

## IAT-1 SOLUTION ADE

**Answer COMPULSARILY**

Simplify the following expression using Quine –McCluskey's method

$$F(A,B,C,D) = \sum m(0,1,2,3,10,11,12,13,14,15).$$

**Solution : Step 1 :** List all minterms in the binary form as shown in table (column (a)).

**Step 2 :** Arrange the minterms according to no. of 1s, as shown in table (column (b)).

Minterms	Binary representation	Minterms	Binary representation
$m_0$	0000	$m_0$	0000 ✓
$m_1$	0001	$m_1$	0001 ✓
$m_2$	0010	$m_2$	0010 ✓
$m_3$	0011	$m_3$	0011 ✓
$m_{10}$	1010	$m_{10}$	1010 ✓
$m_{11}$	1011	$m_{12}$	1100 ✓
$m_{12}$	1100	$m_{11}$	1011 ✓
$m_{13}$	1101	$m_{13}$	1101 ✓
$m_{14}$	1110	$m_{14}$	1110 ✓
$m_{15}$	1111	$m_{15}$	1111 ✓

Step 3 : Compare each binary no. with every term in the adjacent next high category and if they differ only by one position, put a check mark and copy the term in the next column with '-' in the position that they differ.

Step 4 : Apply the same process described in step 3 for the resultant column and continue these cycles until a single pass through cycle yields us further elimination of literals.

Minterms	Binary representation	Minterms	Binary representation
0, 1	0 0 0 - ✓	0, 1, 2, 3	0 0 - -
0, 2	0 0 - 0 ✓	2, 3, 10, 11	- 0 1 -
1, 3	0 0 - 1 ✓	10, 11, 14, 15	1 - 1 -
2, 3	0 0 1 - ✓	12, 13, 14, 15	1 1 - -
2, 10	- 0 1 0 ✓		
3, 11	- 0 1 1 ✓		
10, 11	1 0 1 - ✓		
10, 14	1 - 1 0 ✓		
12, 13	1 1 0 - ✓		
12, 14	1 1 - 0 ✓		
11, 15	1 - 1 1 ✓		
13, 15	1 1 - 1 ✓		
14, 15	1 1 1 - ✓		

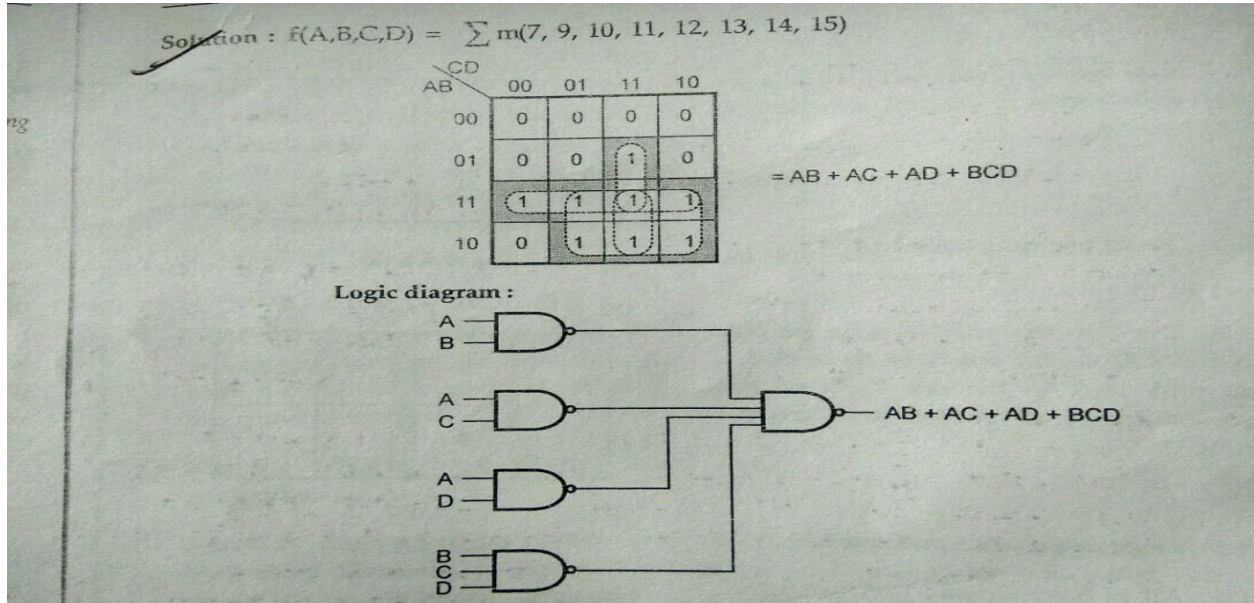
Step 5 : Select the minimum number of prime implicants which must cover all the minterms.

