

1a. What is a data structure? What is a linear and non-linear data structure. What is an algorithm?
Explain the criteria an algorithm must satisfy.

Definition (Data Structure) :-

A logical or mathematical model of a particular organisation of data. Choice is made based on

- (i) Relationships of Data and (Nature of Data)
- (ii) Effectively processing the Data.

Data Structure is a particular way of organisation (organising) data in computer so that it can be used effectively. Any data structure is designed to organise the data to suite a specific purpose so that it can be accessed and worked in appropriate way.

Linear DS:

Traverses the data elements sequentially. One one element can directly be reached. Eg: Arrays, linked list.

Non-linear DS:

A data is connected to several other data items so that a given data structure has the possibility to reach one or more data items. Eg: Trees, Graphs.

Algorithm Specification

Definition: An algorithm is a finite set of instructions that if followed accomplishes a particular task.

All the algorithms must satisfy the following criteria:

① Input: There are zero or more quantities that are externally supplied.

② Output: At least one quantity is produced.

③ Definiteness: Each instruction is clear and unambiguous.

④ Finiteness: For all cases of the problem, the algorithm terminates after a finite number of steps.

⑤ Effectiveness: Every instruction must be basic enough to be carried out and must be feasible.

1b. Write about dynamic memory allocation functions.

malloc(): allocates a continuous block of memory which is not cleared. If enough continuous memory is available and returns a pointer (void *) to the block of memory otherwise returns null.
^{in eq:} p holds the address of the first element of the block of 2n bytes.

```
void *calloc(size_t nmemb, size_t size);  
           ↗ elements ↘ size bytes
```

calloc() allocates memory for an array of n members (nmemb) elements of size bytes each and returns a pointer to the allocated memory and the memory is set to zero.

```
Eq:- char *cp;  
      cp = calloc(100, sizeof(char));
```

```
void * realloc (void * ptr, size_t size);
```

↳ newsize.

[ptr - pointer ~~to~~ pointing to the memory allocated by malloc or calloc().

- realloc() changes the size of memory block pointed to by ptr to size bytes.
- The contents will be unchanged to the minimum of old and new sizes.
- Newly allocated memory will be uninitialised if ptr is NULL, the call is equivalent to malloc of size.
- If the size is equivalent to NULL, then the call is equivalent to free(ptr), otherwise it must have been returned by an earlier call to malloc, calloc, realloc.
- realloc attempts to change size of previous allocated block of memory

```
void free (void * ptr);
```

free() takes a pointer as a argument and free the memory to which the pointer refers.

NOTE :

Once of array of memory is freed, it is improper to use it.

3. a. Consider two polynomials, $A(x)=4x^5+3x^4+5$ and $B(x)=x^4+10x^2+1$

Show diagrammatically how these two polynomials can be stored in a 1-D array. Also give its C representation

3.b. Write a function to add two polynomials

Since each term has two values — one coefficient, other expo

we use structures.
`#define MAX_TERMS 100`
`typedef struct`

```

float coef;
int expon;
} polynomial;

```

→ userdefined type

```

polynomial terms[MAX_TERMS]
int avail = 0;

```

↓ type ↓ name ↓ size

Eg: $A(x) = 4x^5 + 3x^2 - 2$

	term[0]	term[1]	term[2]
coef	4	3	-2
expon	5	2	0

Adding two polynomials:

Eg: $A(x) = 4x^5 + 3x^3 + 5x^2 + 6$

$A(x) = 4x^5 + 3x^3 + 8x + 6$

$B(x) = 2x^5 + 8x$

$B(x) = 2x^5 + 5x^2$

$D(x) = A(x) + B(x)$

$= 6x^5 + 3x^3 + 5x^2 + 8x + 6$

Representing in 1-D Array :-

	startA	finishA	startB	finishB	avail
coef	4	3	8	6	5
expon	5	3	1	0	2
	0	1	2	3	4
	A(x)			B(x)	

```

void attach(float coefficient, int exponent)
{
    if (avail >= MAX_TERMS)
        exit(0);
    terms[avail].coef = coefficient;
    terms[avail].expon = exponent;
    avail++; // increment the index
}

```

```

void addpolynomial(int startA, int finishA, int startB, int
                  int *startD, int *finishD
                  indices of A      indices
                  indices of D.
{
    float coefficient;
    *startD = avail;
    while (startA <= finishA && startB <= finishB)
    {
        if (terms[startA].expon == terms[startB].expon)
            /* Add coefficient
            coefficient = terms[startA].coef + terms[startB].
            if (coefficient)
                attach(coefficient, terms[startA].expon);
            startA++;
            startB++;
        }
        else if (terms[startA].expon > terms[startB].expon)
            /* copy term from B */
            attach(terms[startB].coef, terms[startB].expon);
            startB++;
        }
        else
        {
            attach(terms[startA].coef, terms[startA].expon);
            startA++;
        }
    }
    /* Add in the remaining terms of A */
    for ( ; startA <= finishA; startA++)
        attach(terms[startA].coef, terms[startA].expon);
    /* Add in the remaining terms of B */
    for ( ; startB <= finishB; startB++)
        attach(terms[startB].coef, terms[startB].expo);
    *finishD = avail - 1;
}

