

Internal Assessment Test 1 –November-2025

Sub:	DATA STRUCTURES AND APPLICATIONS					Sub Code:	BCS304	Branch:	AIML & CSML	
Date:		Duration:	90 minutes	Max Marks:	50	Sem/Sec:	III -A, B, C		OBE	
Answer any FIVE FULL Questions								MARKS	CO	RBT
1	a	Define Data Structures. Explain the classification of data structures with a neat diagram. DATA STRUCTURE 1M Data structure is a representation of the logical relationships existing between individual elements of data. A data structure is a way of organizing all data items that considers not only the elements stored but also their relationship to each other. The logical or mathematical model of a particular organization of data is called a data structure. The non-primitive data structures are further classified into 1M(diagram) Data Structure <pre>graph TD DS[Data Structure] --> P[Primitive] DS --> NP[Non-Primitive] P --> Int[Int] P --> Char[Char] P --> Bool[Bool] P --> Float[Float] P --> Pointer[Pointer] NP --> L[Linear] NP --> NL[Non-Linear] L --> S[Static] L --> D[Dynamic] S --> A[Array] D --> LL[Linked List] D --> ST[Stack] D --> Q[Queue] NL --> T[Tree] NL --> G[Graph]</pre> 1. Linear Data Structure 2M 2. Non-linear Data Structure 2M 1. Linear Data Structure: A data structure is said to be linear if its elements form a sequence or a linear list. There are basically two ways of representing such linear structure in memory. a. One way is to have the linear relationships between the elements represented by means of sequential memory location. These linear structures are called arrays. b. The other way is to have the linear relationship between the elements represented by means of pointers or links. These linear structures are called linked lists. The common examples of linear data structure are Arrays, Queues, Stacks, Linked lists 2. Non-linear Data Structure: A data structure is said to be non-linear if the data are not arranged in sequence or a linear. The insertion and deletion of data is not possible in linear fashion. This structure is mainly used to represent data containing a hierarchical relationship between elements. Trees and graphs are the examples of non-linear data structure.						6	1	L2
	b	Differentiate structures and unions. Solution 3M						4	1	L2

	Structure	Union
Definition	uses dif. data type which groups diff. data types into a single entity	at the same memory locat ⁿ
Keyword	struct	union
Size	Sum of the Size of all members used	Size of the largest member
Memory locat ⁿ	each member within a structure is allocated unique storage area of locat ⁿ	memory locat ⁿ is shared by individual members of union.

Date:/...../.....

Data overlap	No data overlap as members are independent	full data overlap as members share the same memory.
--------------	--	---

Accessing Mem	Individual members can be accessed at a time	Only one member can be accessed at a time
---------------	--	---

Example Declaration 1M

structure

```
struct student { int id; char name[20]; float marks; };
```

union

```
union data { int i; float f; char ch; };
```

Explain various memory allocation techniques? Explain how memory can be dynamically allocated using malloc().

Solution

various memory allocation techniques

Static Memory Allocation 1M

Static Memory is allocated for declared variables by the compiler. The address can be found using the address of operator and can be assigned to a pointer. The memory is allocated during

2

a

4

2

L2

	compile time. Memory allocation 1M Memory allocation done at the time of execution(run time) is known as dynamic memory allocation. Functions calloc() and malloc() support allocating dynamic memory. In the Dynamic allocation of memory space is allocated by using these functions when the value is returned by functions and assigned to pointer variables. Dynamic Memory Allocation functions in C: 2M malloc() Related header file is stdlib.h> These functions provide a complete set of memory allocation, reallocation, deallocation, and heap management tools. malloc() Syntax: ptr = (cast-type*) malloc(byte-size); Here, ptr is pointer of cast-type. The malloc() function returns a pointer to an area of memory with size of byte size. If the space is insufficient, allocation fails and returns NULL pointer. Example: ptr = (int*) malloc(100 * sizeof(int));			
b	Write a C Functions to implement pop, push and display operations for stacks using arrays. 2(each function)*3=6 Solution:- 2(each function)*3=6 <pre>int stack[MAX]; int top = -1; // Function to push an element into the stack void push(int value) { if (top == MAX - 1) { printf("Stack Overflow! Cannot push %d\n", value); } else { top++; stack[top] = value; printf("%d pushed into the stack.\n", value); } } // Function to pop an element from the stack void pop() { if (top == -1) { printf("Stack Underflow! No elements to pop.\n"); } else { printf("%d popped from the stack.\n", stack[top]); top--; } } // Function to display all elements in the stack void display() { if (top == -1) { printf("Stack is empty.\n"); } else { printf("Stack elements are:\n"); for (int i = top; i >= 0; i--) { printf("%d\n", stack[i]); } } }</pre>	6	1	L3
3	Evaluate the following postfix expressions, 1. abc+*de/- where a=5, b=6, c=2, d=12, e=4 2. AB+CDE-+/ where A=5, B=5, C=4, D=7, E=1	10	2	L3