Prime Implicants	$m_0$ (col 1)	$m_1$ (col 2)	$m_2$ (col 3)	$m_3$ (col 4)	$m_{10}$ (col 5)	$m_{11}$ (col 6)	$m_{12}$ (col 7)	$m_{13}$ (col 8)	$m_{14}$ (col 9)	$m_{15}$ (col 10)
$\overline{A}B$ 12, 13, 14, 15 ✓							⊙	⊙	⊙	⊙
$\overline{A}C$ 10, 11, 14, 15 ✓					⊙	⊙			⊙	⊙
$\overline{B}C$ 2, 3, 10, 11			•	•	•	•				
$\overline{A}\overline{B}$ 0, 1, 2, 3 ✓	⊙	⊙	⊙	⊙						

$$F(A, B, C, D) = \overline{A}\overline{B} + \overline{A}C + \overline{B}C$$

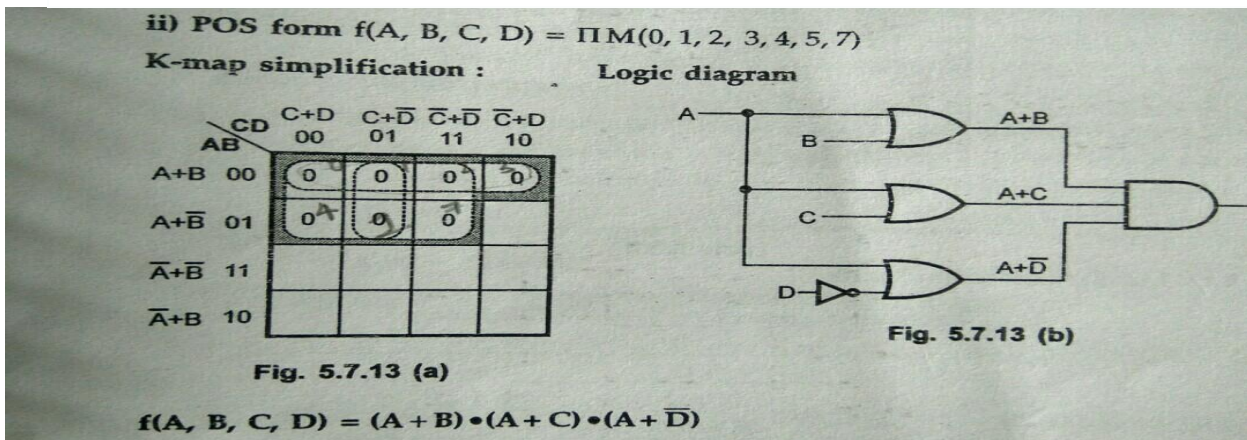
2 A) Obtain the simplified Boolean equation by using K-map method and express it in SOP form. Realize logic circuit by using NAND gates only

