamental storage service for Google's search engine. As the size of the web data that was crawled and saved was quite substantial, Google needed a distributed file system to redundantly store massive amounts of data on cheap and unreliable computers. None of the traditional distributed file systems can provide such functions and hold such large amounts of data. In addition, GFS was designed for Google applications, and Google applications were built for GFS. In traditional file system design, such a philosophy is not attractive, as there should be a clear interface between applications and the file system, such as a POSIX interface.

There are several assumptions concerning GFS. One is related to the characteristic of the cloud computing hardware infrastructure (i.e., the high component failure rate). As servers are composed of inexpensive commodity components, it is the norm rather than the exception that concurrent failures will occur all the time. Another concerns the file size in GFS. GFS typically will hold a large number of huge files, each 100 MB or larger, with files that are multiple GB in size quite common. Thus, Google has chosen its file data block size to be 64 MB instead of the 4 KB in typical traditional file systems. The I/O pattern in the Google application is also special. Files are typically written once, and the write operations are often the appending data blocks to the end of files. Multiple appending operations might be concurrent. There will be a lot of large streaming reads and only a little random access. As for large streaming reads, highly sustained throughput is much more important than low latency.





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