Answer 2: 1 5M

5

1) $abc + *de/ -$
 $562 + *1241 -$

$a=5, b=6, c=2, d=12, e=4$

	Stack
5	5
6	5 6
2	5 6 2
+	5 8 3
*	40
12	40 12 4
/	40 3
-	37

2) $AB + CDE - + /$
 $55 + 471 - + /$

	Stack
5	5
5	5 5
+	10
4	10 4
7	10 4 7
1	10 4 7 1
-	10 4 6
+	10 10
/	1

i) Insert a node at the beginning ii) Delete a node at the front iii) Display

singly linked list explanation **2.5M**

Diagram illustrating a linked list structure. A pointer labeled "NAME or START" points to the first node. The list consists of six nodes, each containing an "Information part" and a "Nextpointer field". The last node's nextpointer field points to a null value (X). Labels indicate the "Nextpointer field of third node" and "Information part of third node".

Fig: Linked list with 6 nodes

```
START=9      INFO[9]=N
LINK[3]=6    INFO[6]=V
LINK[6]=11   INFO[11]=E
LINK[11]=7   INFO[7]=X
LINK[7]=10   INFO[10]=I
LINK[10]=4   INFO[4]=T
LINK[4]= NULL value, So the list has ended
```

Defining a node structure

```
typedef struct listNode *listPointer typedef struct {  
char data[4]; listPointer list;  
} listNode;
```

Create a New Empty list

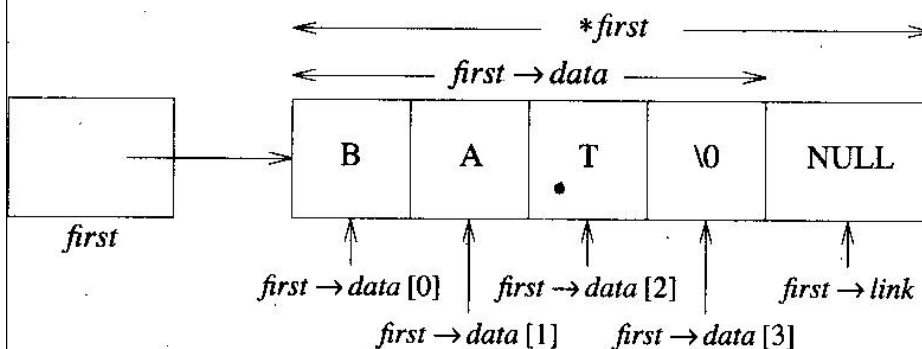
```
listPointer first = NULL
```

To create a New Node

```
MALLOC (first, sizeof(*first));
```

To place the data into NODE

```
strcpy(first-> data,"BAT"); first-> link = NULL
```



Function to insert a node at the beginning-2.5M

```
void insertAtBeginning(int value) {  
    struct Node* newNode = (struct Node*)malloc(sizeof(struct Node));  
    newNode->data = value;  
    newNode->next = head; // Point new node to current head  
    head = newNode;      // Move head to new node  
    printf("%d inserted at the beginning.\n", value);  
}
```