$$F(A,B,C,D) = \sum m(7, 9, 10, 11, 12, 13, 14, 15).$$

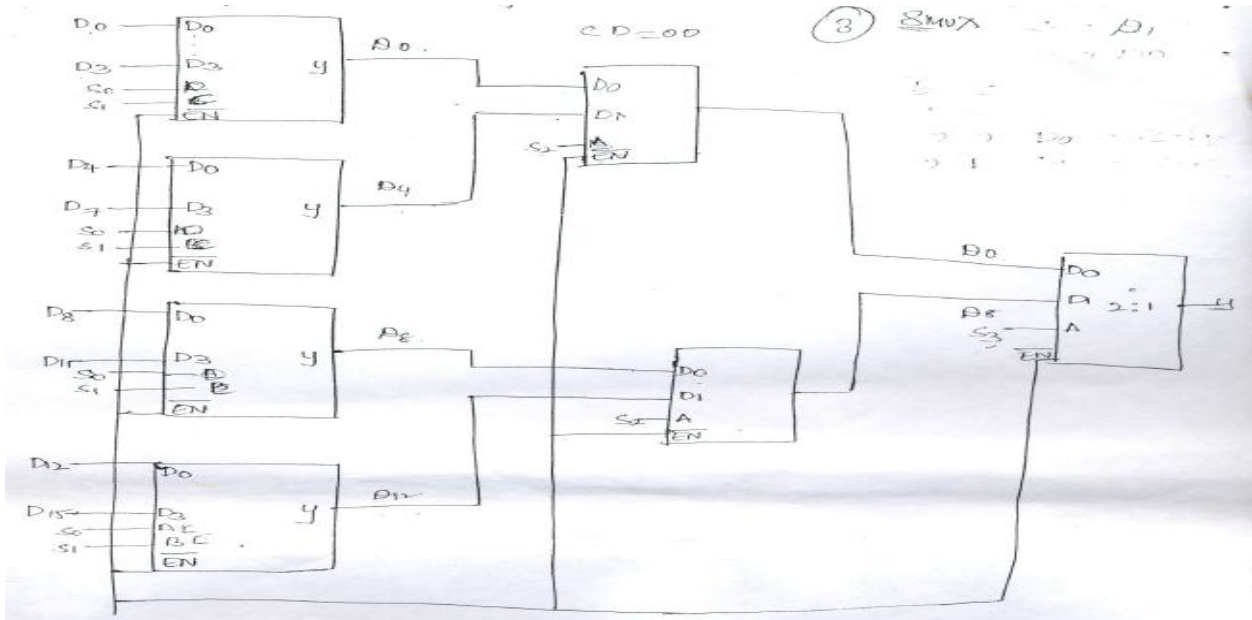


2 B) Find the minimal product using K-map for the Boolean function

$$F(A,B,C,D) = \sum m(6,7,9,10,13) + d(1,4,5,11).$$



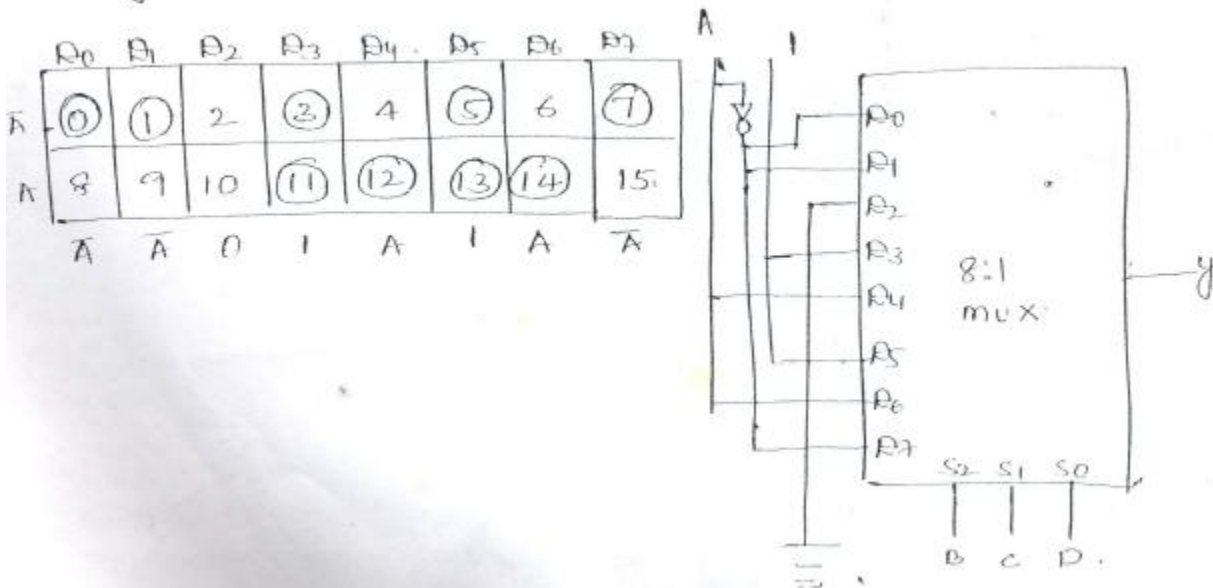
4 (a) Construct a 16:1 multiplexer using 4 to 1 and 2 to 1 multiplexers.



