

For the channel diagram shown in Fig Q1 calculate $H(X)$, $H(Y)$, $H(X,Y)$, $H(X|Y)$, $H(Y|X)$ and $I(X;Y)$ if $P(x_1) = P(x_2) = P(x_3)$

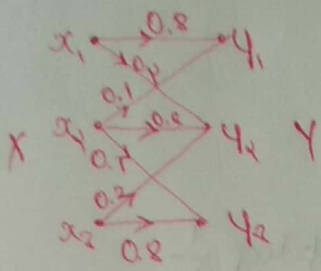


Fig Q1. Noise diagram

The channel matrix

$$\text{Sol}^n \quad P(Y|X) = \begin{matrix} & \begin{matrix} y_1 & y_2 & y_2 \end{matrix} \\ \begin{matrix} x_1 \\ x_2 \\ x_3 \end{matrix} & \begin{bmatrix} 0.8 & 0.2 & 0 \\ 0.1 & 0.9 & 0.1 \\ 0 & 0.2 & 0.8 \end{bmatrix} \end{matrix}$$

$$P(x_1) = P(x_2) = P(x_3) = \frac{1}{3}$$

Joint Probability matrix

$$P(X,Y) = \begin{matrix} & \begin{matrix} y_1 & y_2 & y_2 \end{matrix} \\ \begin{matrix} x_1 \\ x_2 \\ x_3 \end{matrix} & \begin{bmatrix} (0.8)/3 & (0.2)/3 & 0 \\ (0.1)/3 & (0.9)/3 & (0.1)/3 \\ 0 & (0.2)/3 & (0.8)/3 \end{bmatrix} \end{matrix}$$

$$P(y_1) = \frac{0.9}{3} \cdot \frac{1}{3} + \frac{0.2}{3} \cdot \frac{1}{3} = \frac{0.9}{3}$$

$$\begin{aligned} \text{i. } H(X) &= \sum_{i=1}^3 P(x_i) \log_2 \frac{1}{P(x_i)} \\ &= 3 \left(\frac{1}{3} \log_2 3 \right) \\ &= \underline{\underline{1.5849 \text{ bits/sym}}} \end{aligned}$$

$$\begin{aligned} \text{(ii) } H(Y) &= \sum_{j=1}^2 P(y_j) \log_2 \frac{1}{P(y_j)} \\ &= 0.9 \log_2 \frac{3}{0.9} + \frac{1.9}{3} \log_2 \frac{3}{1.9} \\ &= \underline{\underline{1.5709 \text{ bits/sym}}} \end{aligned}$$

$$\begin{aligned} \text{(iii) } H(X,Y) &= \sum_{i=1}^3 \sum_{j=1}^2 P(x_i, y_j) \log_2 \frac{1}{P(x_i, y_j)} \\ &= 2 \left(\frac{0.8}{3} \log_2 \frac{3}{0.8} \right) + 2 \left(\frac{0.2}{3} \log_2 \frac{3}{0.2} \right) \\ &\quad + 2 \left(\frac{0.1}{3} \log_2 \frac{3}{0.1} \right) \\ &= \underline{\underline{2.3746 \text{ bits/sym}}} \end{aligned}$$

$$\begin{aligned} \text{(iv) } H(Y|X) &= H(X,Y) - H(X) \\ &= 2.3746 - 1.5849 \\ &= \underline{\underline{0.7897 \text{ bits/sym}}} \end{aligned}$$

$$\begin{aligned} \text{(v) } H(X|Y) &= H(X,Y) - H(Y) \\ &= 2.3746 - 1.5709 \\ &= \underline{\underline{0.8037 \text{ bits/sym}}} \end{aligned}$$

$$\begin{aligned} \text{(vi) } I(X;Y) &= H(X) - H(X|Y) \\ &= 1.5849 - 0.8037 \\ &= \underline{\underline{0.7812 \text{ bits/sym}}} \end{aligned}$$

Derive an expression for the channel capacity of binary symmetric channel.