Function to delete a node from the front-2.5M

```
void deleteFromFront() {  
    if (head == NULL) {  
        printf("List is empty. Nothing to delete.\n");  
    } else {  
        struct Node* temp = head;  
        head = head->next; // Move head to next node  
        printf("%d deleted from the front.\n", temp->data);  
        free(temp); // Free old head  
    }
```

Function to display the list- 2.5M

```
void display() {  
    if (head == NULL) {  
        printf("List is empty.\n");  
    } else {  
        struct Node* temp = head;  
        printf("Linked List: ");  
        while (temp != NULL) {  
            printf("%d -> ", temp->data);  
            temp = temp->next;  
        }  
    }
```


	<pre> printf("NULL\n"); } } </pre>			
5	What is Circular queue? How insertion () , deletion () and display ()done on circular queue. <u>CIRCULAR QUEUES- 2.5</u> It is “The queue which wrap around the end of the array.” The array positions are arranged in a circle as shown in figure. In this convention the variable front is changed. front variable points one position counterclockwise from the location of the front element in the queue. The convention for rear is unchanged. The diagrams illustrate the operations on a circular queue with 8 slots. In (a) Initial, front points to slot 6 (A) and rear points to slot 2 (C). In (b) Addition, front points to slot 6 (A) and rear points to slot 3 (D). In (c) Deletion, front points to slot 5 (B) and rear points to slot 3 (D). Implementation of Circular Queue Operations When the array is viewed as a circle, each array position has a next and a previous position. The position next to MAX-QUEUE-SIZE -1 is 0, and the position that precedes 0 is MAX-QUEUE-SIZE -1. When the queue rear is at MAX_QUEUE_SIZE-1, the next element is inserted at position 0. In circular queue, the variables front and rear are moved from their current position to the next position in clockwise direction. This may be done using code <pre> if (rear == MAX_QUEUE_SIZE-1) rear = 0; else rear++; </pre> Add to a circular queue 2.5M <pre> void addq(element item) { rear = (rear +1) % MAX_QUEUE_SIZE; if (front == rear) queueFull(); queue [rear] = item; } </pre> Delete from a circular queue 2.5M <pre> element deleteq() { /* remove front element from the queue */ element item; if (front == rear) return queueEmpty(); /* return an error key */ front = (front+1)% MAX_QUEUE_SIZE; return queue[front]; } </pre> Display 2.5M Steps: If queue is empty, display message.	10	2	L2

Start from front, print elements until you reach rear.
Use $(i + 1) \% \text{MAX}$ to move circularly.

$$A = \begin{bmatrix} 0 & 0 & 3 & 0 & 4 \\ 0 & 0 & 5 & 7 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 6 & 0 & 0 \end{bmatrix}$$

Discuss Sparse matrix? Give the triplet form of a given matrix and find its transpose
SPARSE MATRIX REPRESENTATION 2M

- We can classify uniquely any element within a matrix by using the triple. Therefore, we can use an array of triples to represent a sparse matrix.
- Sparse matrix contains many zero entries.
- When a sparse matrix is represented as a 2-dimensional array, we waste space. For ex, if 100×100 matrix contains only 100 entries then we waste 9900 out

	col 0	col 1	col 2
row 0	-27	3	4
row 1	6	82	-2
row 2	109	-64	11
row 3	12	8	9
row 4	48	27	47

(a)

	col 0	col 1	col 2	col 3	col 4	col 5
row 0	15	0	0	22	0	-15
row 1	0	11	3	0	0	0
row 2	0	0	0	-6	0	0
row 3	0	0	0	0	0	0
row 4	91	0	0	0	0	0
row 5	0	0	28	0	0	0

(b)

4M

$$\begin{matrix} & 0 & 1 & 2 & 3 & 4 \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 0 & 3 & 0 & 4 \\ 0 & 0 & 5 & 7 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 6 & 0 & 0 \end{bmatrix} \end{matrix}$$

Triplet

4	5	6
0	2	3
0	4	4
1	2	5
1	3	7
3	1	2
3	2	6

Transpose

5	4	6
1	3	2
2	0	3
2	1	5
2	3	6
3	1	7
4	0	4