(b) Implement the given Boolean function by using 8:1 multiplexer.

$$F(A,B,C,D) = \sum m(0,1,3,5,7,11,12,13,14)$$

$f(A,B,C,D) = \sum(0,1,3,5,7,11,12,13,14)$   
using 8:1 mux



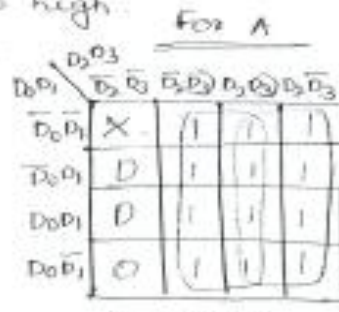
5 Design a 4 bit priority Encoder.

Priority Encoders:- The encoders will operate correctly prior that one & only one decimal i/p is high at any given time.

- Priority Encoder is a logic ckt that responds to just one i/p in accordance with some priority str, among all those i/p's may simultaneously go high.

4 i/p priority Encoder

Inputs				O/p's		
D <sub>0</sub>	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	A	B	V
0	0	0	0	x	x	0
1	0	0	0	0	0	1
x	1	0	0	0	1	1
x	x	1	0	1	0	1
x	x	x	1	1	1	1



$$A = D_3 + D_2$$



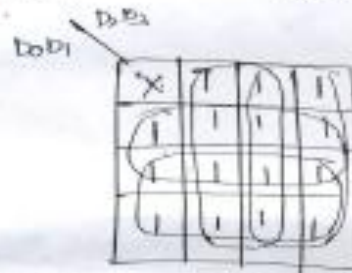
$$B = D_3 + D_1 \bar{D}_2 \bar{D}_3$$

From the T.T we get

$$A = D_3 + D_2$$

$$B = D_3 + D_1 \bar{D}_2 \bar{D}_3$$

$$V = D_0 + D_1 + D_2 + D_3$$

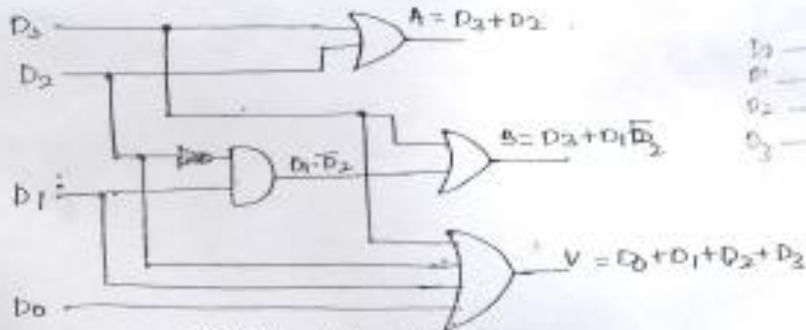


$$V = D_0 + D_1 + D_2 + D_3$$

There are 16 possible i/p combination.

$$A = \sum m(1, 2, 3, 5, 6, 7, 9, 10, 11, 13, 14, 15)$$

$$B = \sum m(1, 3, 4, 5, 7, 9, 11, 12, 13, 15)$$



4 bit priority Encoder

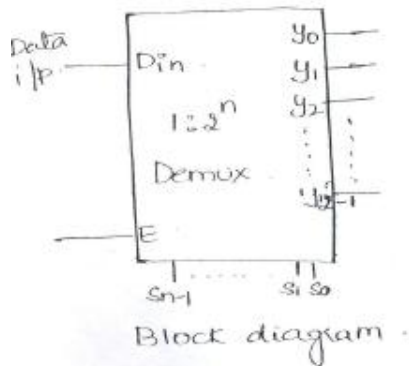


4 = 2<sup>2</sup> = o/p lines of block  
4 = 2<sup>2</sup>

6 (a) Explain the working of 1:4 demultiplexers.

Demultiplexer :- Is a ckt that receives a single i/p & distributes it over several o/p's. (i.e)  $2^n$  possible o/p lines. (1 to N).

- The selection of specific o/p line is controlled by the values of n selection lines.