The use is as shown below

Channel matrix

The joint probability matrix

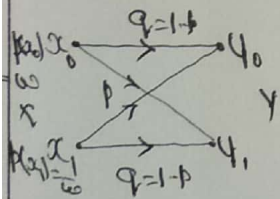


Fig. Channel diagram

$$P(Y|X) = \begin{bmatrix} q & p \\ p & q \end{bmatrix}$$

$$P(X,Y) = \begin{bmatrix} wq & wp \\ wp & wq \end{bmatrix}$$

$$\therefore P(y_0) = wq + wp$$

$$P(y_1) = wp + wq$$

$$\text{w.k.t. } C = \max_{P(x)} I(X;Y) \text{ bits/sec}$$

$$I(X;Y) = H(Y) - H(Y|X)$$

$$H(Y) = \sum_{j=0}^1 P(y_j) \log_2 \frac{1}{P(y_j)}$$

$$= (wq + wp) \log_2 \frac{1}{(wq + wp)} + (wp + wq) \log_2 \frac{1}{(wp + wq)}$$

Conditional entropy

Conditional entropy $H(Y|X)$ is

$$H(Y|X) = \sum_{j=0}^1 \sum_{i=0}^1 p(x_i, y_j) \log_2 \frac{1}{p(y_j|x_i)}$$

$$= \omega q \log_2 \frac{1}{q} + \bar{\omega} q \log_2 \frac{1}{q} + \omega p \log_2 \frac{1}{p} + \bar{\omega} p \log_2 \frac{1}{p}$$

$$= (\omega + \bar{\omega}) q \log_2 \frac{1}{q} + (\omega + \bar{\omega}) p \log_2 \frac{1}{p}$$

$$= p \log_2 \frac{1}{p} + q \log_2 \frac{1}{q}$$

$$\therefore I(X; Y) = \left[(\omega q + \bar{\omega} p) \log_2 \frac{1}{p(\omega q + \bar{\omega} p)} + (\omega p + \bar{\omega} q) \log_2 \frac{1}{q(\omega p + \bar{\omega} q)} \right] - \left[p \log_2 \frac{1}{p} + q \log_2 \frac{1}{q} \right]$$

$$\max \{I(X; Y)\} \Rightarrow \omega = \bar{\omega} = \frac{1}{2}$$

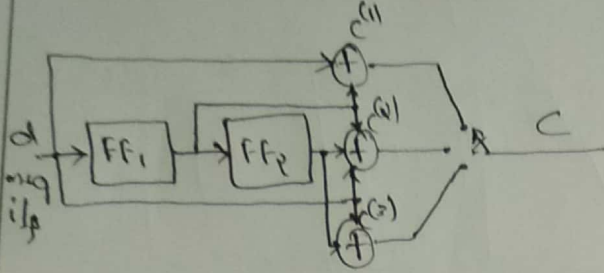
$$\therefore \max I(X; Y) = \left\{ \frac{1}{2} \log_2 2 + \frac{1}{2} \log_2 2 \right\} - p \log_2 \frac{1}{p} + q \log_2 \frac{1}{q}$$

$$I(X; Y) = 1 - (p \log_2 \frac{1}{p} + q \log_2 \frac{1}{q}) \text{ bits/sec}$$

$$\therefore C = \underline{\underline{1 - (p \log_2 \frac{1}{p} + q \log_2 \frac{1}{q})}} \text{ bits/sec}$$

4. Consider the (3,1,4) convolutional code with $g^{(1)} = (110)$, $g^{(2)} = (111)$, $g^{(3)} = (101)$

(a) Convolution Encoder circuit



Encoding Circuit.

(b) obtain the generator matrix

$$[G]_{L \times (Lm)} = \begin{bmatrix} g_1^{(1)} & g_1^{(2)} & g_1^{(3)} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & g_2^{(1)} & g_2^{(2)} & g_2^{(3)} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & g_3^{(1)} & g_3^{(2)} & g_3^{(3)} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_4^{(1)} & g_4^{(2)} & g_4^{(3)} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_5^{(1)} & g_5^{(2)} & g_5^{(3)} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_6^{(1)} & g_6^{(2)} & g_6^{(3)} \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_7^{(1)} & g_7^{(2)} & g_7^{(3)} \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_8^{(1)} & g_8^{(2)} & g_8^{(3)} \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_9^{(1)} & g_9^{(2)} & g_9^{(3)} \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & g_{10}^{(1)} & g_{10}^{(2)} & g_{10}^{(3)} \\ 0 & g_{11}^{(1)} & g_{11}^{(2)} & g_{11}^{(3)} \\ 0 & g_{12}^{(1)} & g_{12}^{(2)} & g_{12}^{(3)} \\ 0 & g_{13}^{(1)} & g_{13}^{(2)} & g_{13}^{(3)} \\ 0 & g_{14}^{(1)} & g_{14}^{(2)} & g_{14}^{(3)} \\ 0 & g_{15}^{(1)} & g_{15}^{(2)} & g_{15}^{(3)} \\ 0 & g_{16}^{(1)} & g_{16}^{(2)} & g_{16}^{(3)} \\ 0 & g_{17}^{(1)} & g_{17}^{(2)} & g_{17}^{(3)} \\ 0 & g_{18}^{(1)} & g_{18}^{(2)} & g_{18}^{(3)} \end{bmatrix}$$