```

4a. Write Knuth Morris pattern matching algorithm

4b. Apply Knuth Morris pattern matching algorithm the same to search pattern 'abcdabcy' in the text 'abcxabcdabxabcdabcdabcy'.

```

int kmpm(char *string, char *pattern)
{
    int i=0, j=0;
    int sl, pl;
    sl = strlen(s);
    pl = strlen(p);
    while (i < sl && j < pl)
    {
        if (s[i] == p[j])
            i++, j++;
        else
        {
            j = FF(j-1);
            if (j == 0)
                i++;
        }
    }
    return ((j == pl) ? (i - pl) : -1);
}

```

Pattern	a	b	c	d	a	b	c	y
F.F	0	0	0	0	1	2	3	0

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22
String:	a	b	c	x	a	b	c	d	a	b	x	a	b	c	d	a	b	c	d	a	b	c	y
Pattern	a	b	c	d	a	b	c	y															
					a	b	c	d	a	b	x	y											
									a	b	x	d	a	b	c	y							
												a	b	c	d	a	b	c	x	y			
																				a	b	c	y

5a. Write a fast transpose algorithm to transpose the given sparse matrix

```

void fastTranspose (term a[], term b[])
{ /* the transpose of a is placed in b */
  int rowTerms [MAX_col], startingPos [MAX_col],
  int i, j, numCols = a[0].col, numTerms = a[0].value;
  b[0].row = numCols; b[0].col = a[0].row;
  b[0].value = numTerms;
  if (numTerms > 0) /* nonzero matrix */
  {
    for (i=0; i < numCols; i++)
      rowTerms[i] = 0;

    for (i=1; i <= numTerms; i++)
      rowTerms[a[i].col]++;
    startingPos[0] = 1;
    for (i=1; i < numCols; i++)
      startingPos[i] = startingPos[i-1] + rowTerms[i-1];
    for (i=1; i <= numTerms; i++) {
      j = startingPos[a[i].col]++; b[j].row = a[i].col;
      b[j].col = a[i].row; b[j].value = a[i].value;
    }
  }
}

```

5b. Express the given sparse matrix as triplets and find its transpose.

sparse matrix :- triple				Transpose of sparse matrix :- tri			
	row	col	value		row	col	value
a[0]	6	6	8	b[0]	6	6	8
a[1]	0	0	15	b[1]	0	0	15
a[2]	0	3	22	b[2]	0	4	91
a[3]	0	5	-15	b[3]	1	1	11
a[4]	1	1	11	b[4]	2	1	3
a[5]	1	2	3	b[5]	2	5	28
a[6]	2	3	-6	b[6]	3	0	22
a[7]	4	0	91	b[7]	3	2	-6
a[8]	5	2	28	b[8]	5	0	-15

6a. Write a recursive algorithm for tower of Hanoi and ackerman's function


```

void #if main()
{
    int n=3; //Read n
    toh(n, 'a', 'b', 'c');
}

void toh(int n, char a, char b, char c)
{
    if(n==0)
        return;
    else
    {
        toh(n-1, 'a', 'c', 'b');
        printf("Move the disks from %c to %c, a, b);", a, b);
        toh(n-1, 'c', 'b', 'a');
    }
}

```

```

int ackerman(int m, int n)
{
    if (m==0) return n+1;
    else if (n==0) return ackerman(m-1, 1);
    else return ackerman(m-1, ackerman(m, n-1));
}

void main()
{
    int m, n;
    // Read m and n
    printf("A(%d, %d) = %d ; m, n, ackerman(m, n));", m, n, ackerman(m, n));
}

```

6b. Convert the infix expression to postfix form

i) $(a+b)*d+e/(f+a*d)+c$, ii) $((a/b)+(d*c))-(a*c)$

$ab+d*e+fad*+c/$

$ab/dc*+ac*-$

7. Write an algorithm to implement a stack using dynamic array whose initial capacity is 1 and array doubling is used to increase the stack's capacity (that is dynamically relocate twice the memory) whenever an element is added to a full stack. Implement the operations-push, pop and display

```

typedef struct
{
    int ele;
} STACK;

```

```

STACK *stack;
stack = malloc(sizeof(*stack))
int capacity = 1;
int top = -1;
void push(int item)
{
    if (top >= capacity - 1)
        stackFull();
    else
        stack[++top] = item;
}
void stackFull()
{
    stack = realloc(stack, 2 * capacity, * sizeof(*stack));
    capacity *= 2;
}
int pop()
{
    if (top == -1)
        stackEmpty();
    else
    {
        return stack[top--];
    }
}
void stackEmpty()
{
    printf("Stack Underflow");
    exit();
}

```

8. List the disadvantages of linear queue and explain how it is solved in circular queue. Give the algorithm to implement a circular queue with suitable example.

```

#define SIZE 8
typedef struct
{
    int ele;
} QUEUE;
QUEUE cqueue[SIZE];
int front = rear = -1;
void addcq (QUEUE item)
{
    if (front == (rear + 1) % SIZE)
        cqueueFull();
    else
    {
        rear = (rear + 1) % SIZE;
        cqueue[rear] = item;
        if (front == -1) /* First insert */
            front = front + 1;
    }
}

```

```

QUEUE deletecq()
{
    QUEUE item;
    if (front == -1)
        queueEmpty();
    else
    {
        item = cqueue[front];
        if (front == rear) /* Only one element */
            front = rear = -1;
        else
            front = (front + 1) % SIZE;
    }
    return item;
}

```

```

void queueFull()
{
    printf("Queue is Full");
    exit();
}

```

```

void queueEmpty()
{
    printf("Queue is Empty");
    exit();
}

```