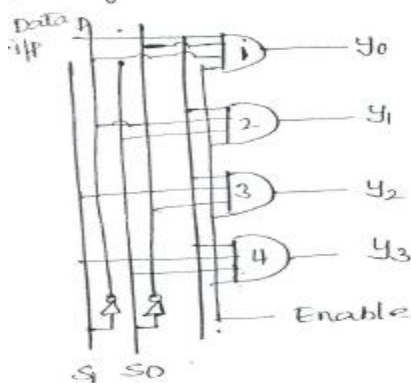


→ Types of Demultiplexer

i] 1:4 Demux :- The single i/p variable  $D_{in}$  has a path to all four o/p's, but the i/p info is directed to only one of the o/p lines depending on select i/p's.

Enable	$S_1$	$S_0$	$D_{in}$	$y_0$	$y_1$	$y_2$	$y_3$
0	X	X	X	0	0	0	0
1	0	0	0	0	0	0	0
1	0	0	1	1	0	0	0
1	0	1	0	0	0	0	0
1	0	1	1	0	1	0	0
1	1	0	0	0	0	0	0
1	1	0	1	0	0	1	0
1	1	1	0	0	0	0	0
1	1	1	1	0	0	0	1

Logic diagram

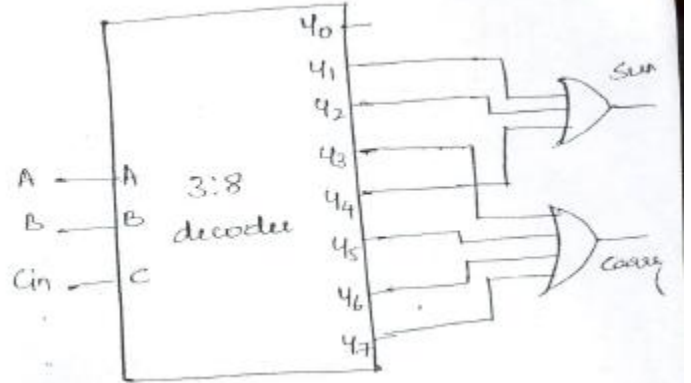


Selected lines		o/p.			
$S_1$	$S_0$	$y_0$	$y_1$	$y_2$	$y_3$
0	0	1	0	0	0
0	1	0	1	0	0
1	0	0	0	1	0
1	1	0	0	0	1

func table 4:1 MUX

(b) Implement a full adder using 3 to 8 decoders

Inputs	Outputs
A B Cin	Carry Sum
0 0 0	0 0
0 0 1	0 1
0 1 0	0 1
0 1 1	1 0
1 0 0	1 0
1 0 1	1 1
1 1 0	1 1
1 1 1	1 1



3 (a) Design a 2 bit magnitude comparator using gates.

2 bit Magnitude Comparator

Let the two 2 bit nos be  $A = A_1A_0$  &  $B = B_1B_0$

1. If  $A_1 = 1$  &  $B_1 = 0$  then  $A > B$  (A)

2. If  $A_1$  &  $B_1$  coincides &  $A_0 = 1$  &  $B_0 = 0$  then  $A > B$ .  
 $A > B: G = A_1 \bar{B}_1 + (A_1 \odot B_1) A_0 \bar{B}_0$

3. If  $A_1 = 0$  &  $B_1 = 1$  then  $A < B$  (B).

4. If  $A_1$  &  $B_1$  coincides &  $A_0 = 0$  &  $B_0 = 1$  then  $A < B$ .

5. If  $A_1$  &  $B_1$  coincides & if  $A_0$  &  $B_0$  coincide  $A = B$ .  
 $E = (A_1 \odot B_1)(A_0 \odot B_0)$

$A_1$	$A_0$	$B_1$	$B_0$	$A > B$	$A = B$	$A < B$
0	0	0	0	0	1	0
0	0	0	1	0	0	1
0	0	1	0	0	0	1
0	0	1	1	0	0	1
0	1	0	0	1	0	0
0	1	0	1	0	1	0
0	1	1	0	0	0	1
0	1	1	1	0	0	1
1	0	0	0	1	0	0
1	0	0	1	1	0	0
1	0	1	0	0	1	0
1	0	1	1	0	0	1
1	1	0	0	1	0	0
1	1	0	1	1	0	0
1	1	1	0	1	0	0
1	1	1	1	0	1	0