5. Encode $d = 1011$ using time domain approach

Given $d = 1011 \therefore L = 4$ bits

o/p bits = $n(L+m) = 3(4+3) = 18$ bits

$$C = [d] * [G]$$

$$[C] = [1011] \begin{bmatrix} 111 & 110 & 011 & 000 & 000 & 000 \\ 000 & 111 & 110 & 011 & 000 & 000 \\ 000 & 000 & 111 & 110 & 011 & 000 \\ 000 & 000 & 000 & 111 & 110 & 011 \end{bmatrix}$$

$$= [111000 \quad 110100 \quad 100001 \quad 001101 \quad 011000]$$

Transform domain approach.

$$C(x) = c^{(1)}(x^n) + x c^{(2)}(x^n) + x^2 c^{(3)}(x^n)$$

$$d(x) = 1 + x^1 + x^3$$

$$g^{(1)}(x) = 1 + x$$

$$g^{(2)}(x) = 1 + x + x^2$$

$$g^{(3)}(x) = 1 + x^2$$

Output from top adder

$$\begin{aligned}
 c^{(1)}(x) &= d(x) * q^{(1)}(x) \\
 &= (1 + x^2 + x^3) * (1 + x) \\
 &= 1 + x^2 + x^3 + x + x^2 + x^4 \\
 &= \underline{\underline{1 + x + x^3 + x^4}}
 \end{aligned}$$

$$\therefore c^{(1)}(x^3) = c^{(1)}(x^3)$$

$$c^{(1)}(x^3) = \underline{\underline{1 + x^2 + x^6 + x^4}}$$

Output from middle adder

$$\begin{aligned}
 c^{(2)}(x) &= d(x) * q^{(2)}(x) \\
 &= (1 + x^2 + x^3) (1 + x + x^4) \\
 &= 1 + x^2 + x^3 + x + x^2 + x^4 + x^3 + x^4 + x^5 \\
 &= \underline{\underline{1 + x + x^5}}
 \end{aligned}$$

$$\therefore c^{(2)}(x^3) = \underline{\underline{x + x^9 + x^{15}}}$$

$$x c^{(2)}(x^3) = \underline{\underline{x^2 + x^4 + x^{16}}}$$

Output from bottom adder

$$\begin{aligned}
 c^{(3)}(x) &= d(x) * q^{(3)}(x) \\
 &= (1 + x^2 + x^3) (1 + x^4) \\
 &= 1 + x^2 + x^3 + x^4 + x^4 + x^5 \\
 &= \underline{\underline{1 + x^2 + x^4 + x^5}}
 \end{aligned}$$

$$\therefore c^{(3)}(x^3) = 1 + x^9 + x^{12} + x^{15}$$

$$x^2 c^{(3)}(x^3) = \underline{\underline{x^4 + x^{11} + x^{17}}}$$

$$\therefore c(x) = 1 + x + x^2 + x^3 + x^4 + x^5 + x^6 + x^7 + x^8 + x^9 + x^{10} + x^{11} + x^{12} + x^{13} + x^{14} + x^{15} + x^{16} + x^{17}$$

$$c = [1111100100001101011]$$

Consider the binary Convolutional Encoder shown in Fig Q6.

message D=101 using code tree.