$A > B$

$A_1 A_0$	$B_1 B_0$	00	01	11	10
00	00	0	0	0	0
01	00	1	0	0	0
11	00	1	1	0	1
10	00	1	1	0	0

$$A > B = A_1 \bar{B}_1 + A_0 \bar{B}_1 \bar{B}_0 + A_1 A_0 \bar{B}_0$$

$A = B$

$A_1 A_0$	$B_1 B_0$	00	01	11	10
00	00	1	0	0	0
01	00	0	1	0	0
11	00	0	0	1	0
10	00	0	0	0	1

$$A = B = \bar{A}_1 \bar{A}_0 \bar{B}_1 \bar{B}_0 + \bar{A}_1 \bar{A}_0 \bar{B}_1 B_0 + A_1 A_0 \bar{B}_1 \bar{B}_0 + A_1 A_0 \bar{B}_1 B_0 + A_1 \bar{A}_0 \bar{B}_1 \bar{B}_0 + A_1 \bar{A}_0 \bar{B}_1 B_0 + \bar{A}_1 \bar{A}_0 B_1 \bar{B}_0 + \bar{A}_1 \bar{A}_0 B_1 B_0 + A_1 \bar{A}_0 B_1 \bar{B}_0 + A_1 \bar{A}_0 B_1 B_0 + \bar{A}_1 A_0 \bar{B}_1 \bar{B}_0 + \bar{A}_1 A_0 \bar{B}_1 B_0 + A_1 A_0 \bar{B}_1 \bar{B}_0 + A_1 A_0 \bar{B}_1 B_0$$

$$= \bar{A}_1 \bar{B}_1 (\bar{A}_0 \bar{B}_0 + A_0 B_0) + A_1 \bar{B}_1 (A_0 \bar{B}_0 + \bar{A}_0 B_0)$$