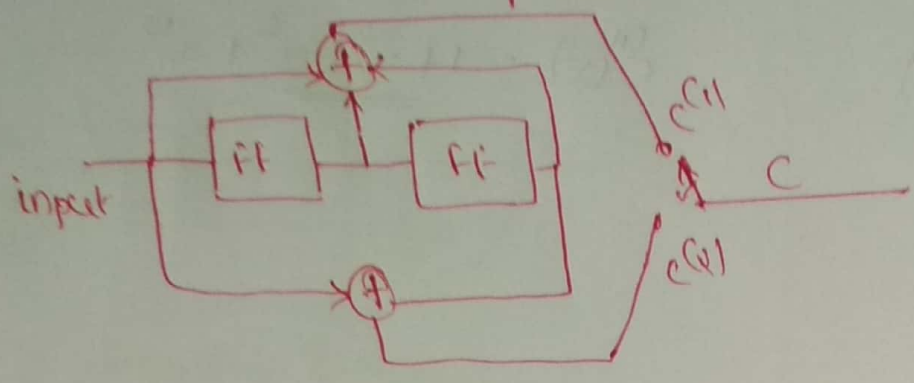


Fig Q6. (2,1,2) Convolutional Encoder

No of states possible = $2^m = 2^2 = 4$ states

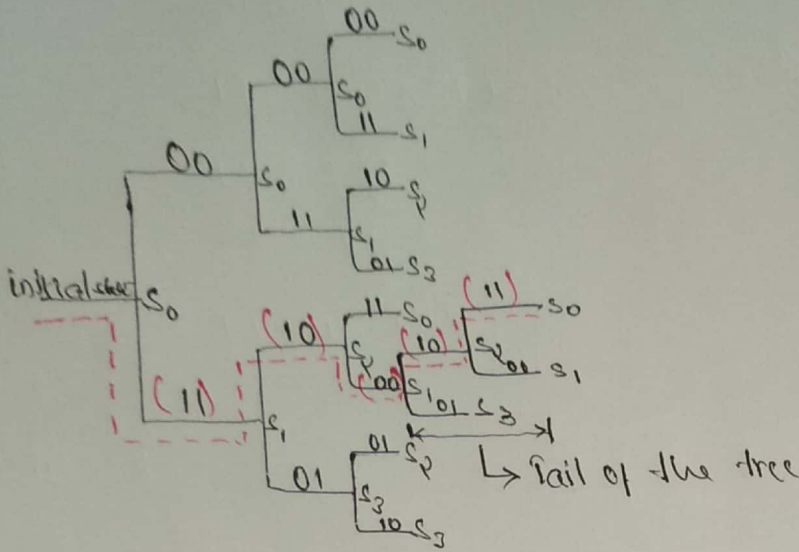
state table:

state	s_0	s_1	s_2	s_3
BD	00	10	01	11

State transition table

PS	BD	i/p	NS	BD	d_{k1}	d_{k-1}	d_{k-2}	$c^{(1)}$	$c^{(2)}$
s_0	00	0	s_0	00	0	0	0	0	0
		1	s_1	10	1	0	0	1	1
s_1	10	0	s_2	01	0	1	0	1	0
		1	s_3	11	1	1	0	0	1
s_2	01	0	s_0	00	0	0	1	1	1
		1	s_1	10	1	0	1	0	0
s_3	11	0	s_2	01	0	1	1	0	0
		1	s_3	11	1	1	1	0	0

Code tree



Msg → 1 0 1 - -

Encoded sequence → [11 10 00 10 11]

2. Find the capacity of the discrete channel shown in Fig Q7 using Shannon's theorem.

Solⁿ: Channel capacity $C = \sum_{i=1}^3 \log_2 (1 - P_i)$ bits/sec

$$\begin{bmatrix} 0.2 & 0.1 & 0.1 \\ 0.4 & 0.6 & 0.4 \\ 0.4 & 0.4 & 0.6 \end{bmatrix} \begin{bmatrix} P_1 \\ P_2 \\ P_3 \end{bmatrix} = \begin{bmatrix} 0.2 \log_2 \frac{1}{0.2} + 0.1 \log_2 \frac{1}{0.1} \\ 0.4 \log_2 \frac{1}{0.4} + 0.6 \log_2 \frac{1}{0.6} \\ 0.4 \log_2 \frac{1}{0.4} + 0.6 \log_2 \frac{1}{0.6} \end{bmatrix}$$

$$0.2P_1 + 0.1P_2 + 0.1P_3 = 0.9719$$

$$0.4P_1 + 0.6P_2 + 0.4P_3 = 1.371$$

$$0.4P_1 + 0.4P_2 + 0.6P_3 = 1.371$$

$$\therefore C = \log_2 \left[2^{-0.9772} + 2^{-1.5703} + 2^{-1.5707} \right]$$

$$= \underline{\underline{0.359 \text{ bits/sec}}}$$