$$= (\bar{A}_0 \odot B_0) (A_1 \odot B_1)$$

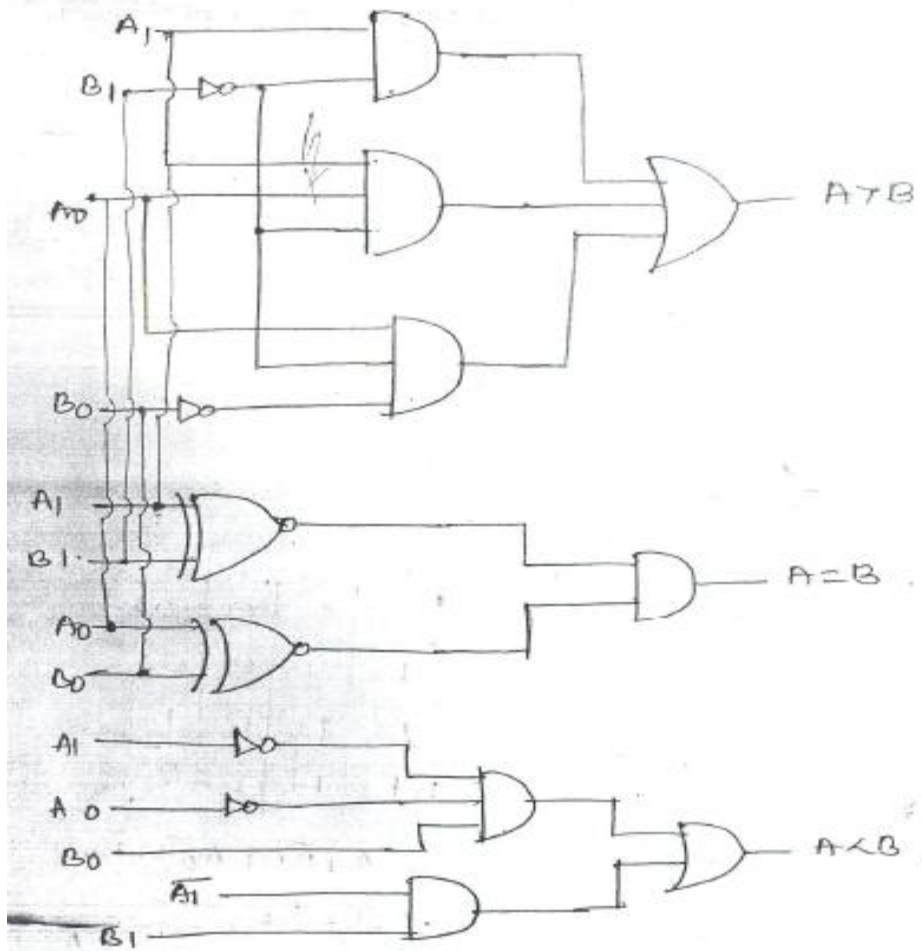
$A < B$

$A_1 A_0$	$B_1 B_0$	00	01	11	10
00	00	0	1	1	1
01	00	0	0	1	1
11	00	0	0	0	0
10	00	0	0	1	0

$$A < B = \bar{A}_1 B_1 + \bar{A}_0 B_1 B_0 + \bar{A}_1 \bar{A}_0 B_0$$

$A_1 A_0$   
 $A_3 A$   
 $\uparrow$   
 $\uparrow$





7. Implement a PLA circuit having 3 inputs, 3 product terms and two outputs for the given Boolean functions  $f_1 = \sum m(1,3,5)$  and  $f_2 = \sum m(5,6,7)$ .

(A) A combinational ckt is defined by the func.

$$F_1 = \sum m(1, 3, 5)$$

$$F_2 = \sum m(5, 6, 7)$$

Implement the ckt with a PLA having 3x3x2 ofp's

Sol<sup>n</sup> Kmap simplification.

For  $F_1$

	BC	$\overline{B}\overline{C}$	$\overline{B}C$	$B\overline{C}$	$BC$
$\overline{A}$			1	1	
A		1			

$$F_1 = \overline{B}C + \overline{A}C$$

For  $F_2$

	BC	$\overline{B}\overline{C}$	$\overline{B}C$	$B\overline{C}$	$BC$
$\overline{A}$					
A		1	1	1	

$$F_2 = AC + AB$$

$\therefore$  we get 4 product terms it cannot be used, examine grouping the maxterms

For  $F_1$

	BC	$\overline{B}\overline{C}$	$\overline{B}C$	$B\overline{C}$	$BC$
$\overline{A}$		0			0
A		0		0	0

$$F_1 = \overline{C} + AB$$

$$F_1 = \overline{B}\overline{C} + \overline{A}C + AB$$

For  $F_2$

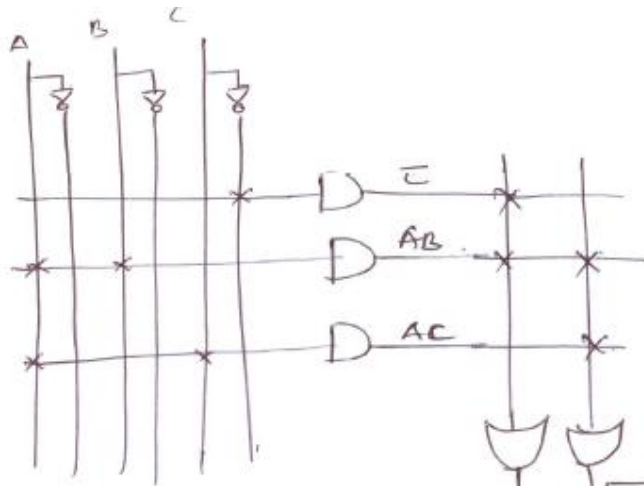
	BC	$\overline{B}\overline{C}$	$\overline{B}C$	$B\overline{C}$	$BC$
$\overline{A}$		0	0	0	0
A		0			

$$F_2 = \overline{B}\overline{C} + \overline{A}$$

needed terms

Funcs  $F_1$  &  $F_2$  have one common product term having 3 product terms.

P.T	ofp's			ofp's	
	A	B	C	$F_1$	$F_2$
$\overline{C}$	-	-	0	1	-
AB	1	1	-	1	1
AC	1	-	1	-	1
				<u>C</u>	<u>T</u>



link open  
to get  
complemented  
o/p.

link closed  
to uncomplemented  
o/p.

1 → Ex-OR gate is programmed  
to invert the  $f_{in}^c$